

S O K E N D A I

NII



Formal Methods for Cyber-Physical Systems, **at Work**

Ichiro Hasuo

National Institute of Informatics & SOKENDAI
Research Director, ERATO MMSD Project, JST

Outline

- * Quality assurance of **cyber-physical systems**
 - * **Formal methods at work**,
coping with uncertainties
- * Introducing the ERATO MMSD project
- * Tech showcase:
“formal methods that are **down-scalable**”
 - * Monitoring
 - * Search-based testing
- * Applications, “**power of math**”

Cyber-Physical Systems : Control Theory and Formal Methods/Software Science

* Cyber-Physical System (CPS)

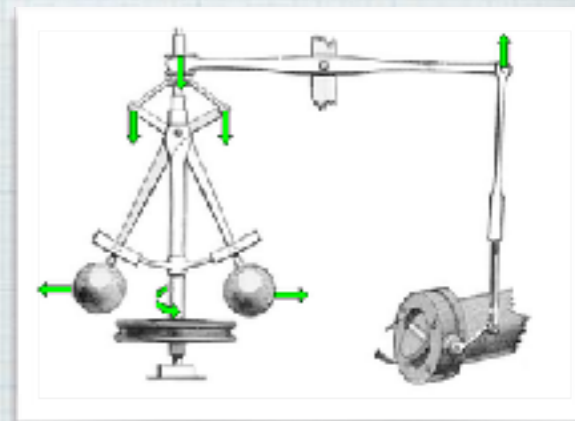
* “A mechanism that is controlled or monitored by computer-based algorithms, tightly integrated with the Internet and its users” (Wikipedia)

* Physical plant (continuous)

+

Digital control (discrete)

* In US: NSF Key Area of Research (2006-)



Cyber-Physical Systems : Control Theory and Formal Methods/Software Science

- * **Formal methods:**

- Logical proofs** for “correctness” of (discrete) programs

- * **Model checking**

- (SPIN, NuSMV, Uppaal, PRISM, ...) [Pnueli, Clarke, Emerson, Sifakis, ...]

- * **Theorem Proving**

- (Coq, Agda, ...) [Milner, Coquand, Leroy, Voevodsky, ...]



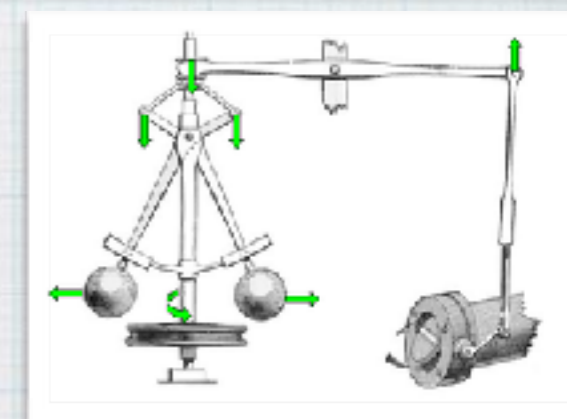
- * **Control Theory:**

- Analysis of continuous dynamics**

- * **Stability, Lyapunov function, ...**

- * Their **similarity** is widely recognized

- * e.g. HSCC, one of the main conferences of annual CPS Week



CPS Research, So Far (the V&V Aspect)



CPS
(esp. hybrid systems)

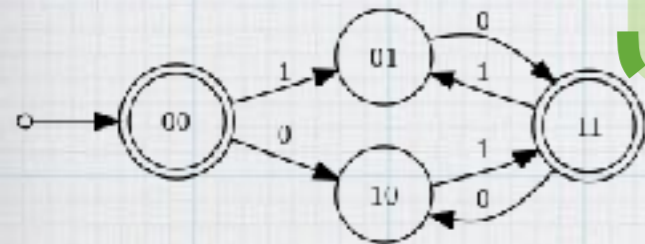
Analysis

Formal Methods

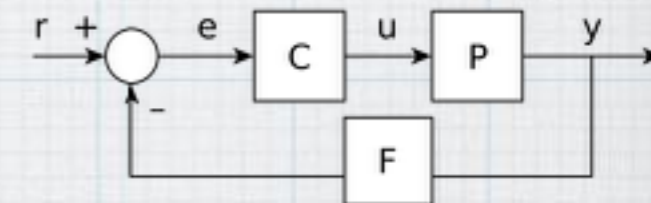
Control Theory

Collaboration

$$\square(p \Rightarrow \diamond q)$$



$$x' = f(x, u)$$



* Problem: **practical applicability**. Why?

* Scalability

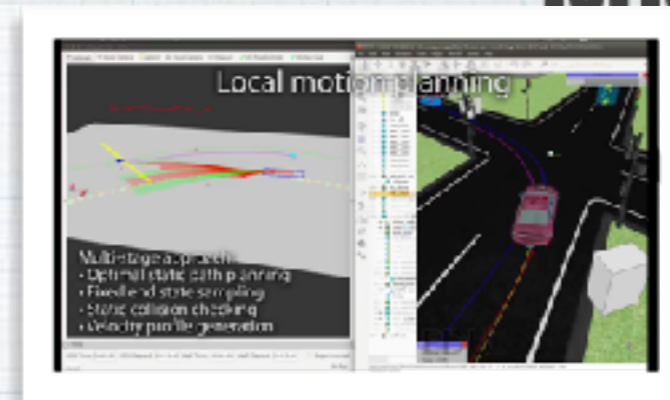
* Uncertainties.

Physical environments, black-box models, statistical AI/ML, ...



On ERATO MMSD

- * JST ERATO Project, 2016/10-2022/03
<https://group-mmm.org/eratommsd>



- * Our goal:
formal methods for **cyber-physical systems (CPS)**
 - * Extend **formal methods**, from software to CPS
 - * Safety, reliability, V&V (Verification & Validation).
“**Check if a system behaves as expected**”
 - * Emphasizing **industry collaboration**, also for scientific inspirations
 - * **Automated driving** as a strategic target domain.
Collaboration with U Waterloo:
www.autonomoose.net



- * Our team:
 - * 28 researchers, > 20 students (as of 2019/09)
 - * **International** and **scientifically diverse**

Our Organization

International and multi-disciplinary. “creative chaos”



Kyoto U IS Site:
Advanced Deductive Verification
Leader:
Kohei Suenaga

Kyoto U RIMS Site:
Categorical Infrastructure
Leader:
Masahito Hasegawa

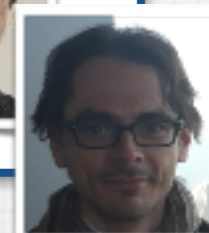
Group 0 @ NII:
Metatheoretical Integration
Leader: Shin-ya Katsumata

Topics:
Programming Languages,
Formal Semantics,
Categorical Models,
Mathematical Logic, ...



Group 3 @ NII:
Formal Methods and Intelligence

Leader: Fuyuki Ishikawa
Subleader: Paolo Arcaini
Topics:
Software Engineering,
Formal Modeling,
Testing, Safe & Explainable AI



Kyushu U Site:
Optimization for CPS V&V
Leader:
Hayato Waki

Osaka U Site:
Control Theory for CPS
Leader:
Toshimitsu Ushio

Group 1 @ NII:
Heterogeneous Formal Methods

Leader: Ichiro Hasuo
Subleader: Masako Kishida
Topics:
Automata Theory,
Control Theory,
Formal Verification,
Proof Assistants,
Automated Deduction,
Runtime Verification



Group 2 @ U Waterloo:
Formal Methods in Industry
Leader: Krzysztof Czarnecki

Topics:
Automated Driving, Software Engineering,
Machine Learning



Members

2019.6, Tokyo NII Site

ERATO 運尾メタ数理システムデザインプロジェクト
ERATO Metamathematics for Systems Design Project

国立情報学研究所 & 科学技術振興機構

National Institute of Informatics & Japan Science and Technology Agency



G0

G1

G3

Group Leaders,
Assoc. Prof's



Shin-ya
Katsumata



Ichiro
Hasuo



Masako
Kishida



Etienne
Andre



Fuyuki
Ishikawa



Paolo
Arcaini

Assistant Prof's
(and similar)



Jeremy
Dubut



Taro
Sekiyama



Akihisa
Yamada

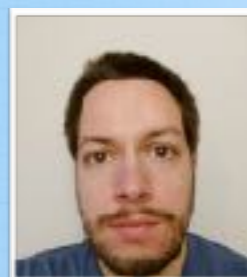


Ahmet
Cetinkaya



Tsutomu
Kobayashi

Post-docs
(and similar)



David
Sprunger



Clovis
Eberhart



Kenta
Cho



Sasinee
Pruekprasert



Toru
Takisaka



Chao
Huang



Xiaoyi
Zhang

Grad Students,
RAs



Satoshi
Kura



Yuichi
Komorida



Takamasa
Okudono



Sang-Hwa
Lee



Liye
Guo



Masaki
Waga



Zhenya
Zhang

- Top researchers from the world, selected out of > 150 applications

Group Leaders,
Assoc. Prof's



Shin-ya Katsumata



Ichiro Hasuo



Etienne Andre



Fuyuki Ishikawa



Paolo Arcaini

FR

IT

G1

TUR
(PhD, TokyoTech)

Assistant Prof's
(and similar)

FR



Jeremy Dubut



Taro Sekiyama



Akil Yam



Ahmet Cetinkaya



Tsutomu Kobayashi

US

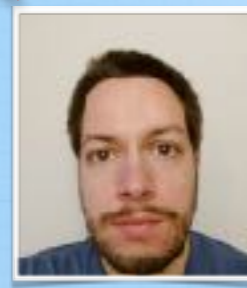
FR

TH
(PhD, Osaka)

CN

CN

Post-docs
(and similar)



David Sprunger



Clovis Eberhart



Kenta Cho



Pr



Ika



Chao Huang



Yoyi ang

KR

CN

CN

Grad Students,
RAs



Satoshi Kura



Yuichi Komorida



Takamasa Okudono



Sang-Hwa Lee



Liye Guo

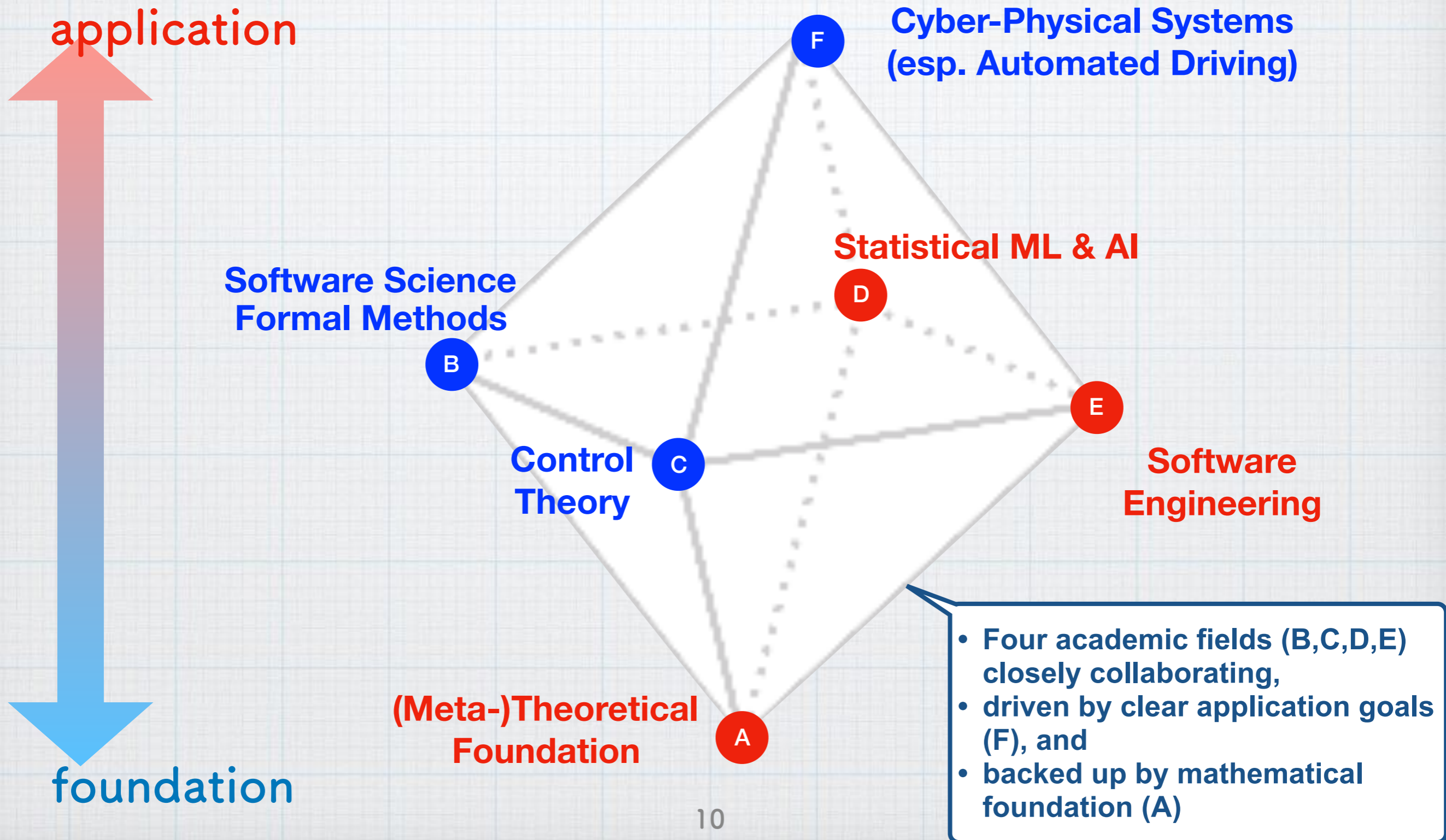


Masaki Waga



Zhenya Zhang

Interdisciplinary Efforts towards CPS: Six Scientific Fields

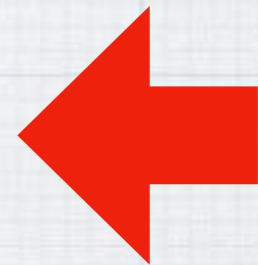


Academic Achievements

- * International visibility:
 - * **CORE rank A*** (top 4%, LICS, CAV, POPL, ...):
> 10 papers
 - * **CORE rank A** (top 5~18%, ATVA, TACAS, GECCO, ICECCS, ...):
> 20 papers
 - * **Best paper awards** :
ICECCS'18, FoSSaCS'19 (CORE rank A), FORMATS'19
 - * **At LICS'19 (CORE rank A*):**
6 out of 60 accepted papers were coauthored by us

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 - * Monitoring
 - * Search-based testing
- * Applications, “**power of math**”



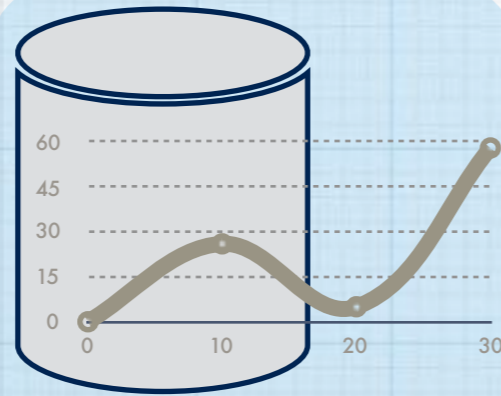
Pattern Matching against Timed Automata, Monitoring

Specification

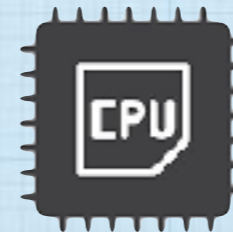
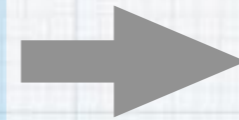
“Frequent gear changes within 3 sec after shifting up to 4th”



Running
system



Log



Monitor



Monitoring
result

“From 18.9 sec. to
23.2 sec.”

- * Runtime verification, monitoring
- * Not straightforward, esp. when specs involve timing constraints
 - * Speed requirements (GBs of log per second)
 - * Computing resource (embedded)
 - * ...
- * Industry needs
- * Technically: theory of **(parametrized) timed automata**



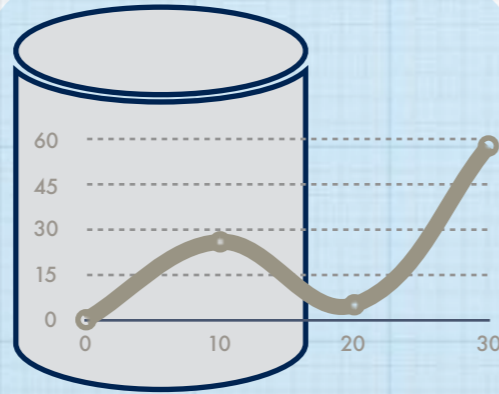
Pattern Matching against Timed Automata, Monitoring

Specification

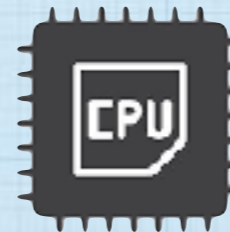
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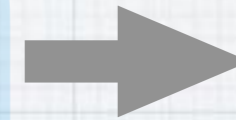
Running
system



Log

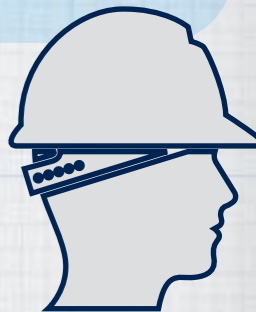


Monitor



Monitoring
result

“From 18.9 sec. to
23.2 sec.”



* Use cases

- * “Here is 1 PB of log, and I want to extract its relevant parts”
- * “Raise an alert if this specific type of anomaly occurs”

Monitoring: Problem Formulations

- * Given: a log, discrete time $w = abaaacb\dots bbc$
 a spec φ "no occurrence of c for 6 steps after b"
 Answer: all subsequences of w that satisfy φ

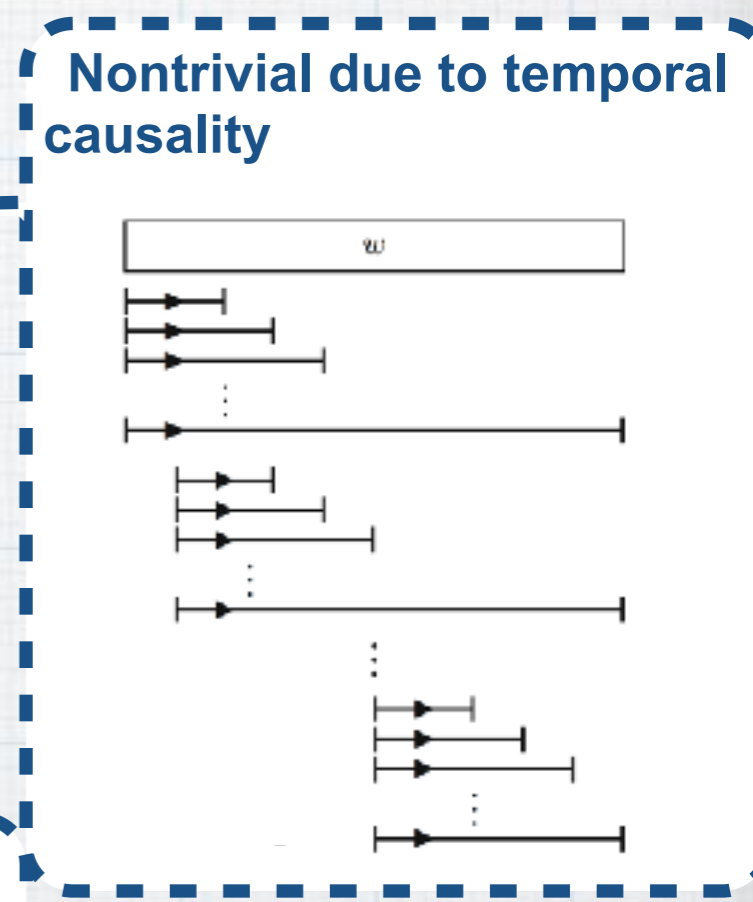
- * Given: a log, **continuous** time $w = (a, 0.12) (b, 1.28) \dots$
 a spec φ "no occurrence of c for 6 **seconds** after b"
 Answer: all subsequences of w that satisfy φ
 [Ulus, CAV'17] [Waga+, FORMATS'17]

Infinitely many such subsequences
 (starting at $t=1$? $t = 1.01$? $t = 1.001$? ...)
 → Efficient representation & computation by **zones**

- * Given: a log, **continuous** time $w = (a, 0.12) (b, 1.28) \dots$
 a **parametrized** spec $\varphi(p)$
 "no occurrence of c for **p** seconds after b" "b occurs with a period of **p** seconds"
 Answer: all the pairs of $(p, (\text{a subseq. of } w \text{ that satisfies } \varphi))$

[Andre+, ICECCS'18] [Waga+, NFM'19] [Waga+, CAV'19]

In industry, fixing a spec is a big challenge.
 Parameters → **flexibility** in specs



Monitoring: Our Achievements

- * Given: a log, discrete time $w = \text{abaaacb...bbc}$
a spec φ "no occurrence of c for 6 steps after b"
Answer: all subsequences of w that satisfy φ

- * Given: a log, **continuous** time $w = (a, 0.12) (b, 1.28) \dots$
a spec φ "no occurrence of c for 6 **seconds** after b"
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Answer: all the pairs of $(p, (\text{a subseq. of } w \text{ that satisfies } \varphi))$
[Andre+, ICECCS'18] [Waga+, NFM'19] [Waga+, CAV'19] ...

Efficient algorithm from theory of **timed automata**.
Processes ~ 1M events/second (laptop).

[Waga+, FORMATS'17]
<https://github.com/maswag/monaa>

Also implemented on
Renesas RH850

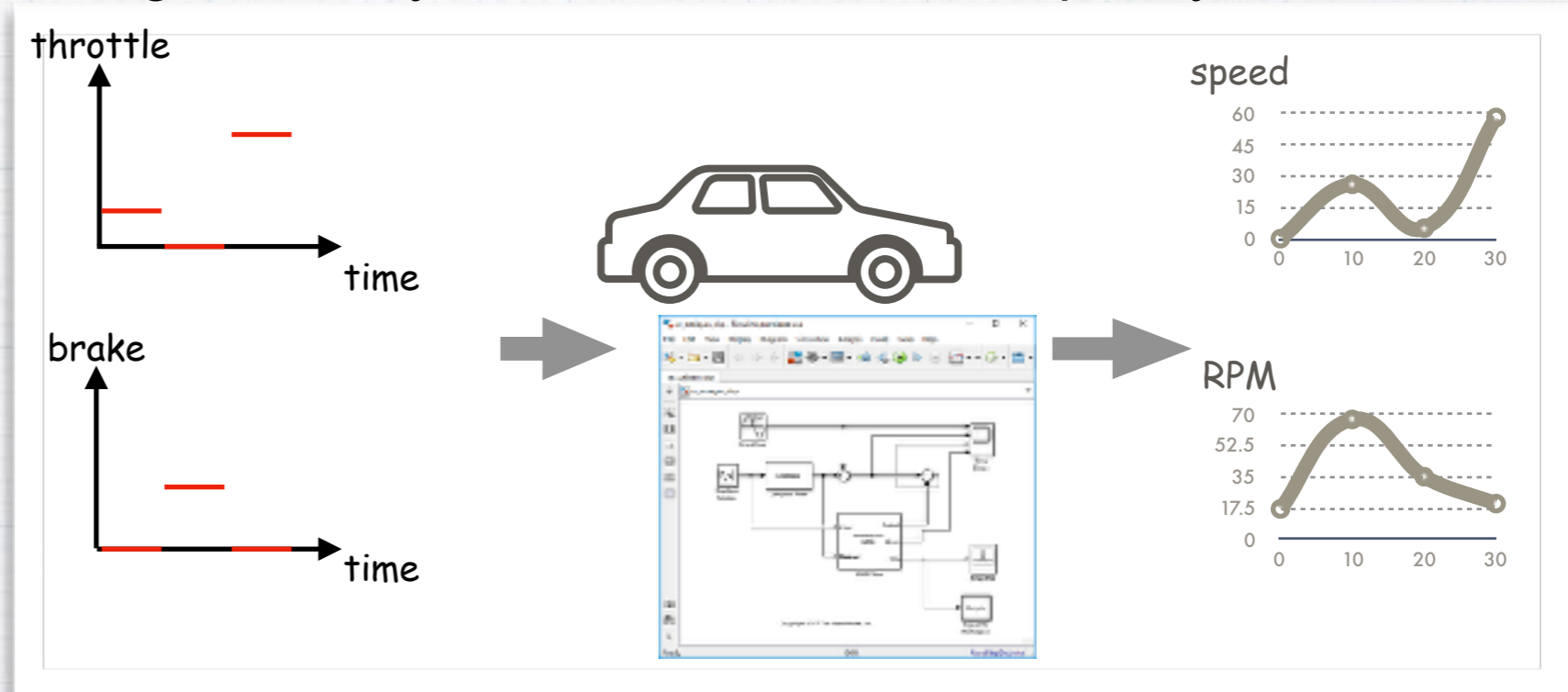
Efficient algorithm using **parametrized timed automata**.
Processes ~ 10K events/second (laptop). [Waga+, NFM'19]
<https://github.com/maswag/symon>

Testing by Reinforcement Learning

[Zhang+, EMSOFT'18] [Zhang+, CAV'19] [Fainekos & Pappas, TCS'09]

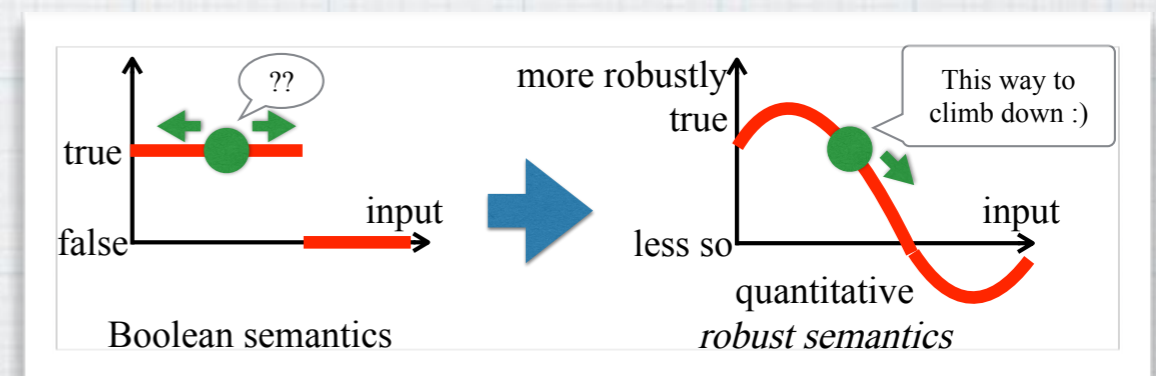
* **Black-box**, search-based testing. Actively search for erroneous input by:

- * Try an input signal
- * Observe the system's behavior
- * Choose the next input that is likely to be erroneous



* [Fainekos & Pappas, TCS'09]
A **reinforcement learning problem**,
by moving

- * from the **Boolean** semantics (erroneous or not)
- * to **quantitative** “robust semantics” (how far from being erroneous)



Search-Based Testing by Reinforcement Learning

[Akazaki & Hasuo, CAV'15]

[Zhang, Ernst, Sedwards, Arcaini & Hasuo, EMSOFT'18]

[Zhang, Hasuo & Arcaini, CAV'19] ...

- * Survey: [Kapinski+, IEEE Control Syst. '16]

J. Kapinski, J. V. Deshmukh, X. Jin, H. Ito, and K. Butts, "Simulation-based approaches for verification of embedded control systems: An overview of traditional and advanced modeling, testing, and verification techniques," IEEE Control Syst., vol. 36, no. 6, pp. 45–64, Dec. 2016.

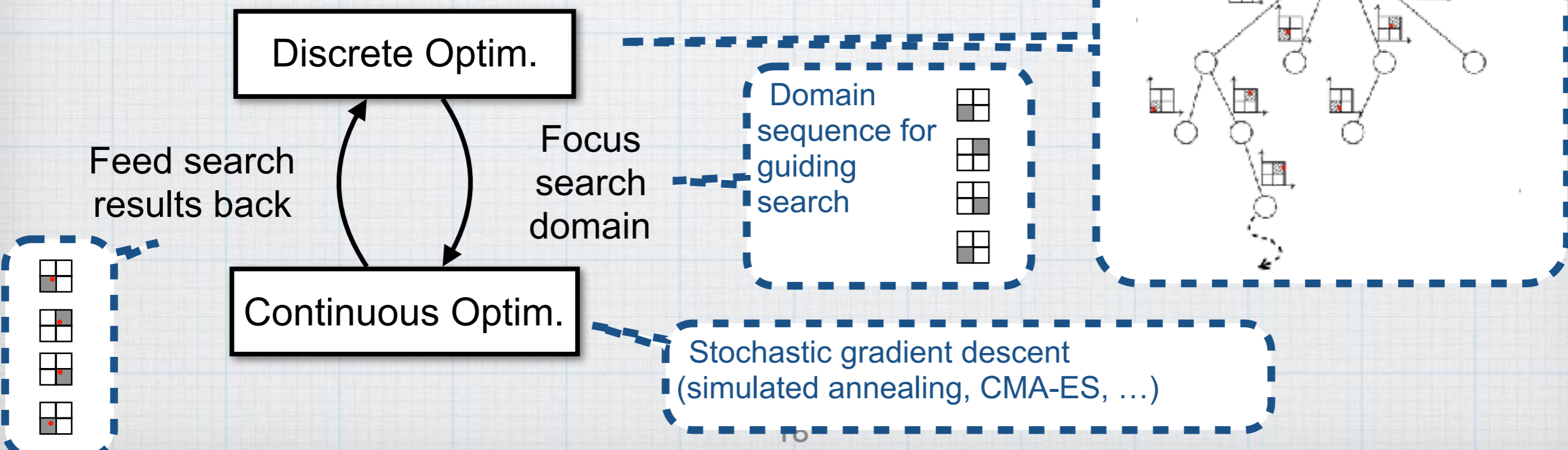
- * **Our Contribution** In [Zhang+, EMSOFT'18] [Zhang+, CAV'19]

- * Exploit **discrete structures** in **stochastic gradient descent**

- * → efficient and extensive search

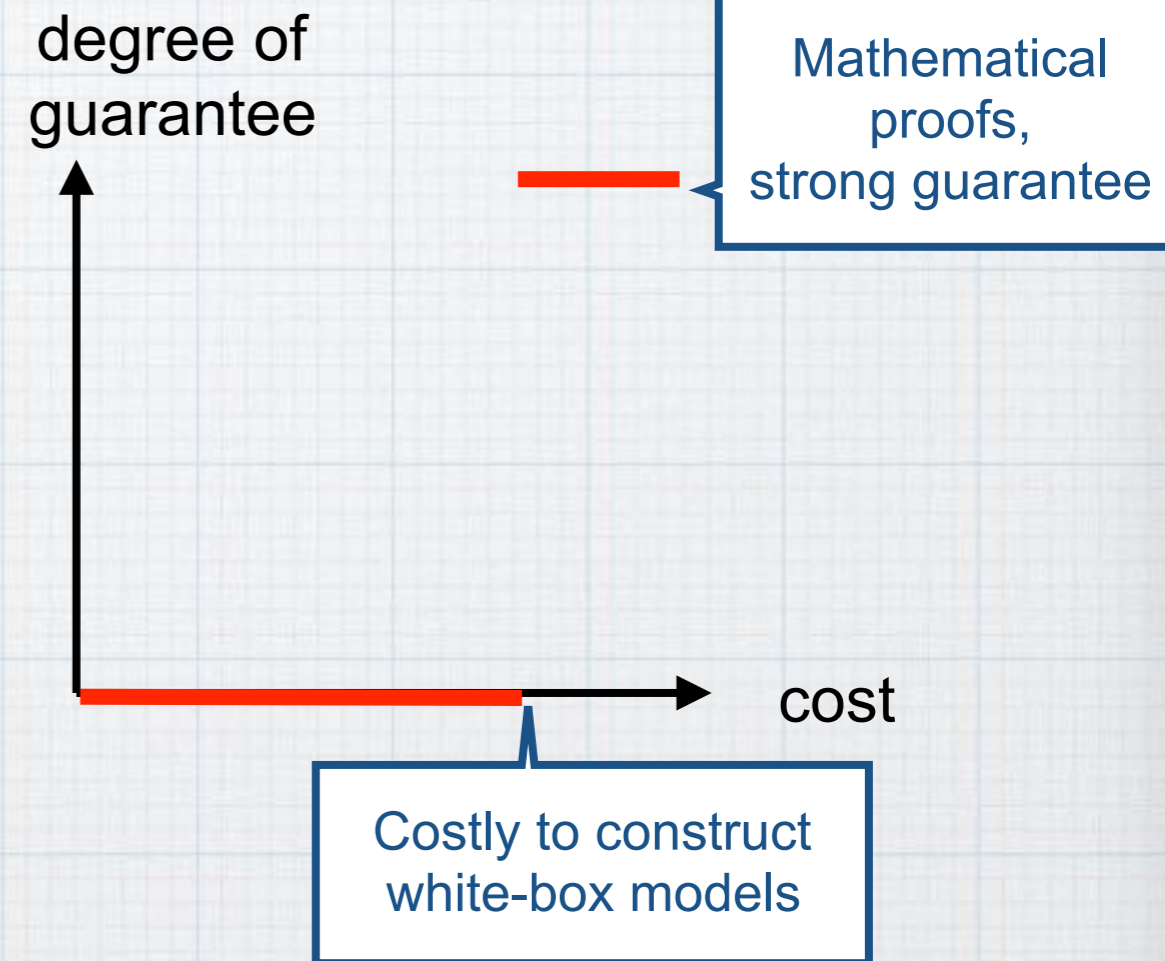
- * → interpretable testing

- * In the form of **hierarchical optimization**



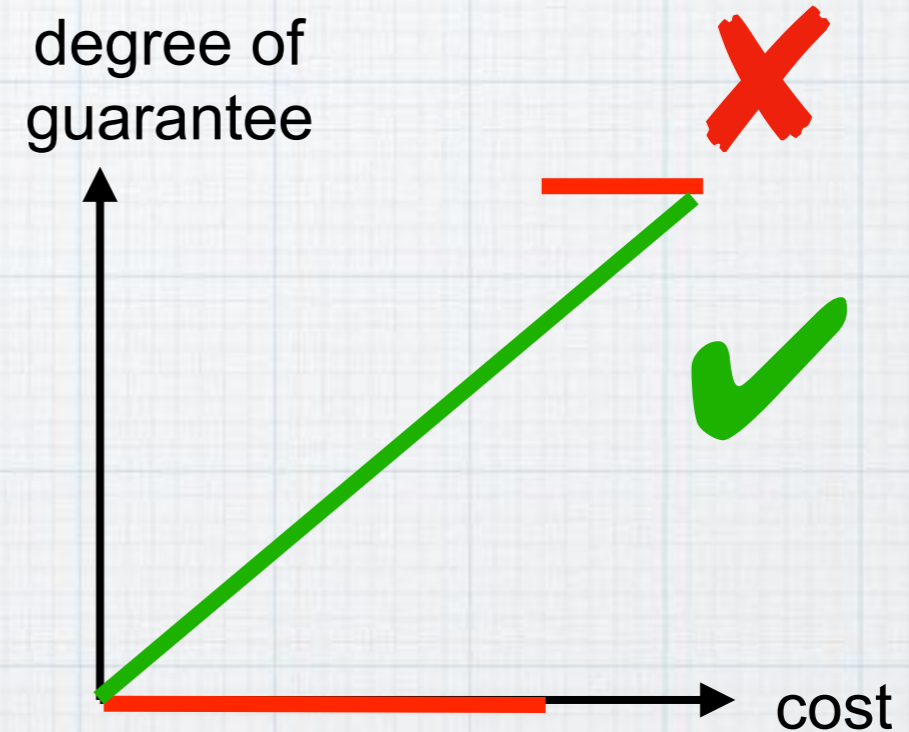
Towards Formal Methods That Are Down-Scalable

- * Many FM methods are not **down-scalable**
 - * They demand **white-box models**
 - * → huge cost before non-zero benefit
- * Example applications to CPS:
airplanes (Airbus), space (NASA), ...
 - * Failure is so expensive
→ worth verification efforts
 - * Still takes decades
 - * Integrating continuous dynamics is hard
→ They focus on software



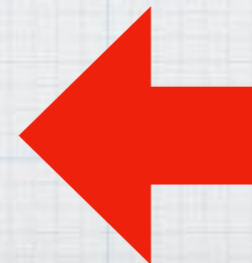
Towards Formal Methods That Are Down-Scalable

- * Real-world applications (esp. automotive domain), need **down-scalability**
 - * Even if we can only afford **half the cost**,
 - * degree of guarantee does not become **zero**, but **half**
- * Our efforts at ERATO MMSD
 - * Combine **testing** and **formal verification**.
Search-based testing, monitoring
 - * **Contract-based verification**
→ **formal safety architecture**
(contract-based verification, see reserve slides)



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Industry Collaboration

- * Mainly **automotive**, but not exclusively
 - * Production systems
 - * AI startups
- * Inspiring **scientific research**, too
 - * New problems
 - * Solve a problem
 - generalize the solution
 - **scientific novelty**
- * A lot of **industry needs**
 - * Methods applicable off-the-shelf (monitoring, search-based testing)
 - * High demand for formal verification
 - * “What is an expected level of quality assurance?”

Collaboration with U Waterloo

test tool

env. input

(road shape,
pedestrians,
other cars, ...)

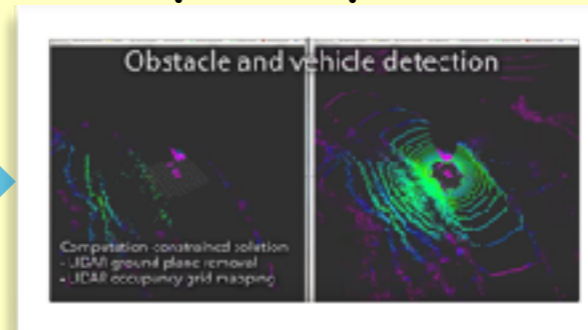
ego car's
position & status

Our goal: comprehensive software stack
with perception units, controllers,
simulators, and testing tools

(cf. ROS, robot operating system)

- serving academic users
- used as (part of) prototype products
- testbed for technical components

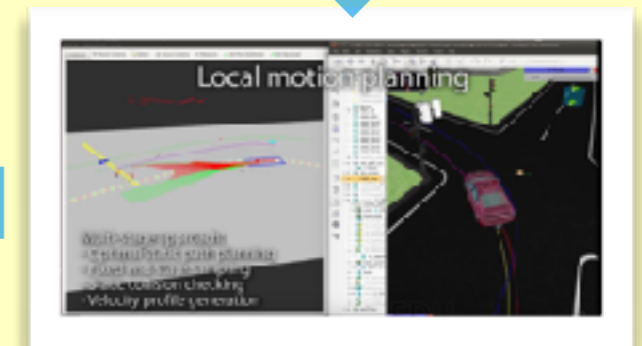
perception



object recognition



local planning
path tracking



behavior planning
path planning

Simulation environment

ERATO 蓮尾メタ数理システムデザインプロジェクト
ERATO Metamathematics for Systems Design Project

国立情報学研究所 & 科学技術振興機構

National Institute of Informatics & Japan Science and Technology Agency

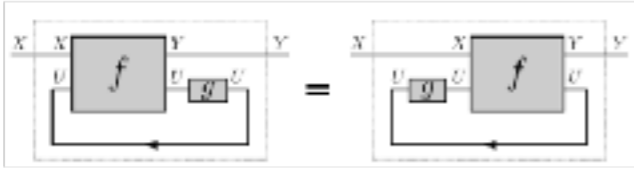


logic, algebra,
 category theory,
 ...



Traced monoidal categories

$$\text{Tr}_{X,Y}^U : \mathbf{C}(X \otimes U, Y \otimes U) \rightarrow \mathbf{C}(X, Y)$$



[Sprunger & Katsumata, LICS'19]

Abstract Technique

$$T[_]$$

$$A ::= \text{true} \mid \text{false} \mid A_1 \wedge A_2 \mid \neg A \mid a_1 < a_2 \mid \forall x \in {}^*\mathbb{N}. A \mid \forall x \in {}^*\mathbb{R}. A$$

$$\begin{array}{ccc} FX \xrightarrow{F\text{beh}_c} FZ & & FX \xrightarrow{Ff} FY \\ c \uparrow & \uparrow \text{final} & c \uparrow \quad \exists \quad \uparrow d \\ X \xrightarrow{\text{beh}_c} Z & & X \xrightarrow{f} Y \end{array}$$

system behavior simulation

Identify
 "mathematical
 essence"

Choose
 parameter e_1

Choose
 parameter e_2

Existing Technique

$$T_1 = T[e_1]$$

Novel Technique

$$T[e_2]$$

knot invariants, QFTs,
 semantics of recursion, ...

back-propagation algorithm
 for recurrent neural networks

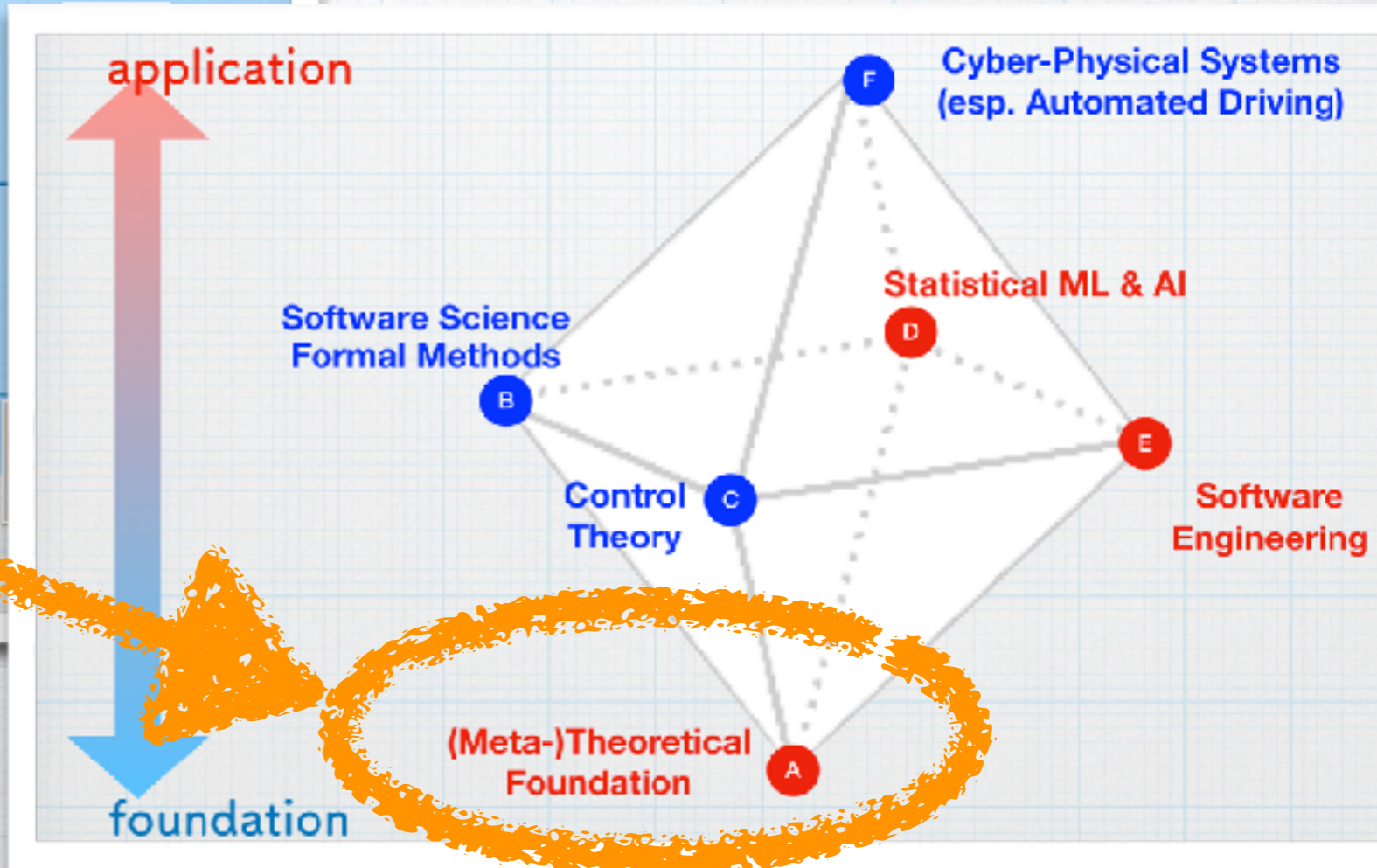
Reachability in
 Automata

Result that stemmed from
 "deep learning reading group" at ERATO MMSD



Power of Math

Group Leaders, Assoc. Prof's	 Shin-ya Katsumata	 Ichiro Hasuo	 Masako Kishida	 Etienne Andre	 Fuyuki Ishikawa	 Paolo Arcaini
Assistant Prof (and similar)	 Jeremy Dubut	 Taro Sakiyama	 Akhisa Yamada	 Ahmet Cetinkaya		
Post-docs (and similar)	 David Sprunger	 Clovis Eberhart	 Kenta Cho	 Sasine Pruekprasert	 Teru Takasaka	 Chao Huang
Grad Students RAs	 Satoshi Kura	 Yulchi Komorida	 Takamasa Okudono	 Sang-Hwa Lee	 Uyi Guo	



- * Contribution of “**theoreticians and listeners**” have been remarkable, direct and indirect
 - * “Metatheoretical transfer”
 - * Interpreters between disciplines
 - * Great “reserve forces”

Conclusions

- * Quality assurance of **cyber-physical systems**
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Thanks for your attention!
<https://group-mmm.org/eratommsd>

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Reserve Slides

Formal Safety Architecture

- * Example: simplex architecture (right)

- * **AC** is complex, performance-oriented, black-box
- * **BC** is simple, safety-oriented, (hopefully) white-box

- * Idea: we can verify the whole system **even if AC is a black-box!**

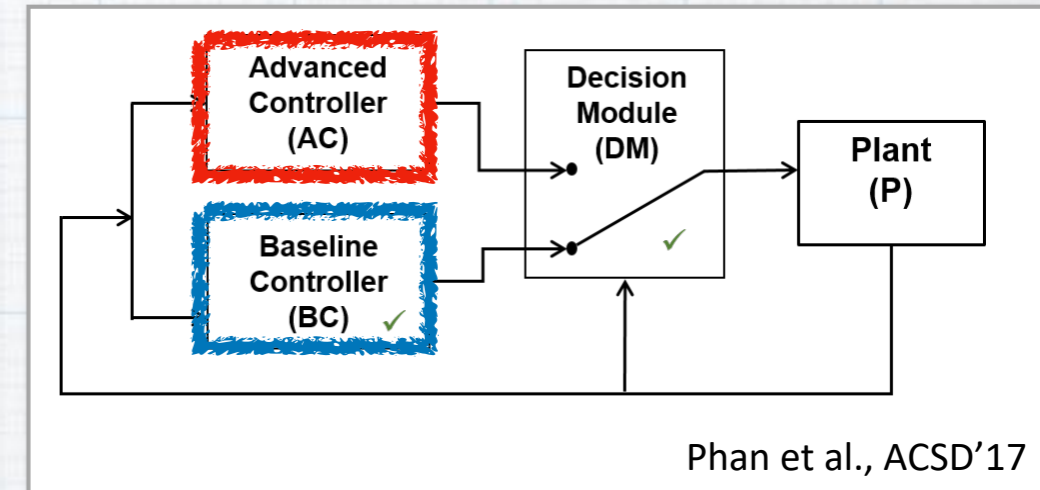
- * Enough to show: safety of BC, and correctness of DM
- * We impose certain **contracts** on AC

- * Related

- * FM4AI (ML/AI components as AC)
- * Monitoring (checking contracts on AC)

- * One promising way to **make formal verification down-scalable**

- * Weaker contracts on AC
→ weaker safety guarantee (but hopefully non-zero)



Formal Safety Architecture

- * At ERATO MMSD:
we **formalize, verify and refine** safety architectures
- * **Event-B: a formal modeling language**
[Abrial, “The Event-B Book”, 2010 CUP] [Kobayashi+, ICFEM’18]
Based on state transition systems.
A tool Rodin supports:
 - * **Safety proofs**
 - * **Incremental modeling by refinements**
 - * Flexibility in choosing model fidelity → **down-scaling**
 - * From (our) general model to (industry partner’s) individual model

