

# Coalgebraic Representation Theory of Fractals

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Radboud Univ. Nijmegen, NL

**Milad Niqui**

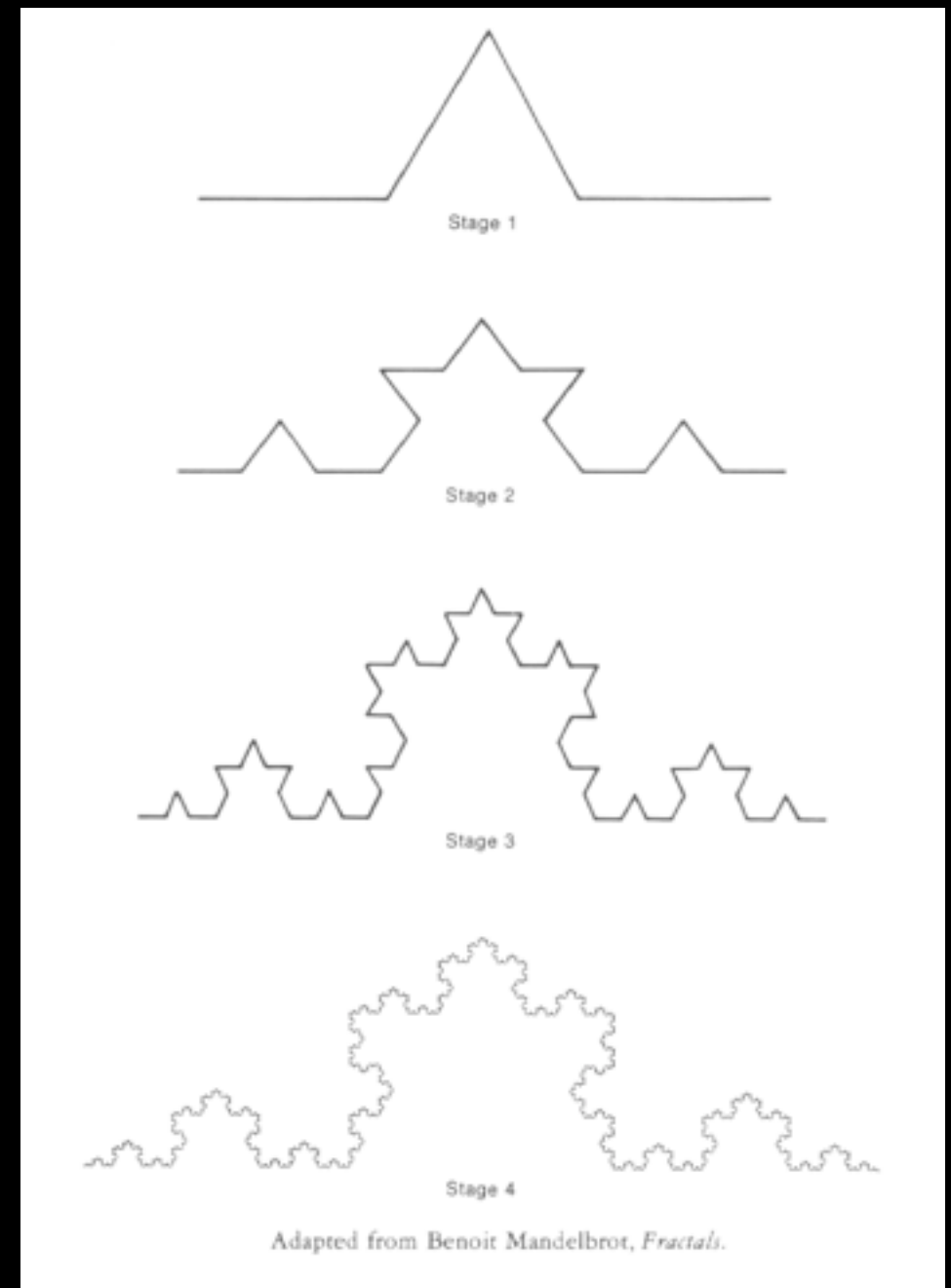
CWI, NL

30 min talk, 5 min Q&A

# Fractals

*the Koch curve*

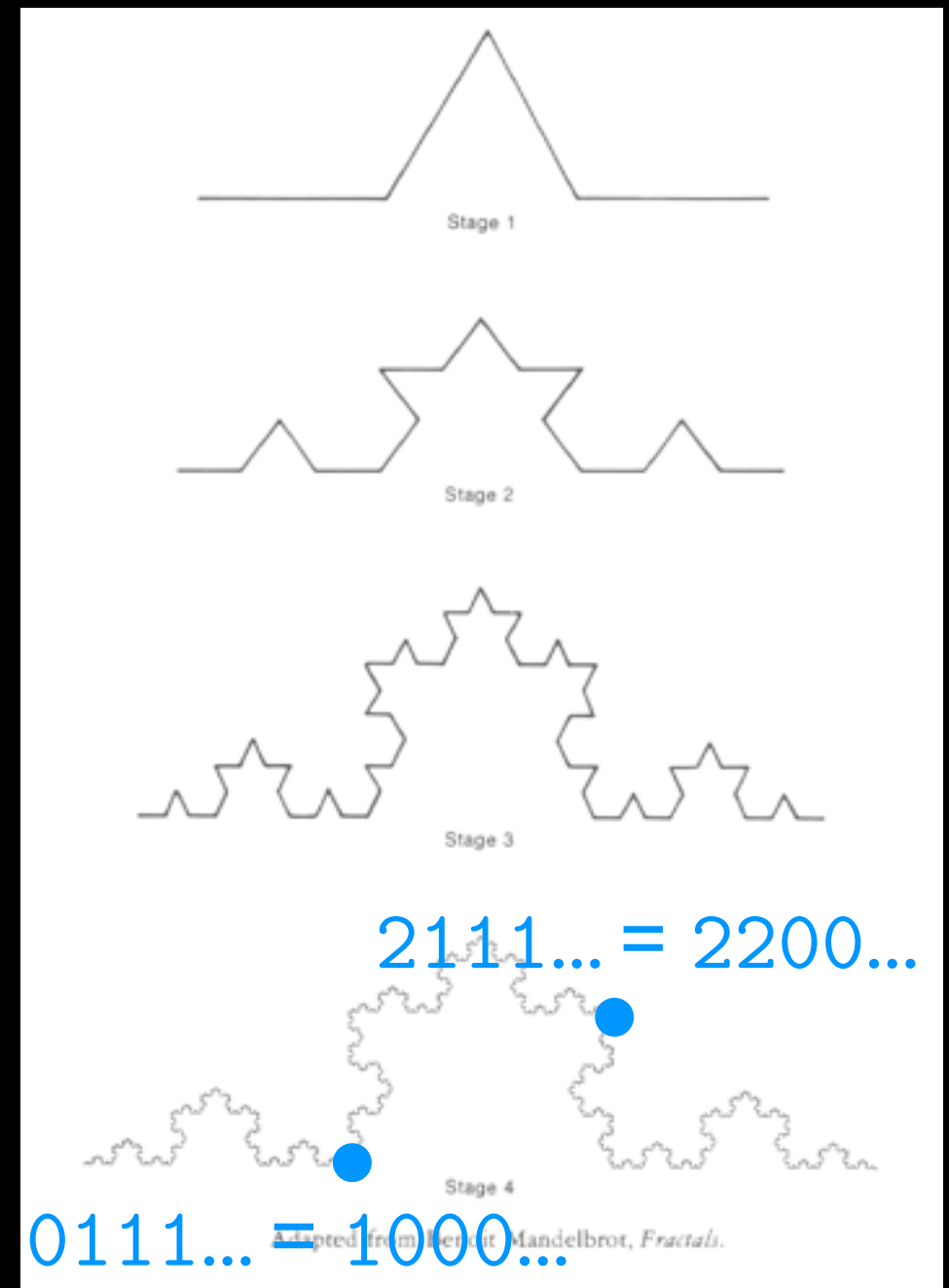
“A rough or fragmented geometric shape that can be split into parts, each of which (at least approximately) a reduced size copy of the whole” [Mandelbrot]



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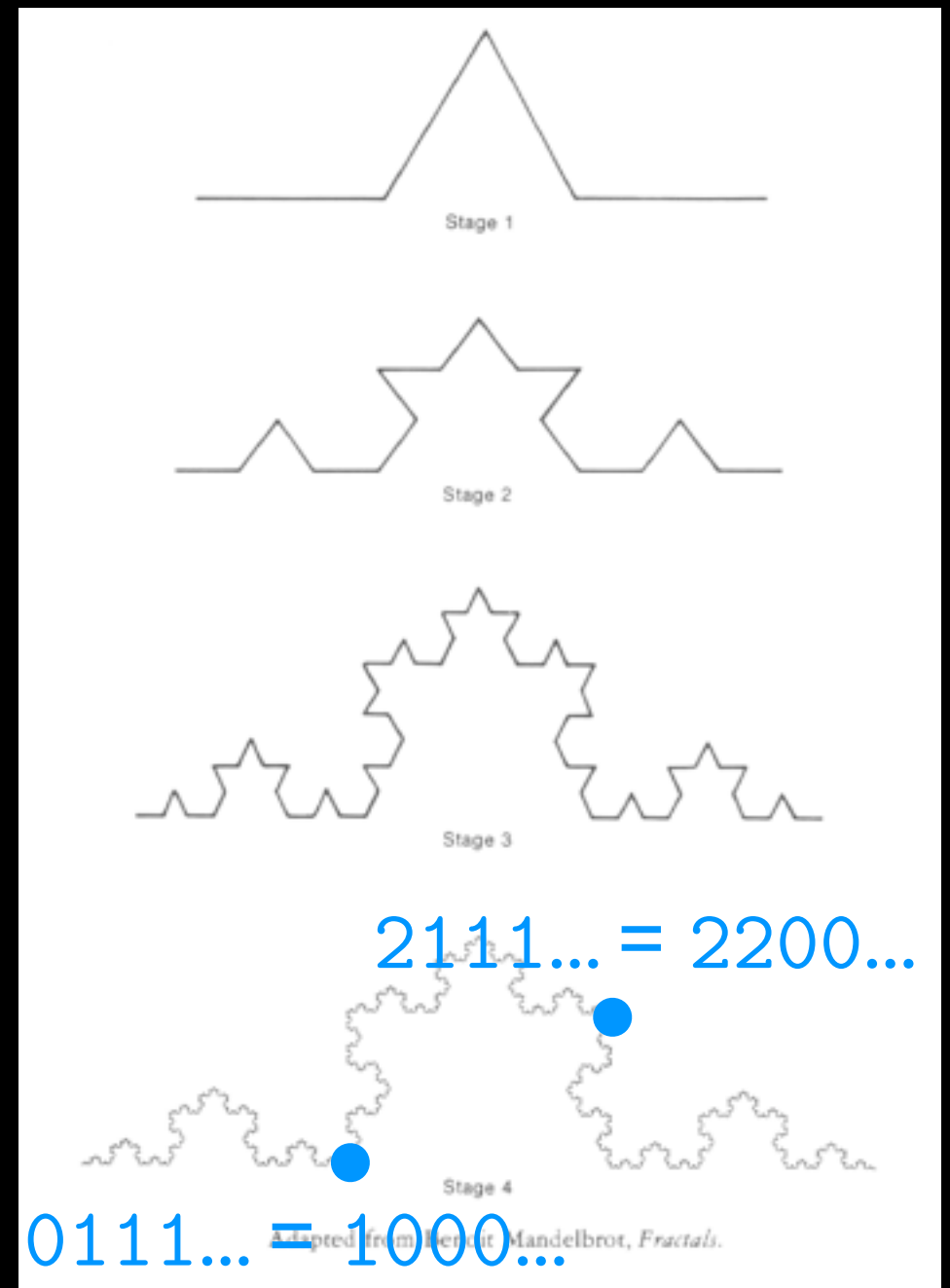
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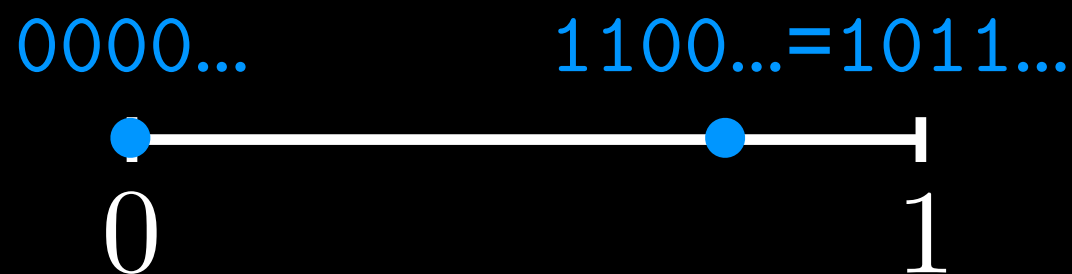
*the unit interval*  $\mathbb{I}$

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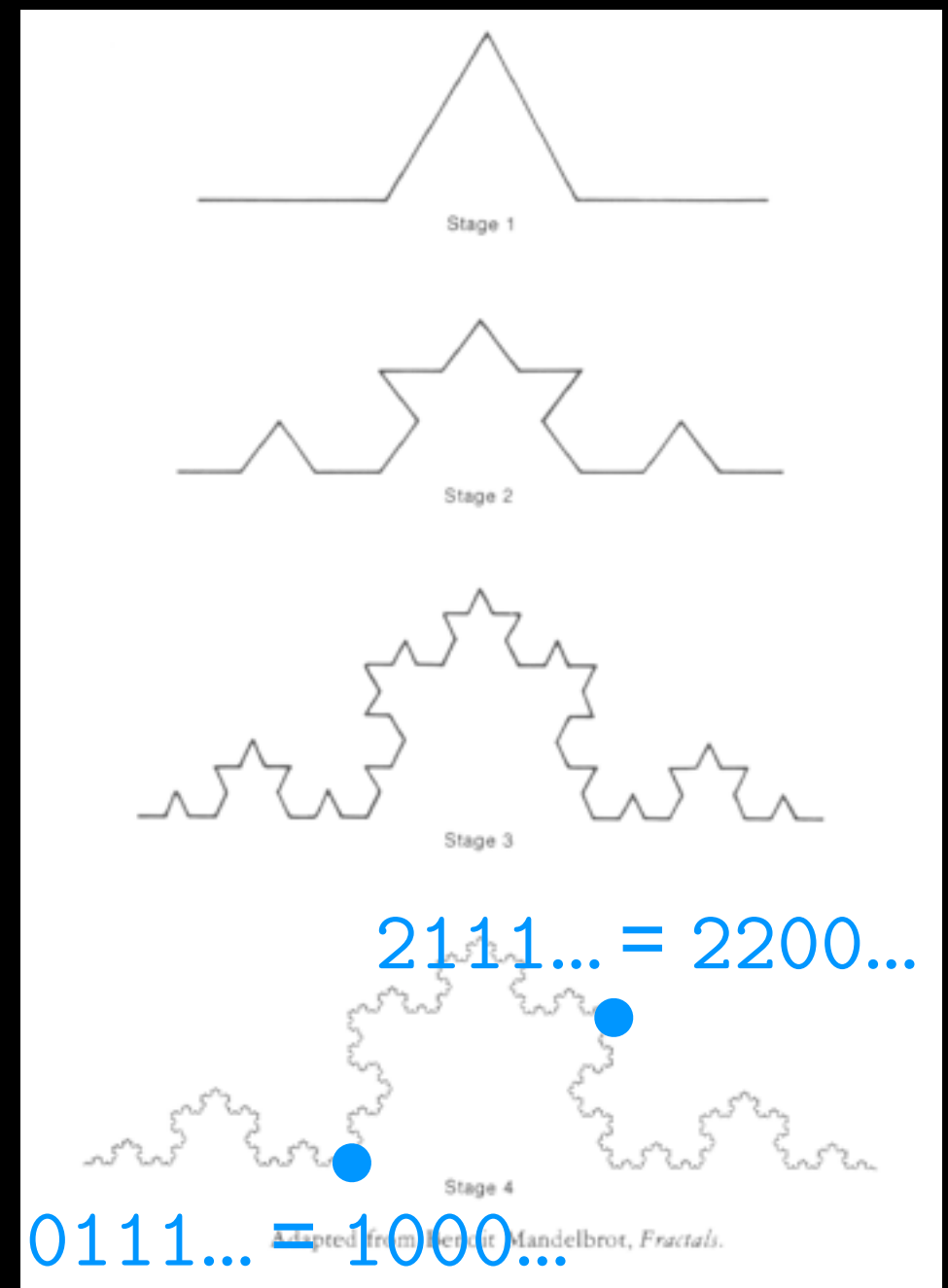
# Fractals

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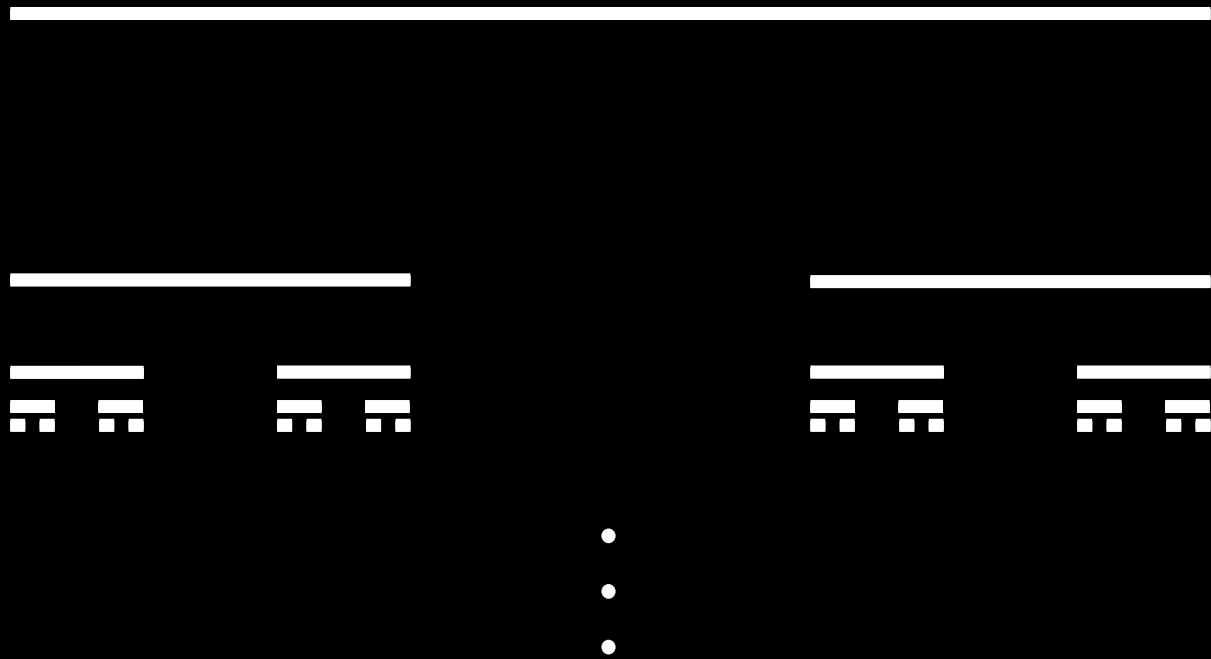


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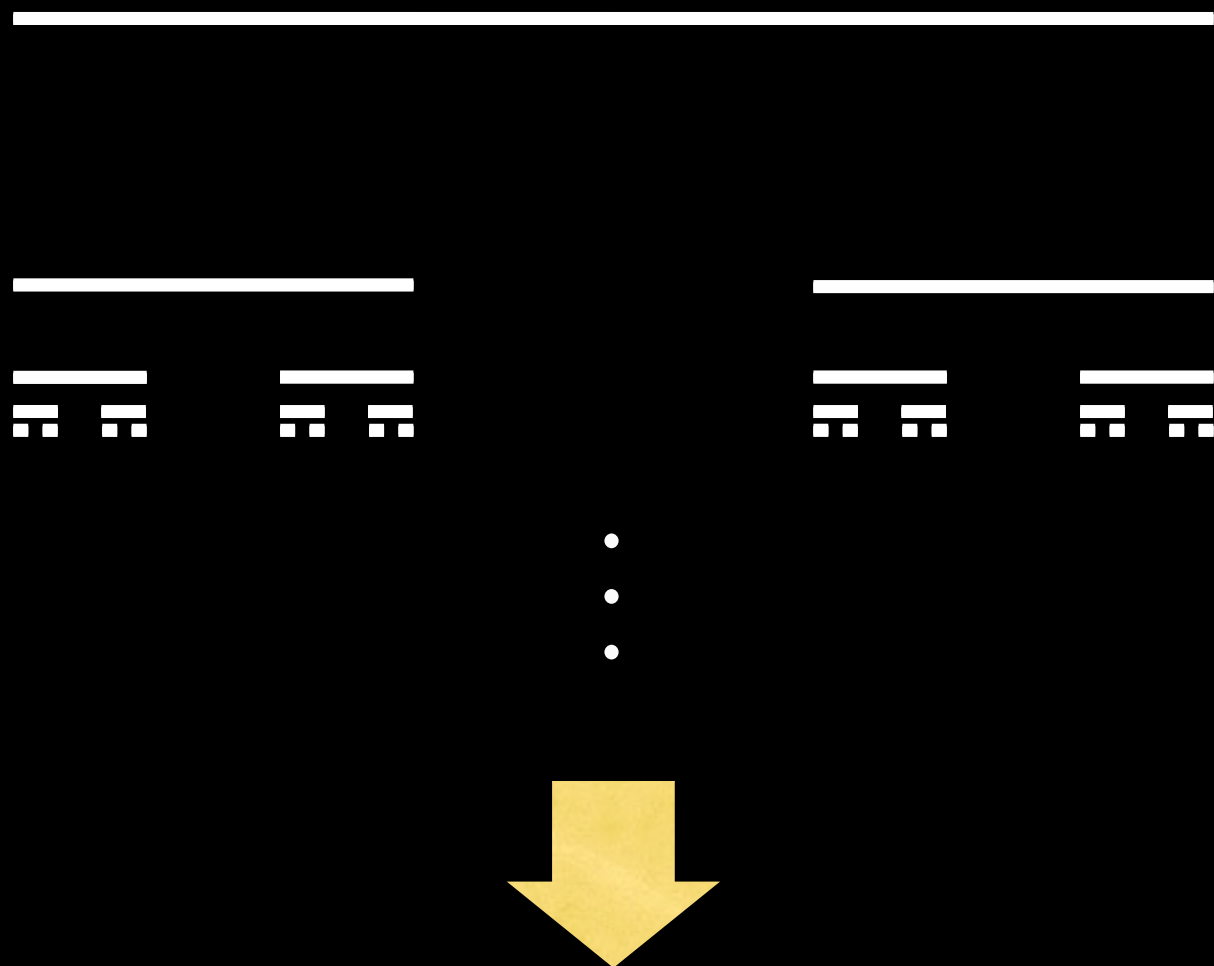
# The Cantor Set



the Cantor set  $C$

# The Cantor Set

*Iterated function system  
(IFS)*



the Cantor set  $C$

$$\begin{aligned} \varphi_0, \varphi_1 &: \mathbb{I} \longrightarrow \mathbb{I} , \\ \varphi_0(x) &= \frac{x}{3} , \\ \varphi_1(x) &= \frac{2+x}{3} . \end{aligned}$$

$C$  as the unique *attractor*:

$$C = \varphi_0(C) \cup \varphi_1(C)$$

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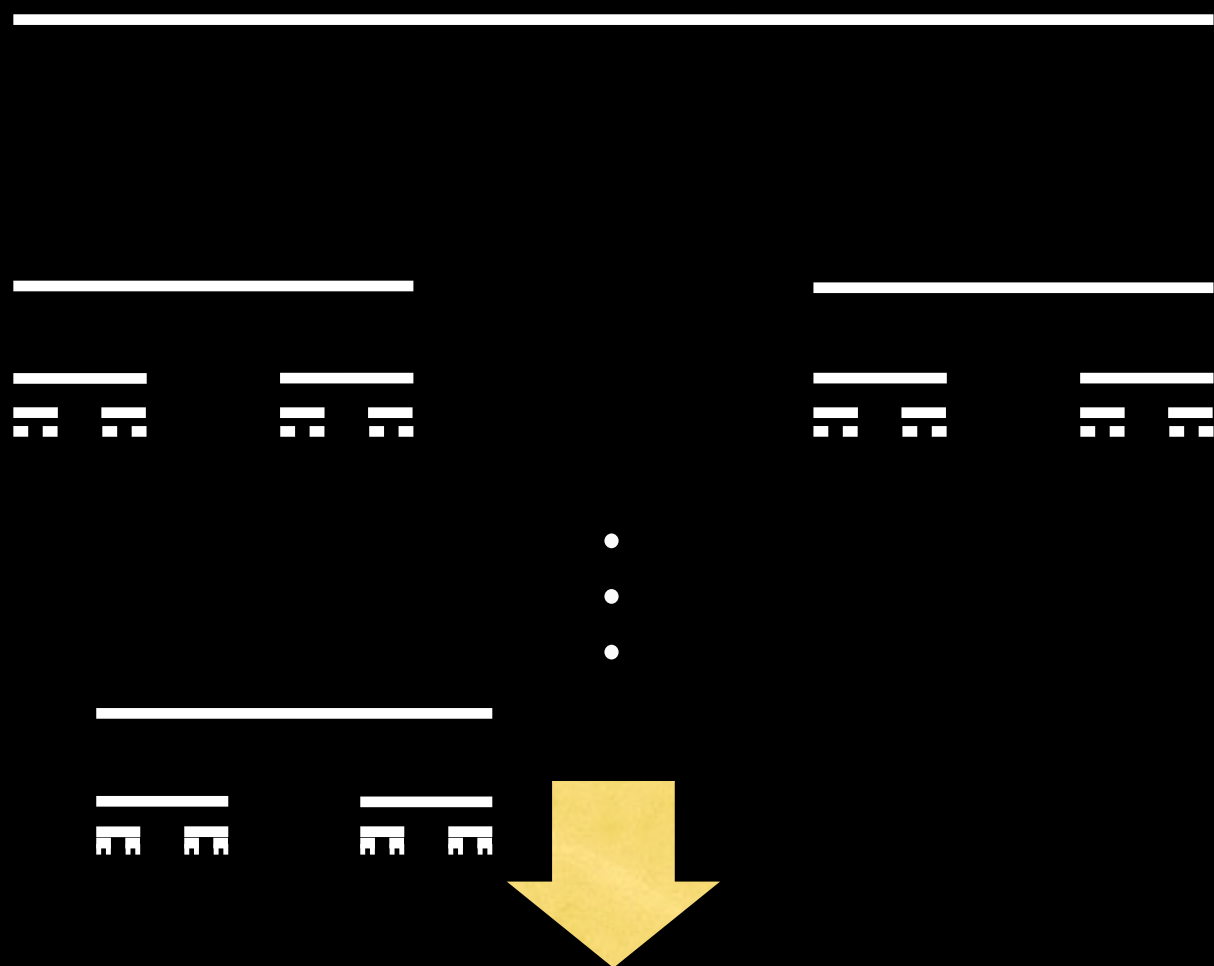
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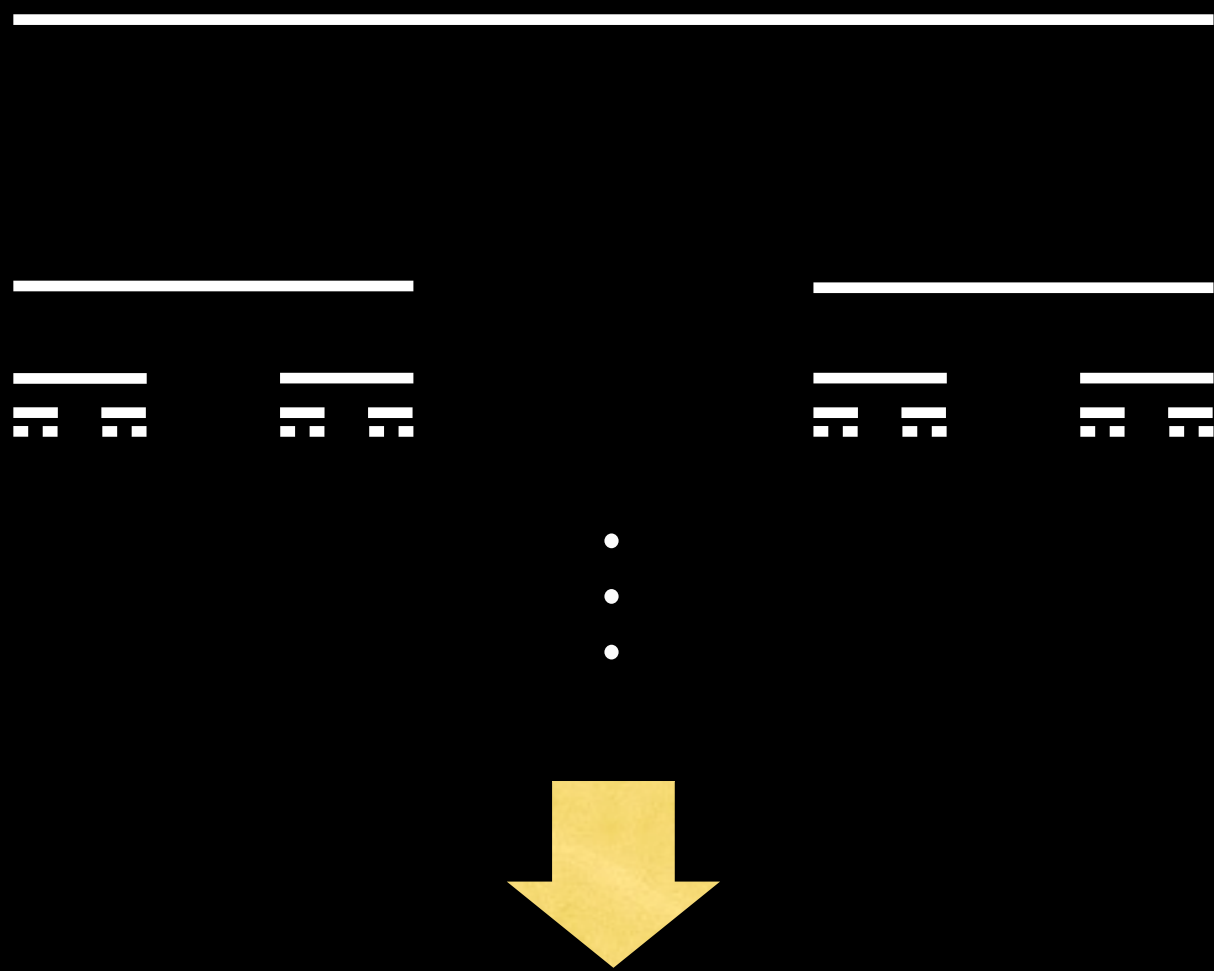
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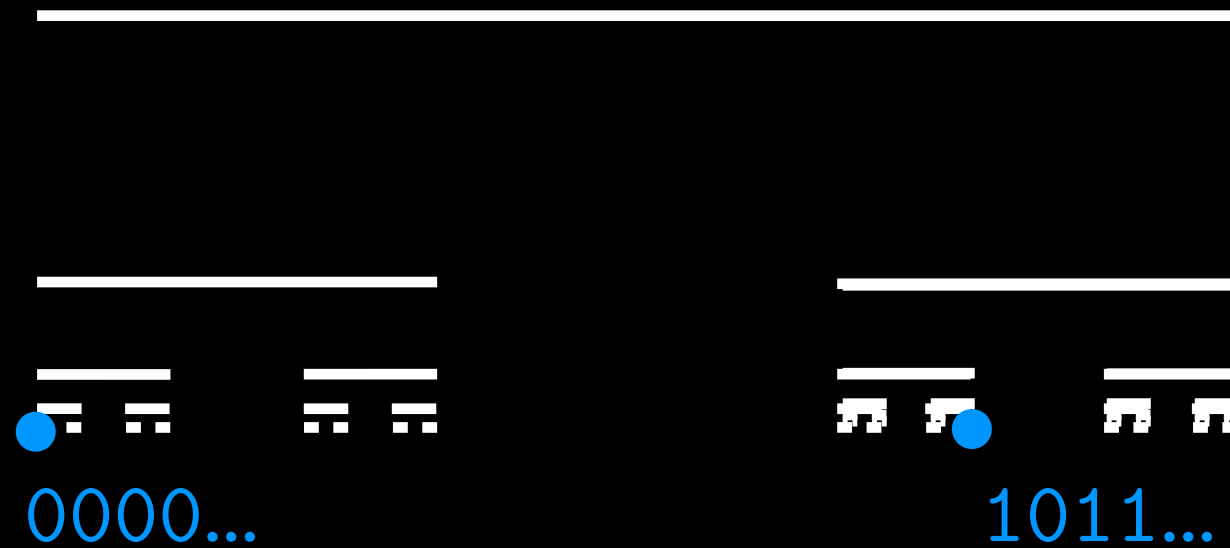
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# Symbolic Representation



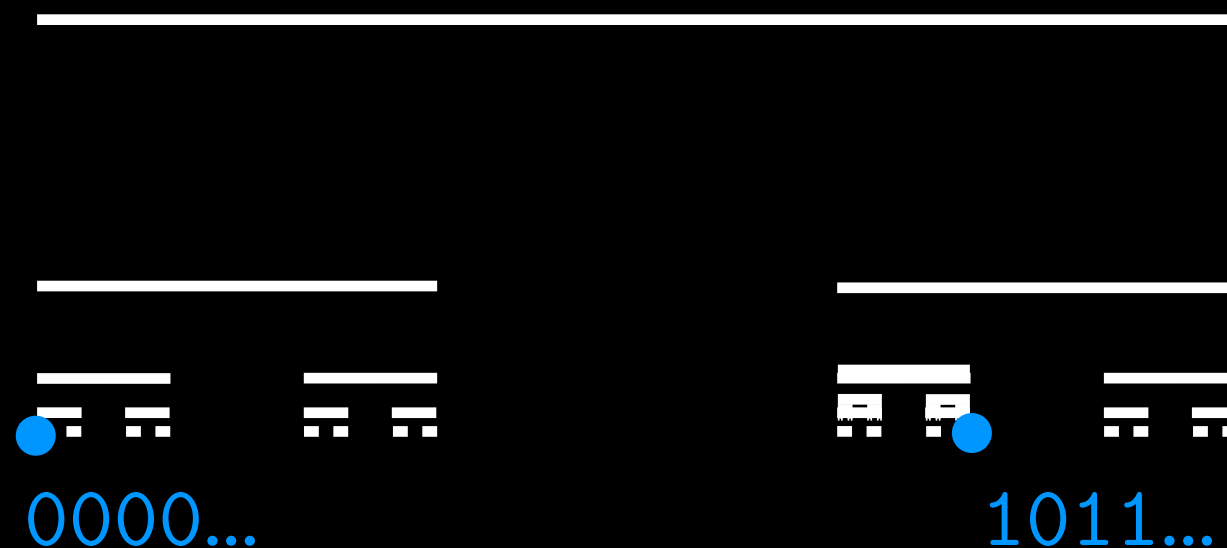
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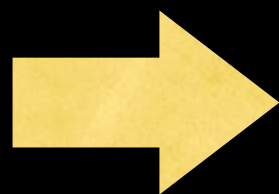
# Symbolic Representation



# Symbolic Representation



fractal  
structure



symbolic representation of the shape,  
by infinite streams

# Observation 1

Symb. representatives

$$2^\omega = \{0, 1\}^\omega$$

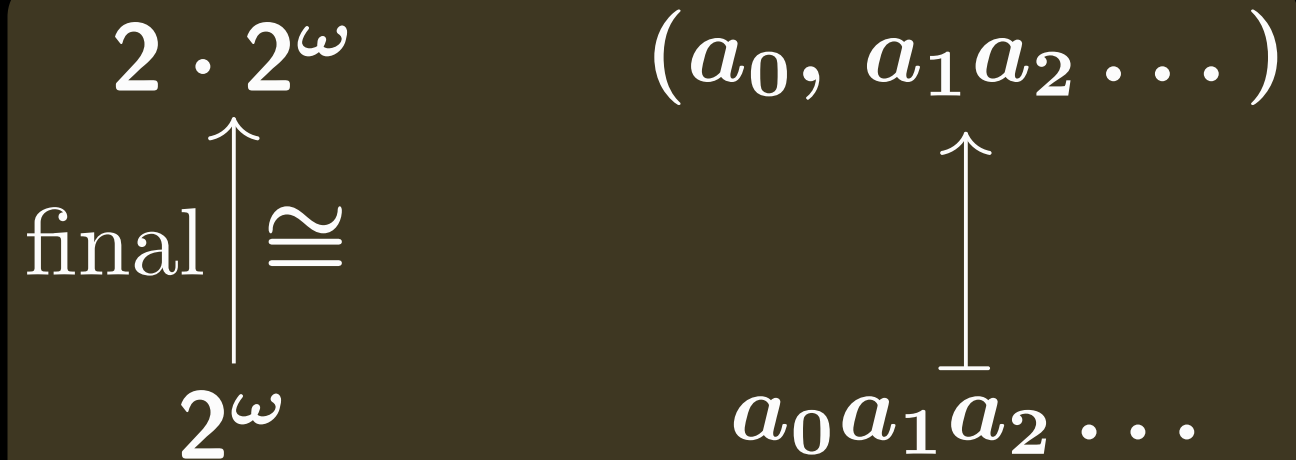
carry a **final coalgebra**

# Observation I

Symb. representatives

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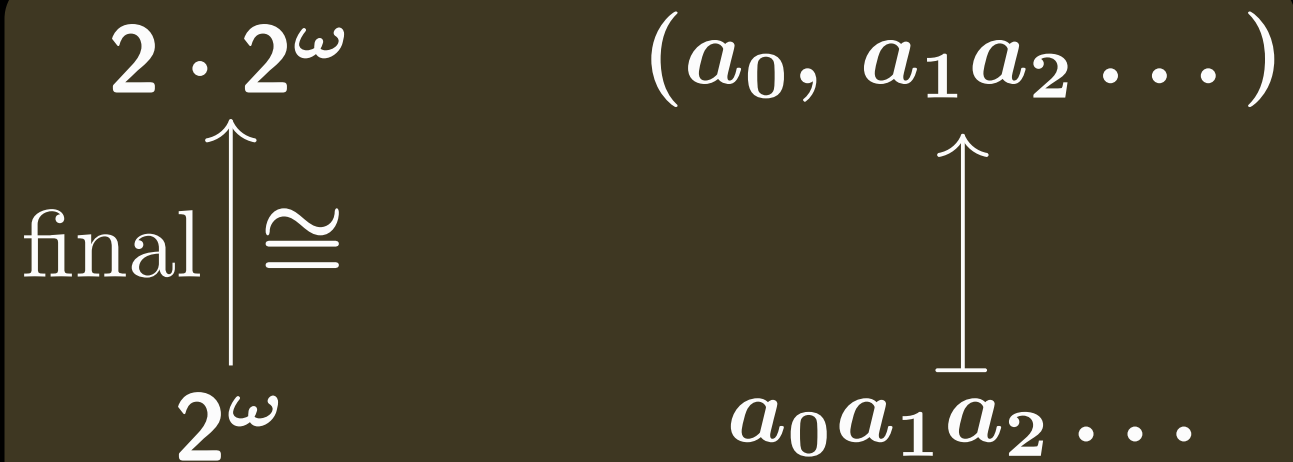
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carry a **final coalgebra**

- final coalgebra as a “fractal”



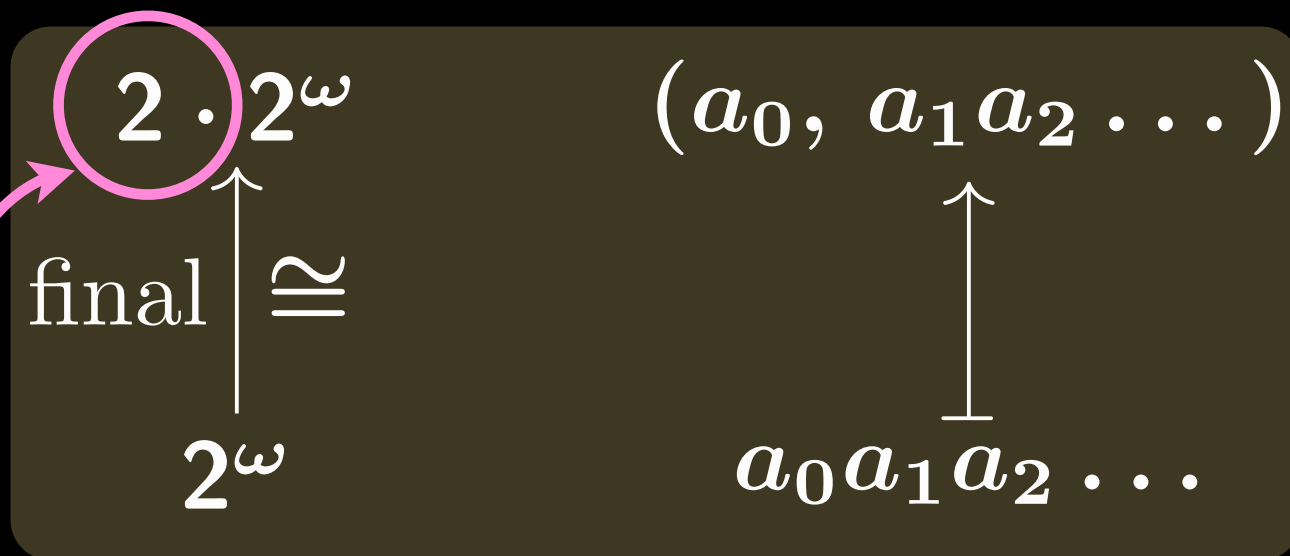
# Observation I

Symb. representatives

$$2^\omega = \{0, 1\}^\omega$$

carry a **final coalgebra**

- final coalgebra as a “fractal”



- *combinatorial specification of the Cantor set:*

$$2 \cdot (\_) : \mathbf{Sets} \rightarrow \mathbf{Sets}$$

# Observation 2

$$\varphi_0, \varphi_1 : \mathbb{I} \longrightarrow \mathbb{I} ,$$

$$\varphi_0(x) = \frac{x}{3} ,$$

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$2 \cdot \mathbb{I}$

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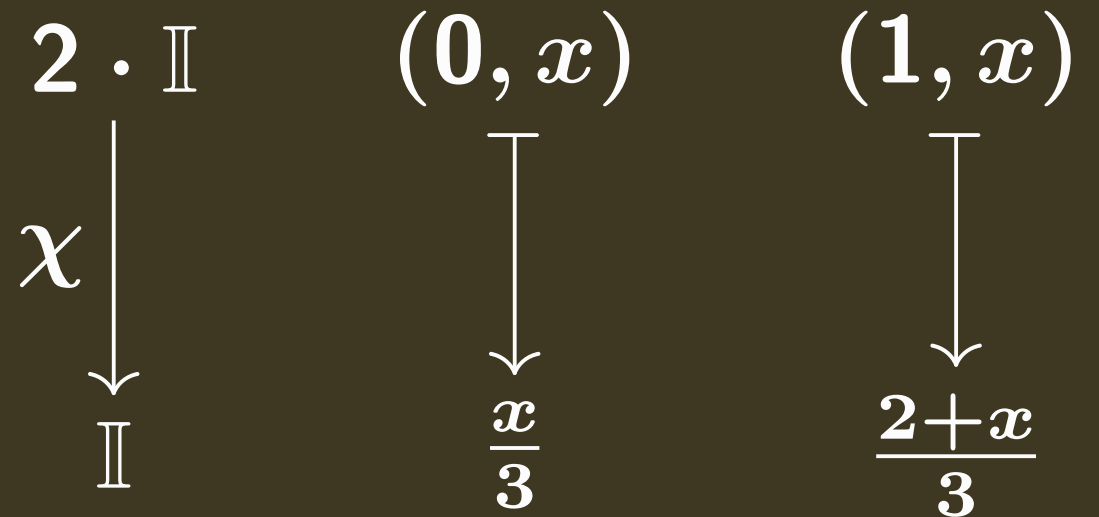
An IFS is an **algebra**

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An IFS is an **algebra**

- for the same functor  $2 \cdot (\_ ) : \mathbf{Sets} \longrightarrow \mathbf{Sets}$

# Observation 3

$$\begin{array}{c} 2 \cdot 2^\omega \\ \uparrow \cong \text{final} \\ 2^\omega \end{array}$$

$$\begin{array}{c} 2 \cdot \mathbb{I} \\ \downarrow \chi \\ \mathbb{I} \end{array}$$

# Observation 3

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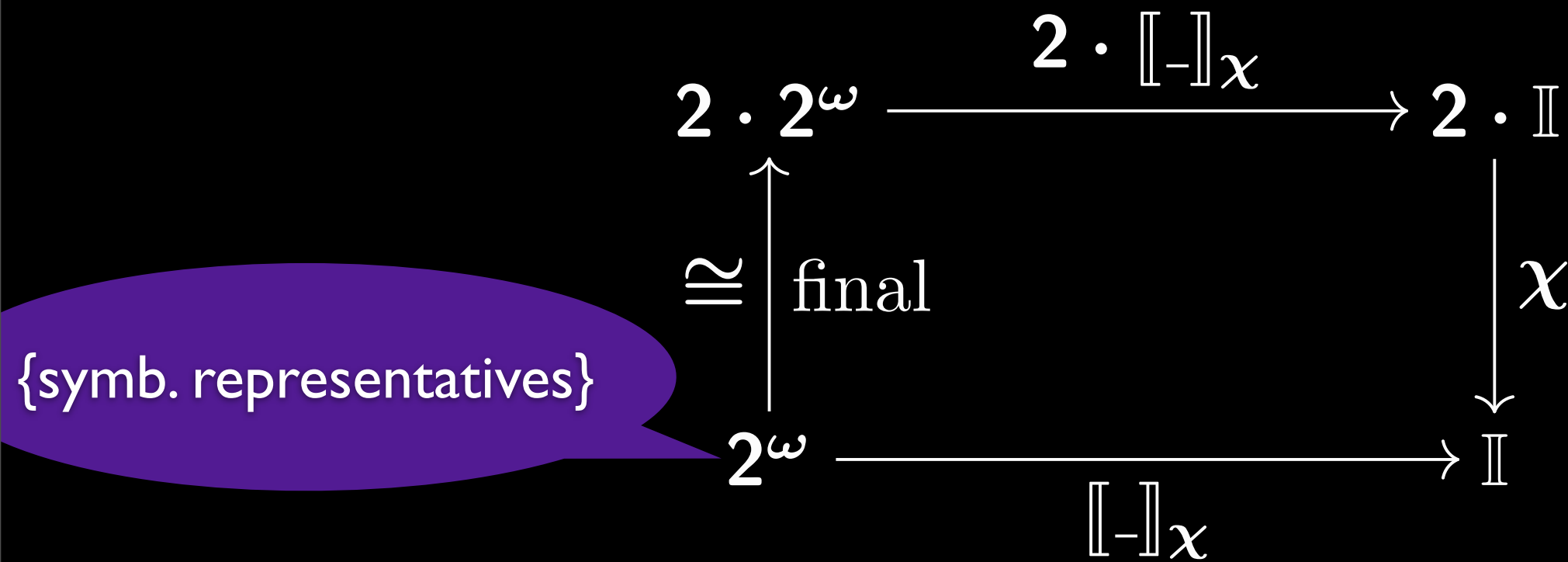
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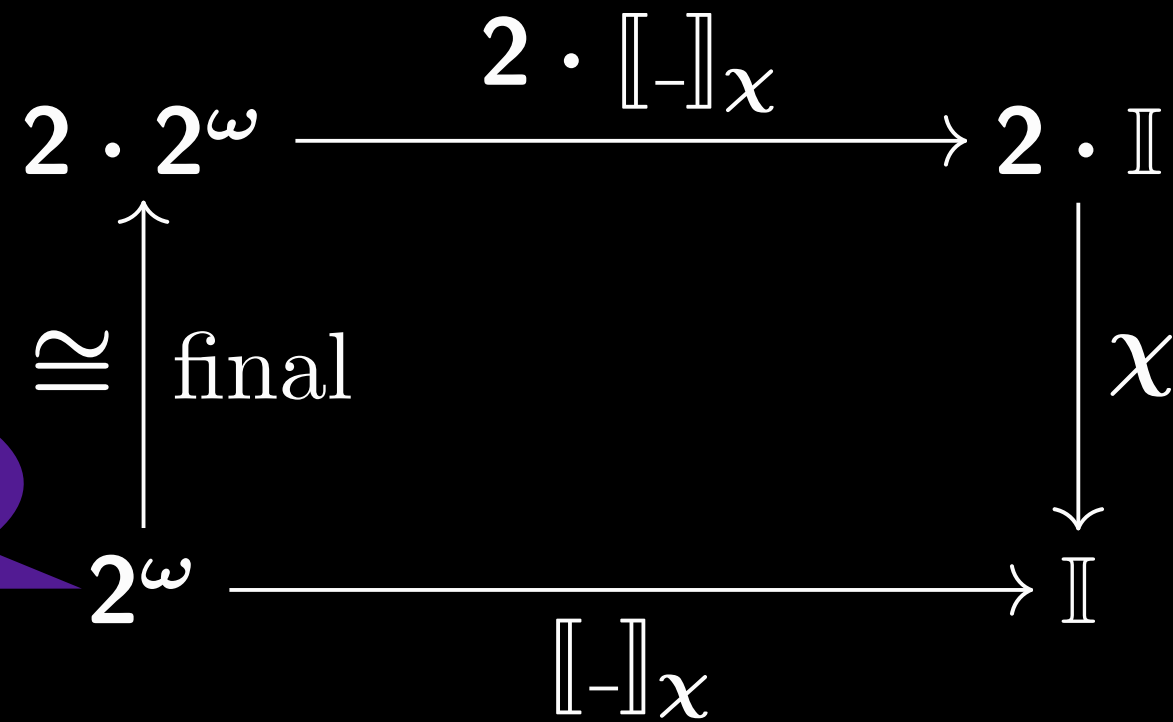


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0000...



1011...



# Observation 3

$$\begin{array}{ccc} 2 \cdot 2^\omega & \xrightarrow{2 \cdot \llbracket - \rrbracket_\chi} & 2 \cdot \mathbb{I} \\ \cong \uparrow \text{final} & & \downarrow \chi \\ 2^\omega & \xrightarrow{\llbracket - \rrbracket_\chi} & \mathbb{I} \end{array}$$

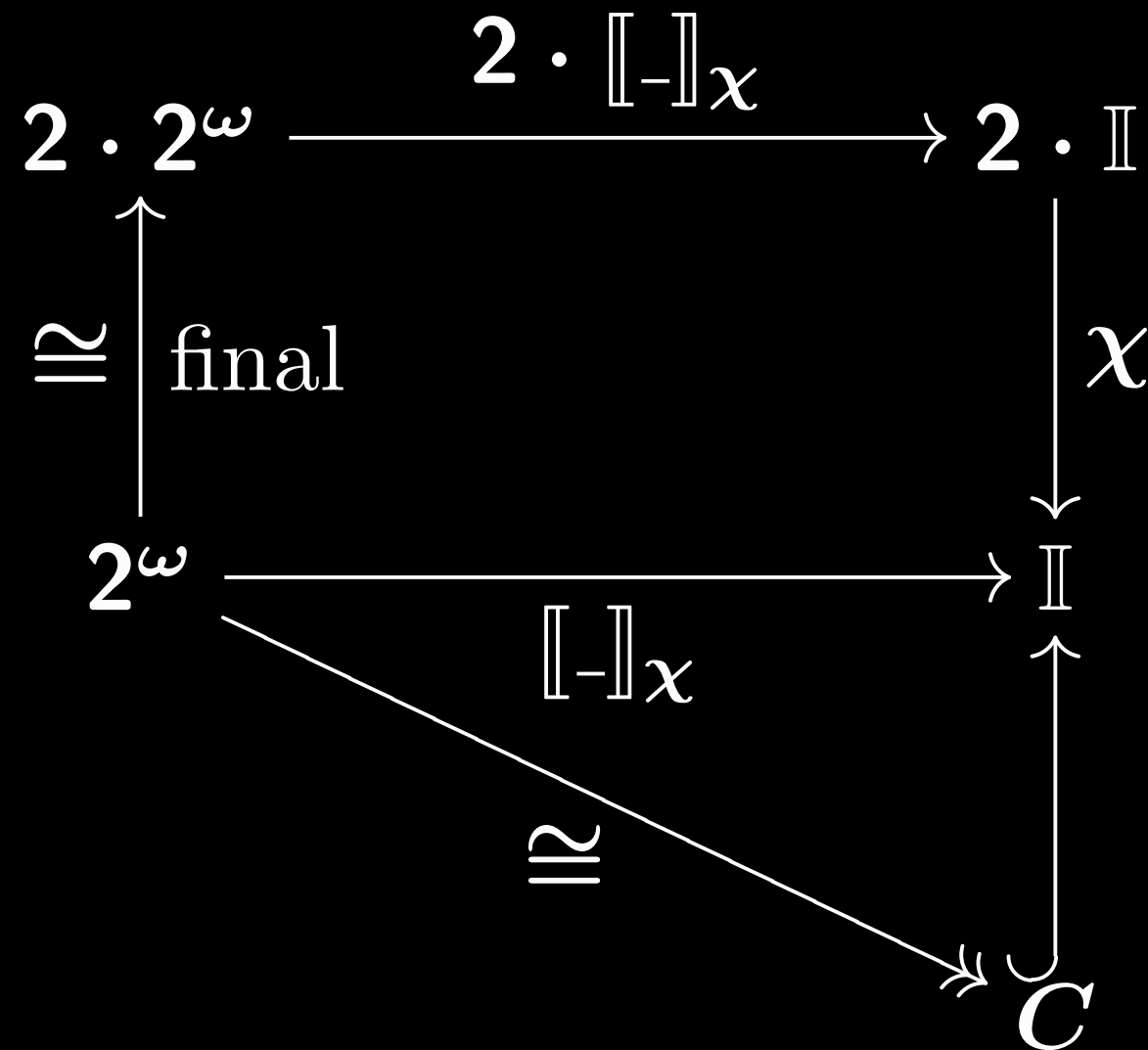
- Such  $\llbracket - \rrbracket_\chi$  uniquely exists.

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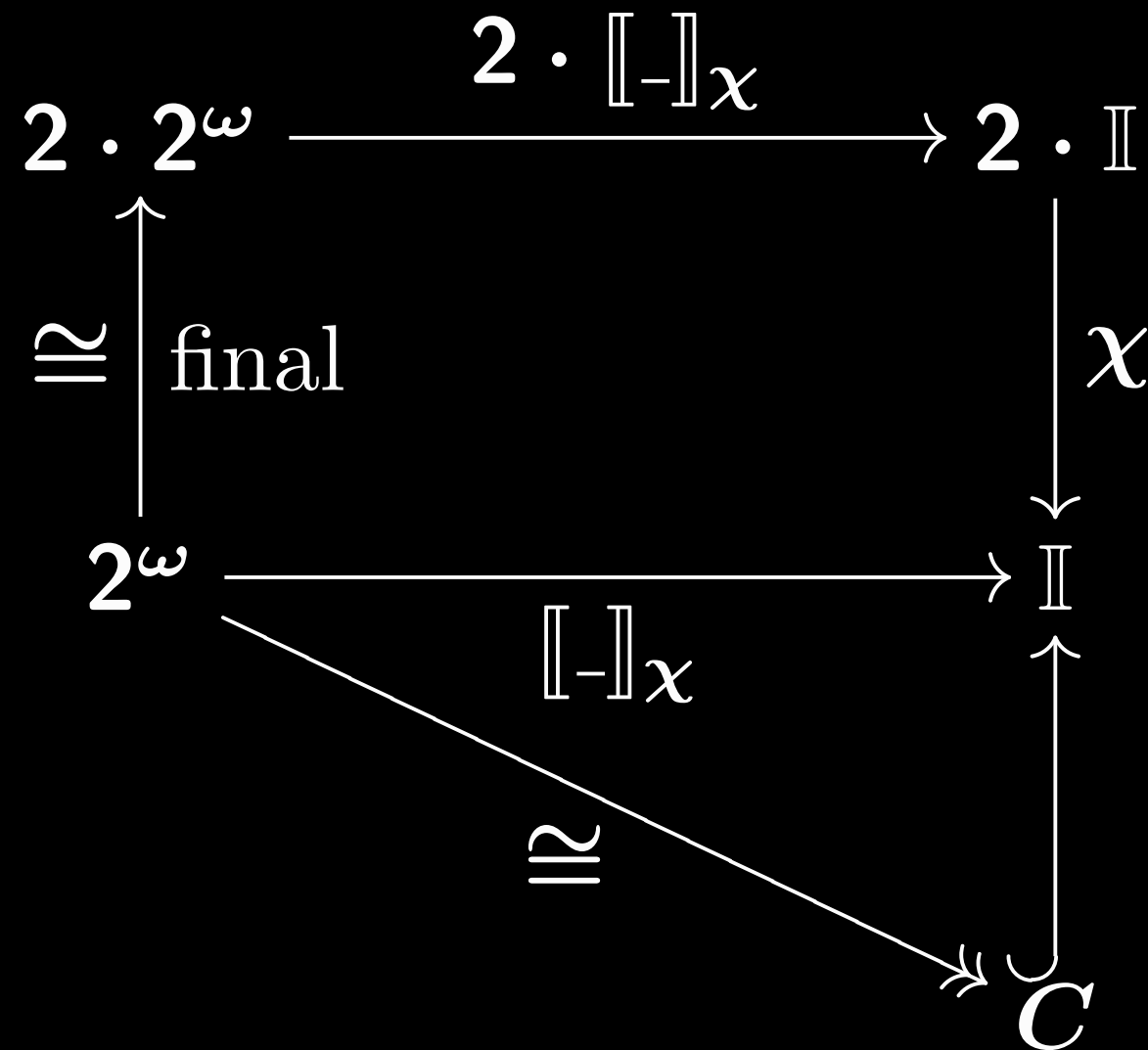
- $\llbracket - \rrbracket_\chi$  is injective.

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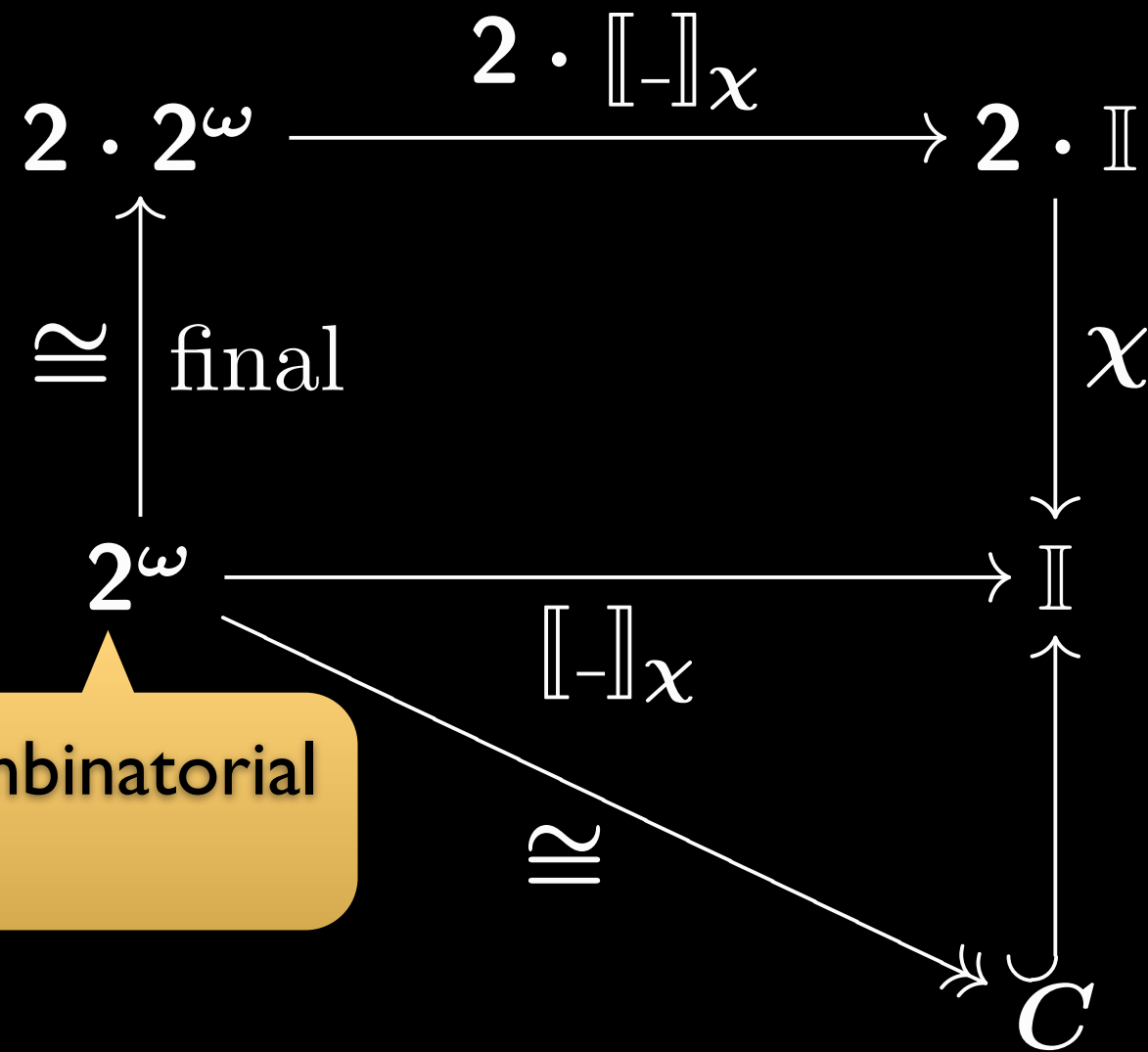


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# “Representation Theory”



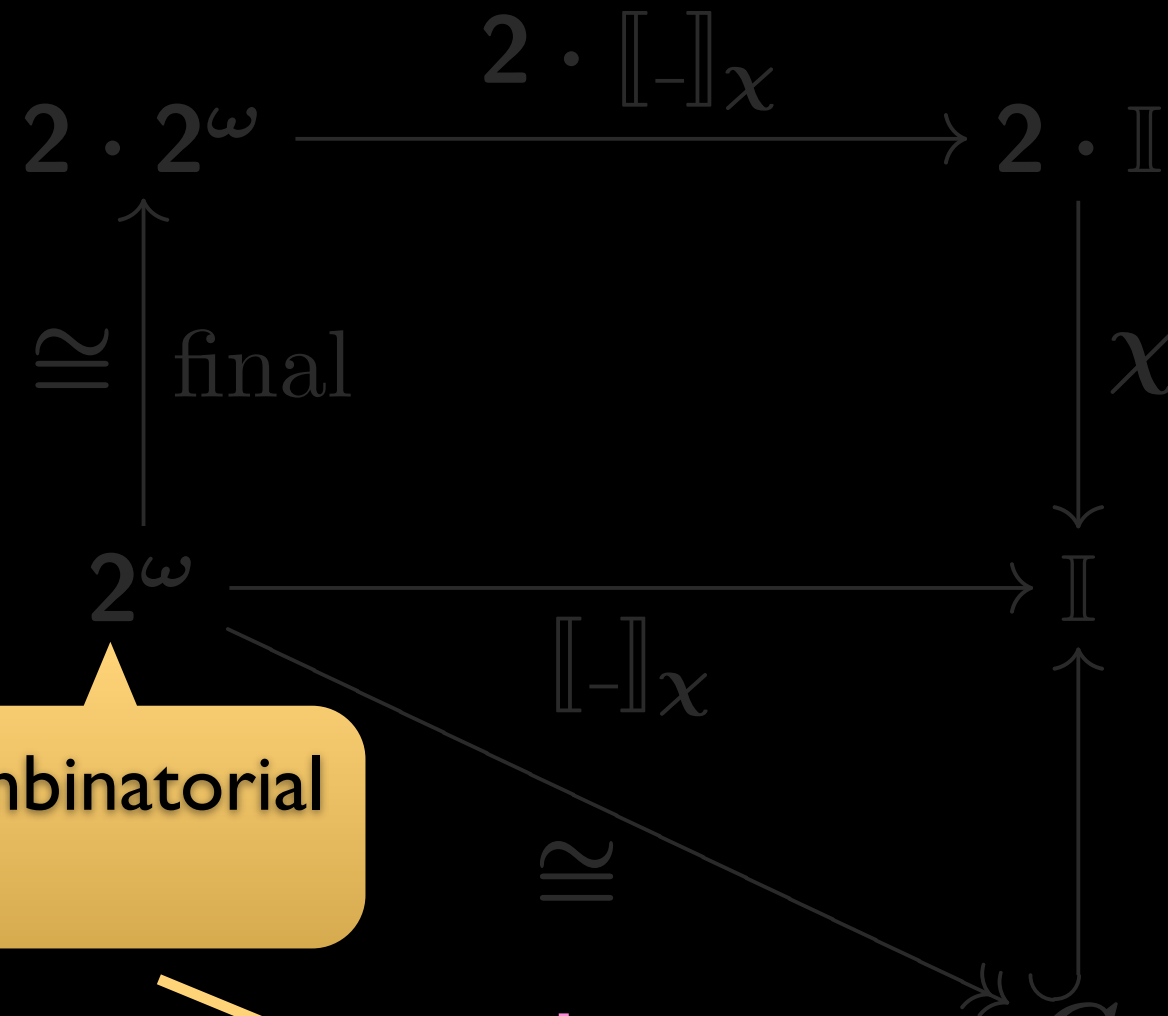
# “Representation Theory”



coalgebraic/symbolic/combinatorial fractal

metric fractal

# “Representation Theory”



coalgebraic/symbolic/combinatorial  
fractal

denotation map

representation map

metric fractal



# Axiomatic Domain Theory

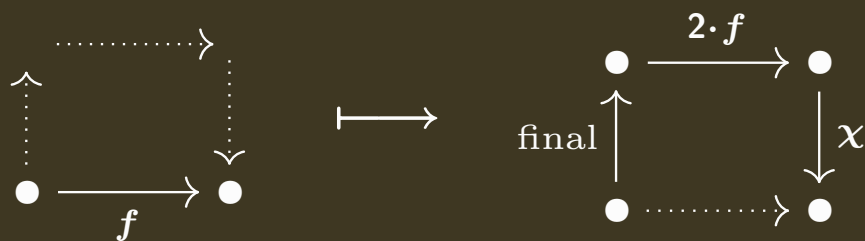
- **Thm.**  $\exists! \llbracket - \rrbracket_{\chi}$  s.t.

$$\begin{array}{ccc}
 2 \cdot 2^{\omega} & \xrightarrow{\quad 2 \cdot \llbracket - \rrbracket_{\chi} \quad} & 2 \cdot \mathbb{I} \\
 \cong \uparrow \text{final} & & \downarrow \chi \\
 2^{\omega} & \xrightarrow{\quad \llbracket - \rrbracket_{\chi} \quad} & \mathbb{I}
 \end{array}$$

## Proof

- $\text{Sets}(2^{\omega}, \mathbb{I})$  is a complete metric space (CMS).

- $\Phi : \text{Sets}(2^{\omega}, \mathbb{I}) \longrightarrow \text{Sets}(2^{\omega}, \mathbb{I})$



is a contracting map.

- Use the Banach fixed pt. thm.

# Axiomatic Domain Theory

- **Thm.**  $\exists! \llbracket - \rrbracket_{\mathcal{X}}$  s.t.

cf. *initial algebra-final coalgebra coincidence*

- solving domain equation [Freyd]  

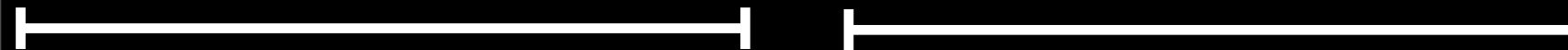
$$X = (X \Rightarrow X)$$
- coalgebraic trace semantics [IH-Jacobs-Sokolova]
- *corecursive algebra* [Capretta-Uustalu-Vene]
- typical in enrichment with “approximation structure”
  - of *infinitary* data, by *finitary* ones
  - order/CPO [Smyth-Plotkin], complete metric [America-Rutten]

$$\begin{array}{ccc}
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# Unit interval II



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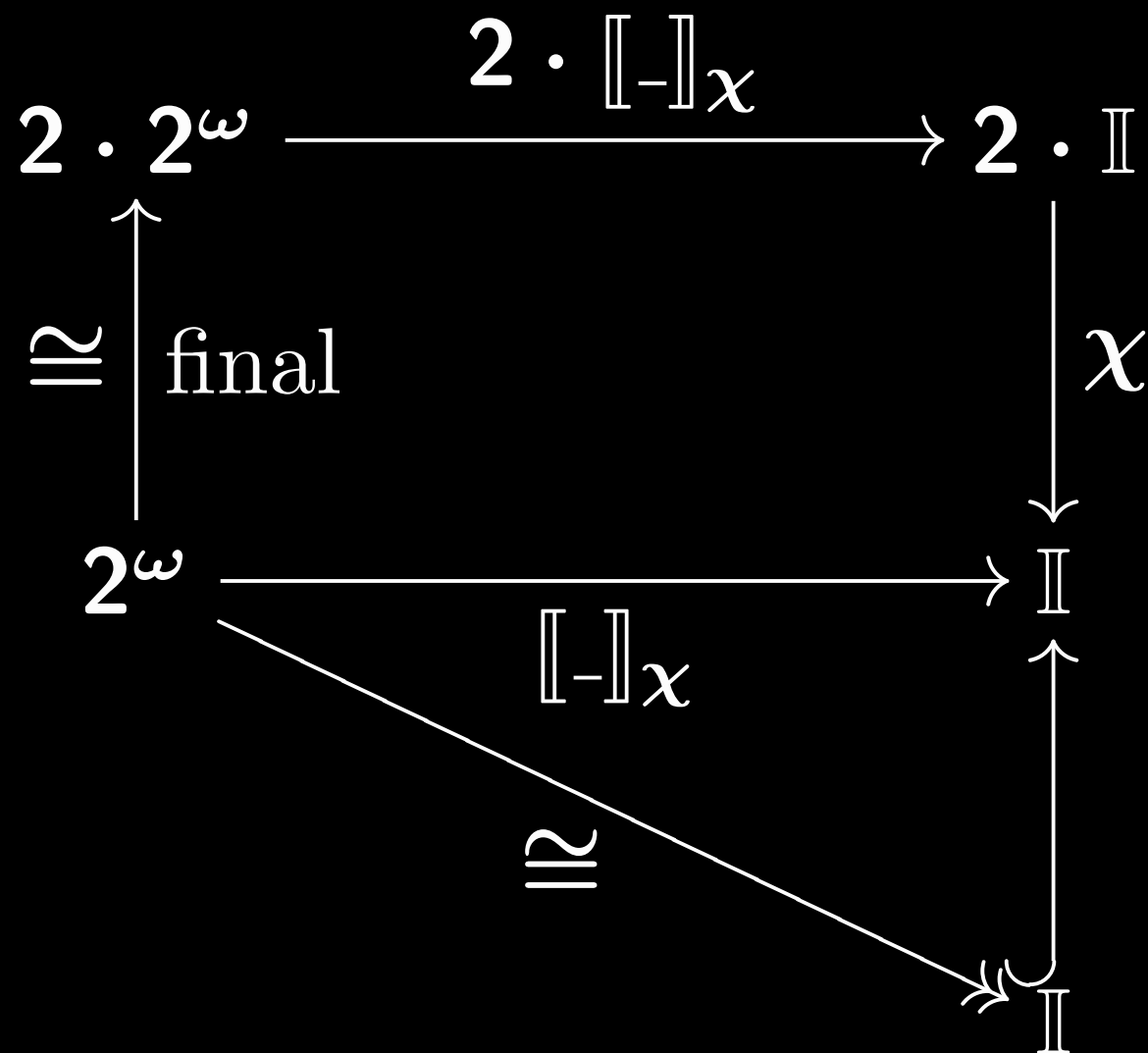


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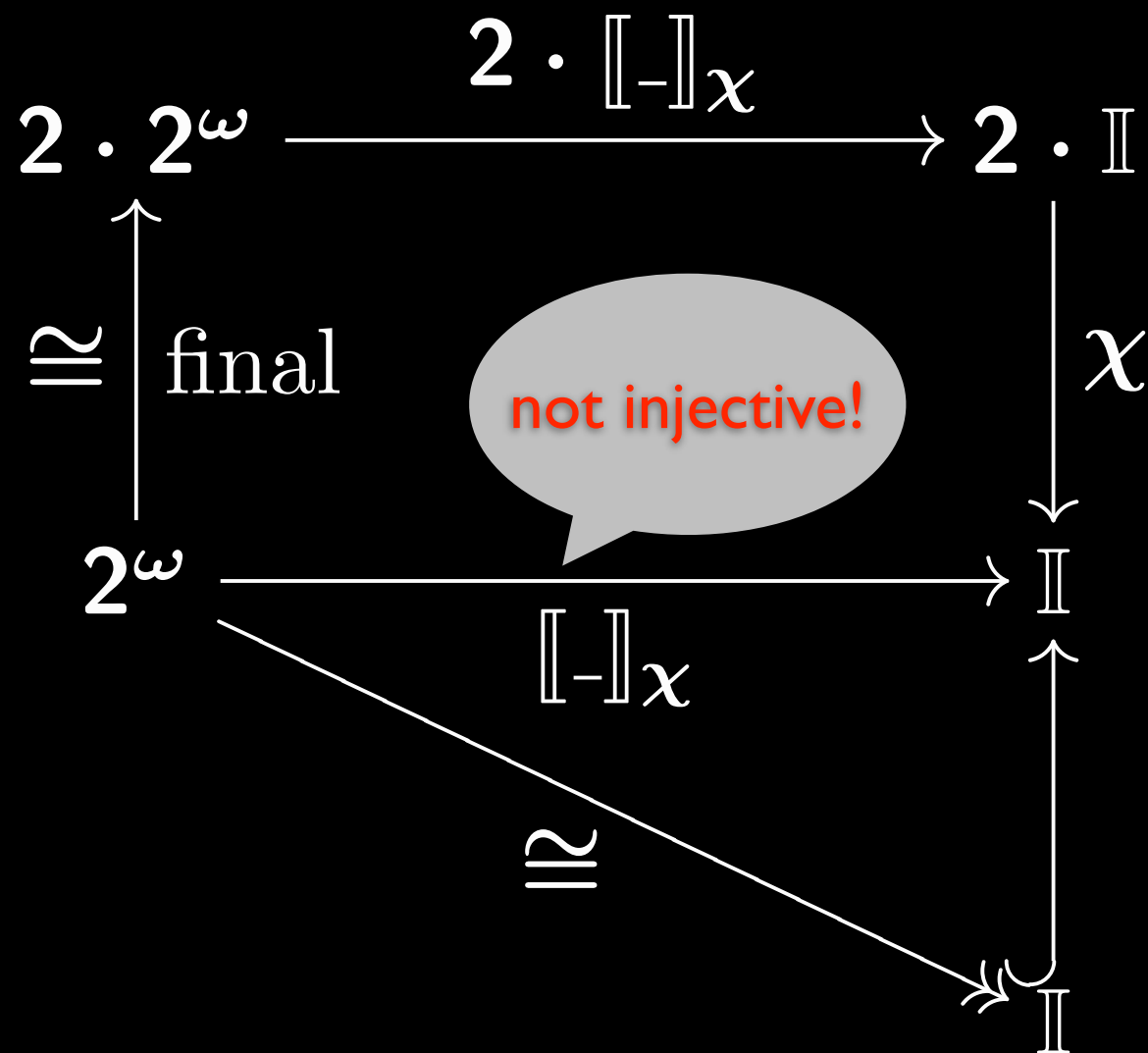
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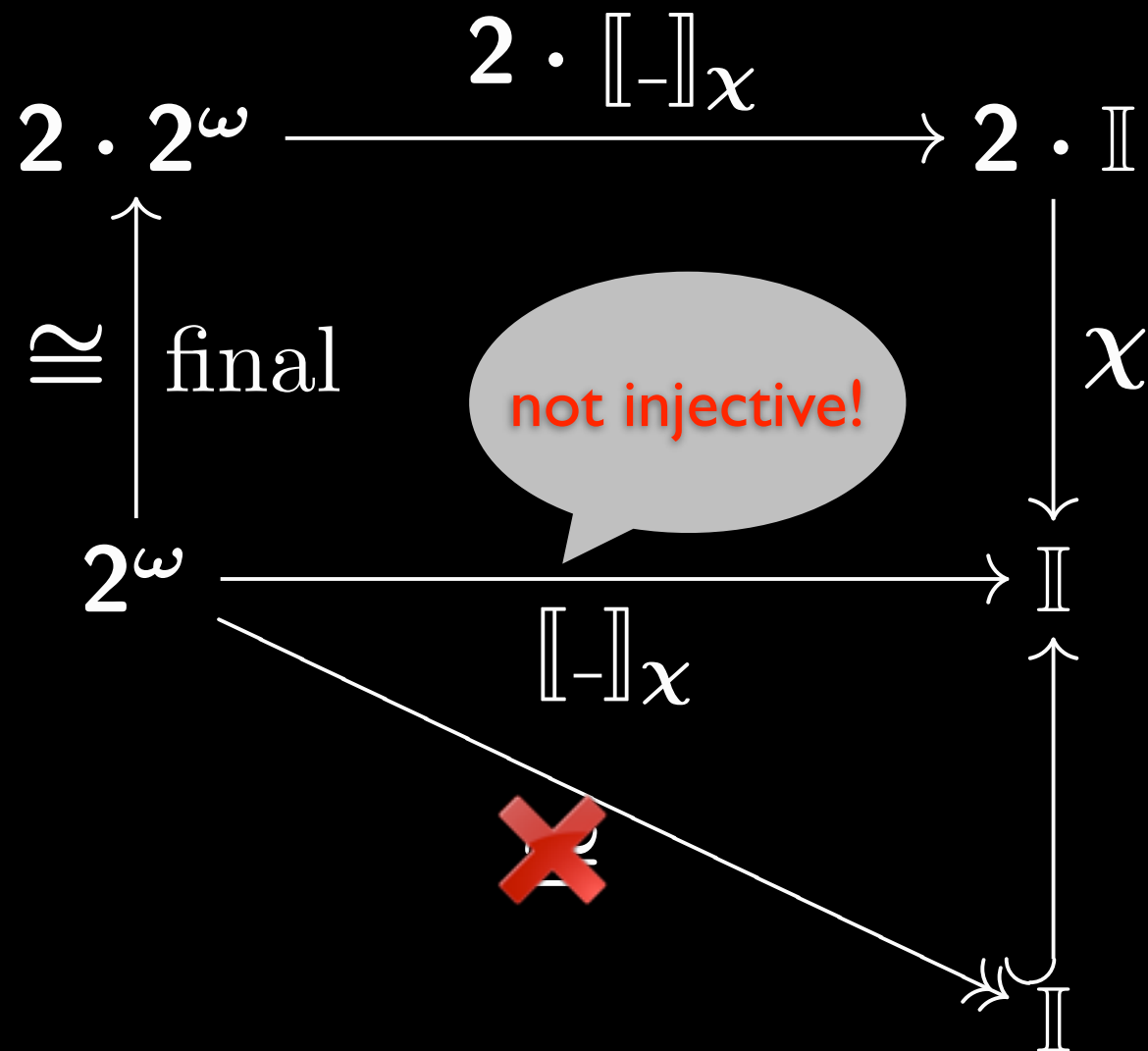
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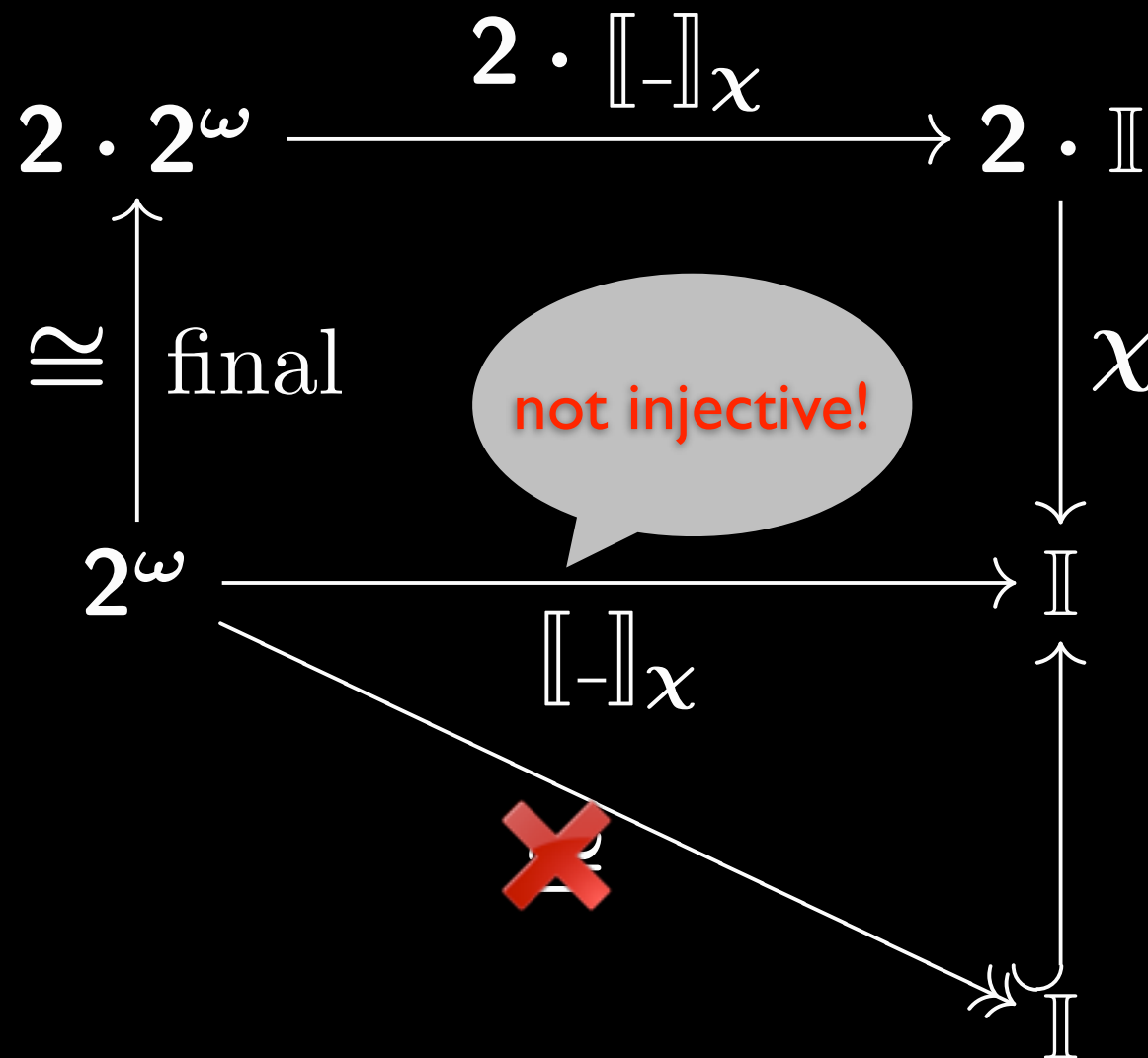
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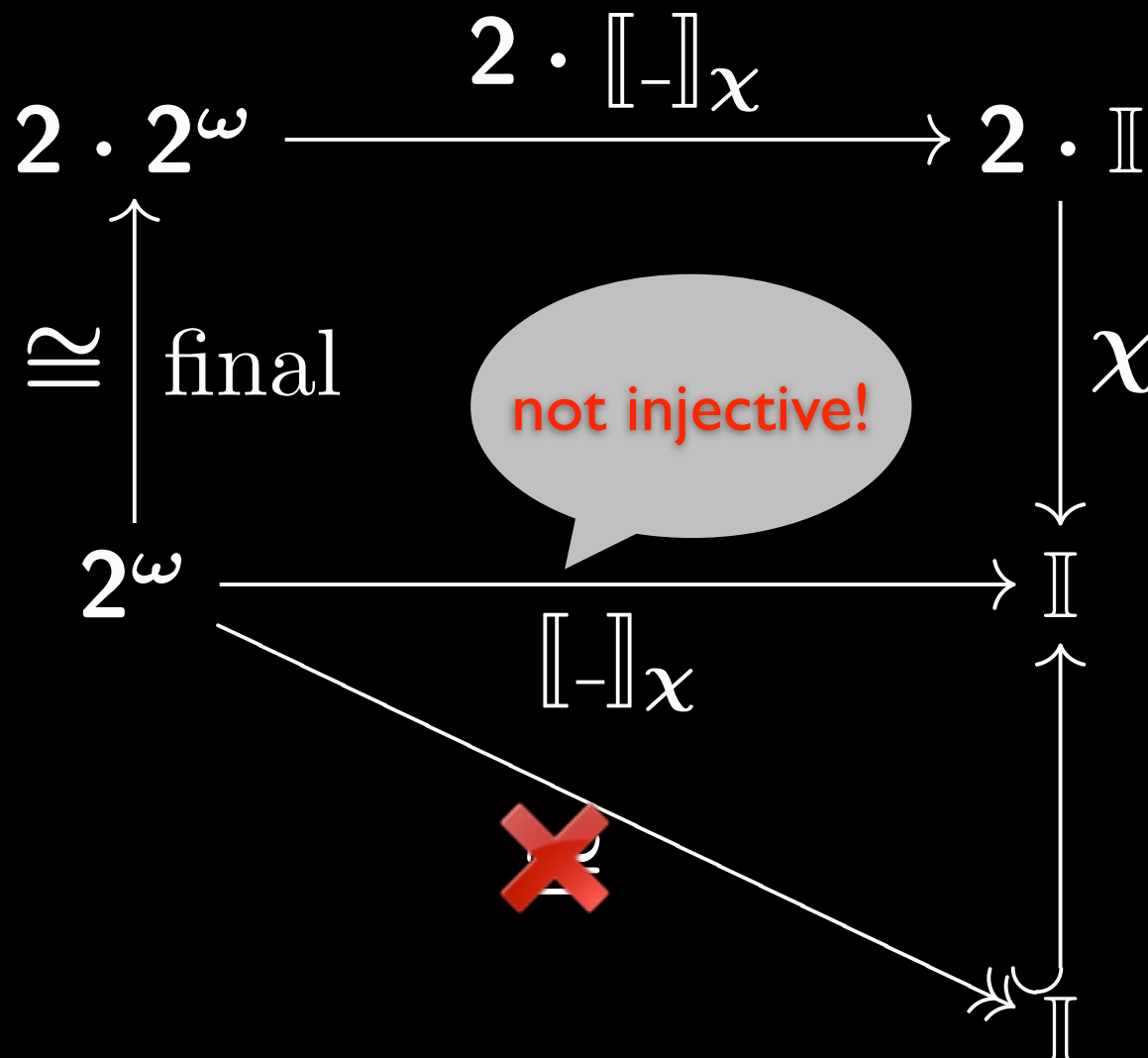
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# Unit interval II

gluing



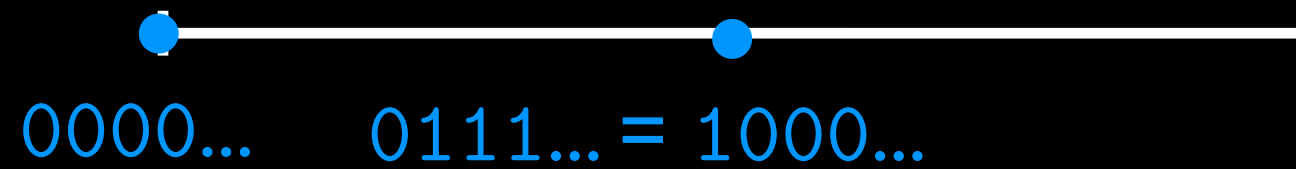
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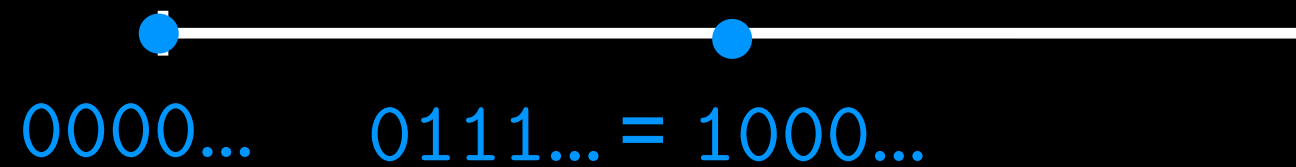
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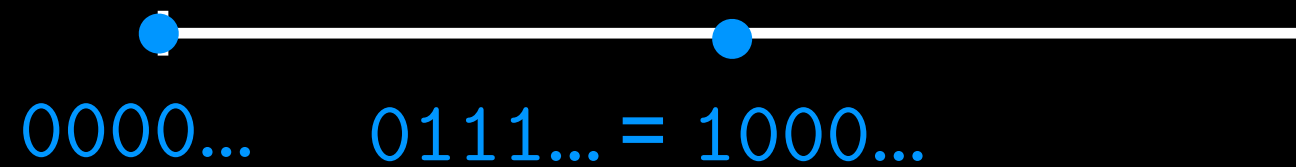


# Unit interval II



$$\{ \text{symbolic representatives} \} = 2^{\omega} / \sim$$

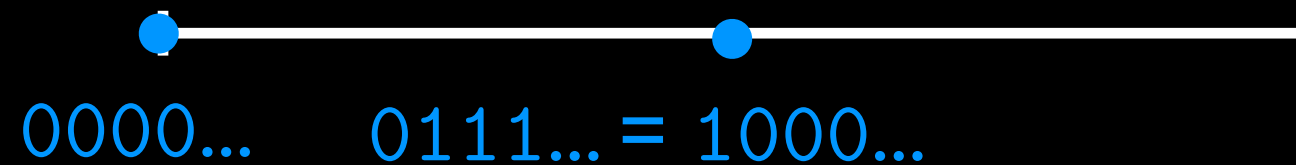
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categorycally?

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- Use of *presheaves* and *modules/distributors/profunctors*
  - Freyd's observation
  - Tom Leinster. A general theory of self-similarity I, II. In arXiv.



# Related Work

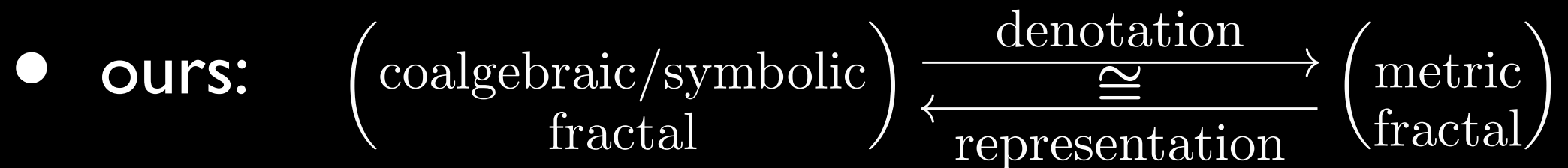
- Real number computation, exact arithmetic  
[Edalat-Heckmann, Kreitz-Weihrauch, Weihrauch, ...]

- “representation”: 
$$\begin{array}{ccc} \mathbb{N}^\omega & & \\ \uparrow & & \\ \mathbb{A} & \longrightarrow & \mathbb{I} \end{array}$$

- **ours:** 
$$\left( \begin{array}{c} \text{coalgebraic/symbolic} \\ \text{fractal} \end{array} \right) \begin{array}{c} \xrightarrow{\text{denotation}} \\ \xleftarrow{\text{representation}} \end{array} \begin{array}{c} \text{metric} \\ \text{fractal} \end{array}$$

# Related Work

- Real number computation, exact arithmetic  
[Edalat-Heckmann, Kreitz-Weihrauch, Weihrauch, ...]



- Real lines  $\mathbb{R}, \mathbb{I}, [0, 1), (0, 1), \dots$   
categorically/coalgebraically [Escardo-Simpson, Pavlovic-Pratt, ...]
- introducing the *fractal* point of view

# Related Work

- **Domain-theoretic approach to fractals**  
[Hayashi, Edalat, Scriven, Coquand, ...]

domain theory

semantics of  
recursion/“infinity”

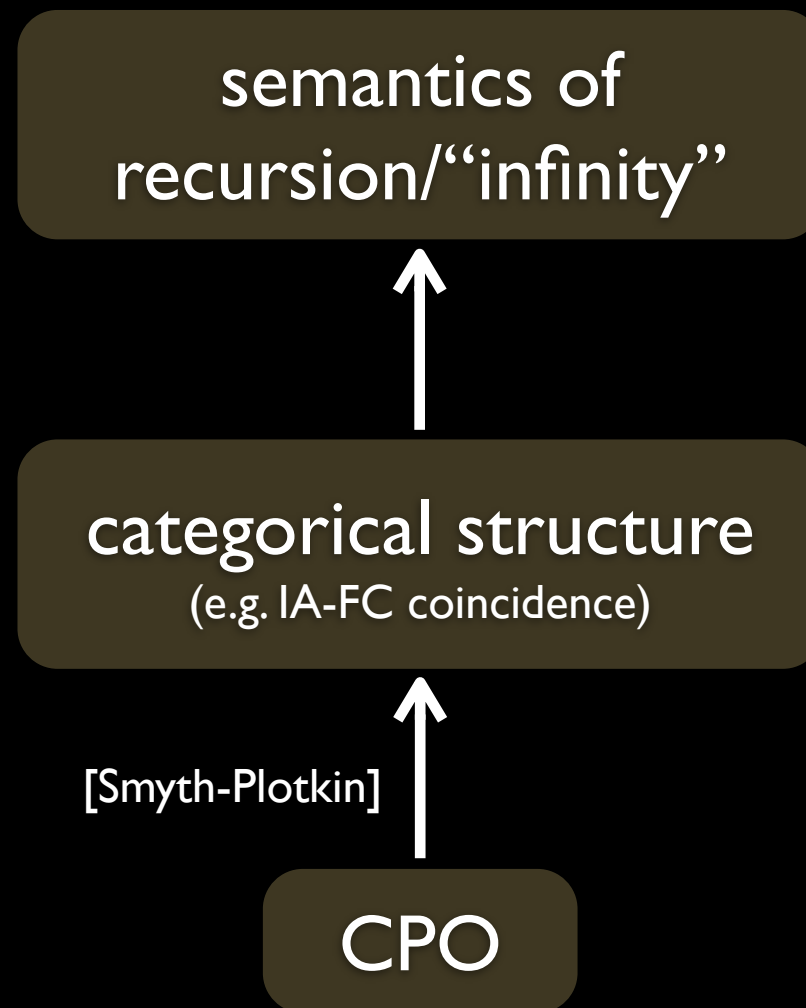
CPO



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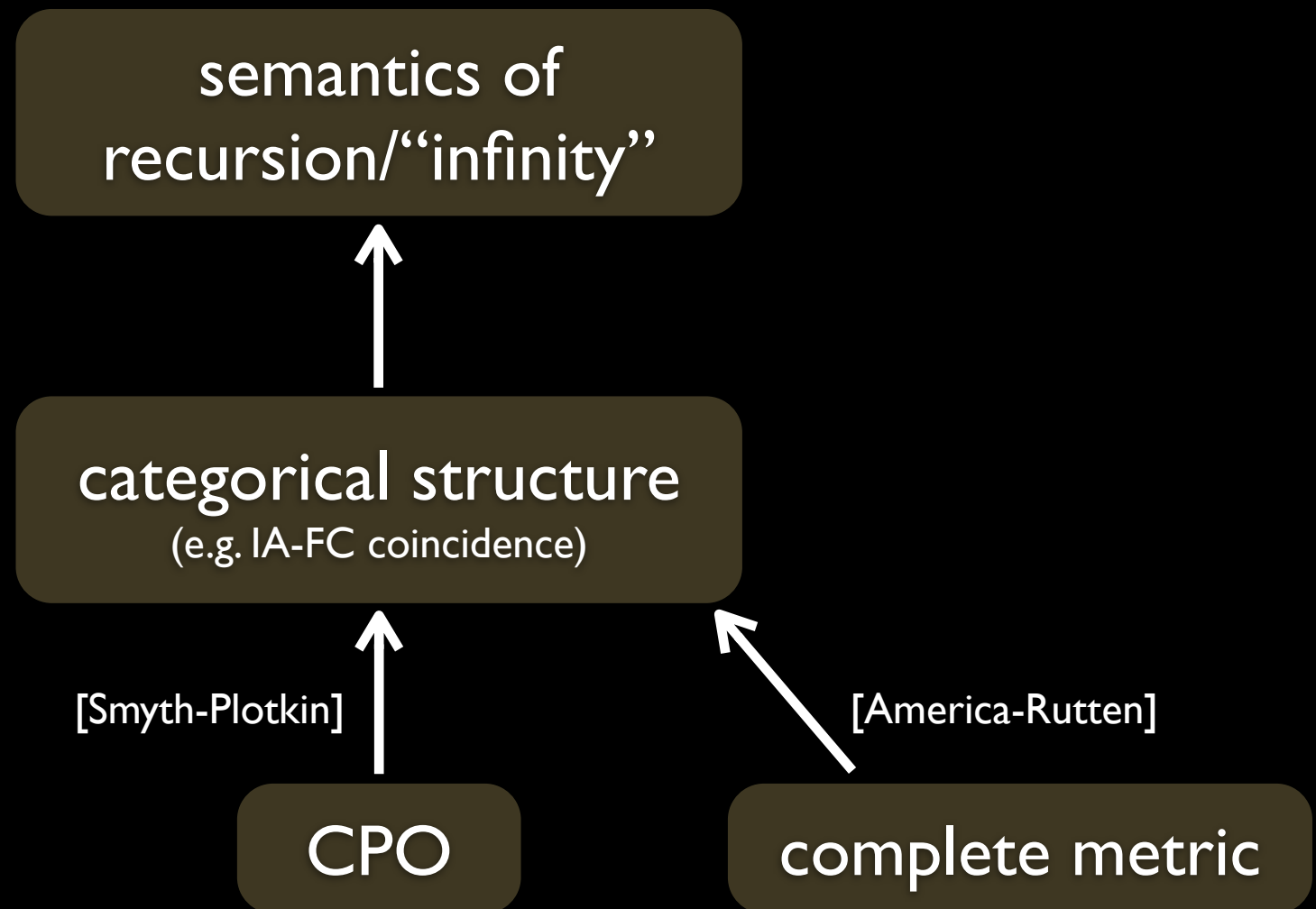
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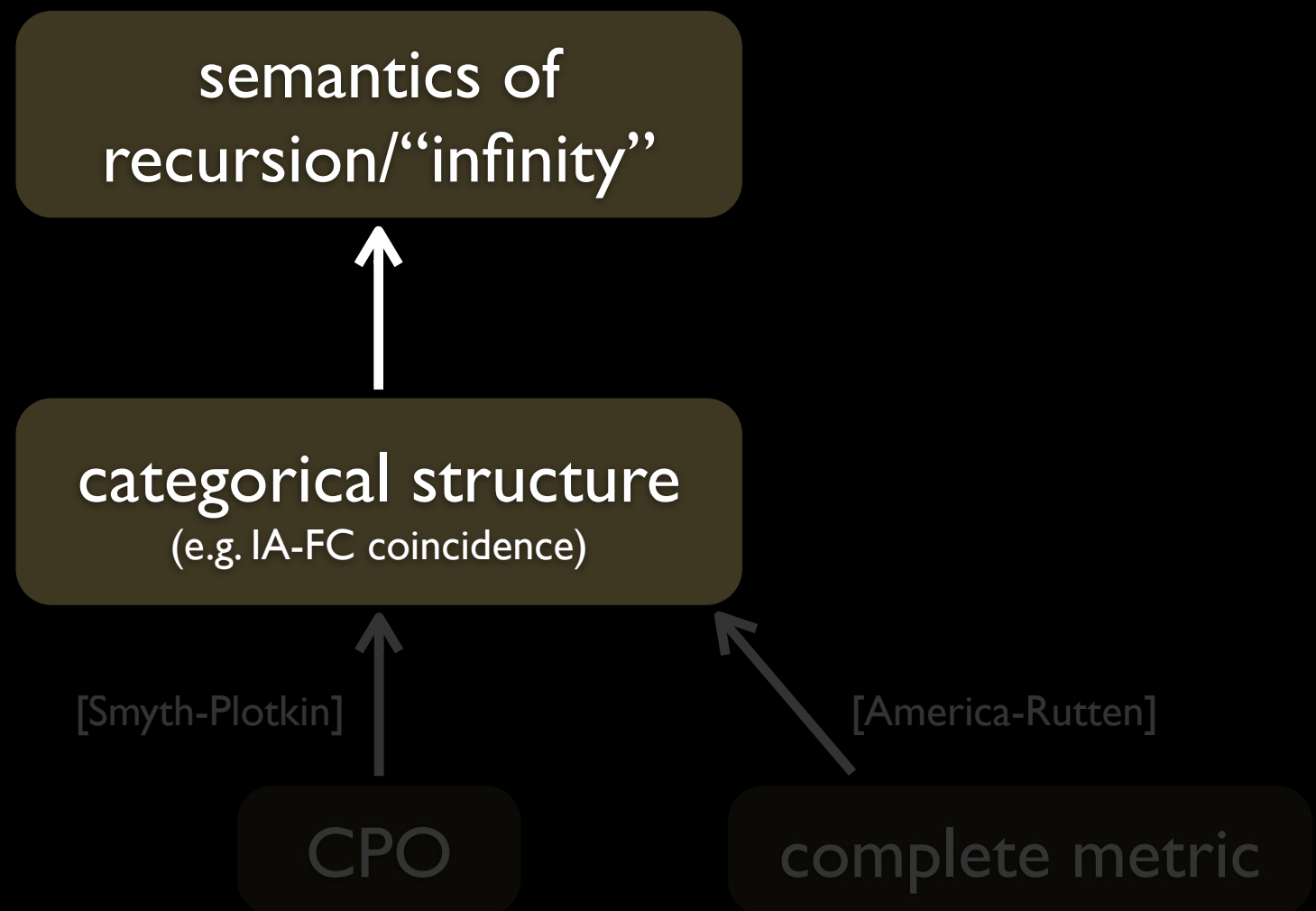
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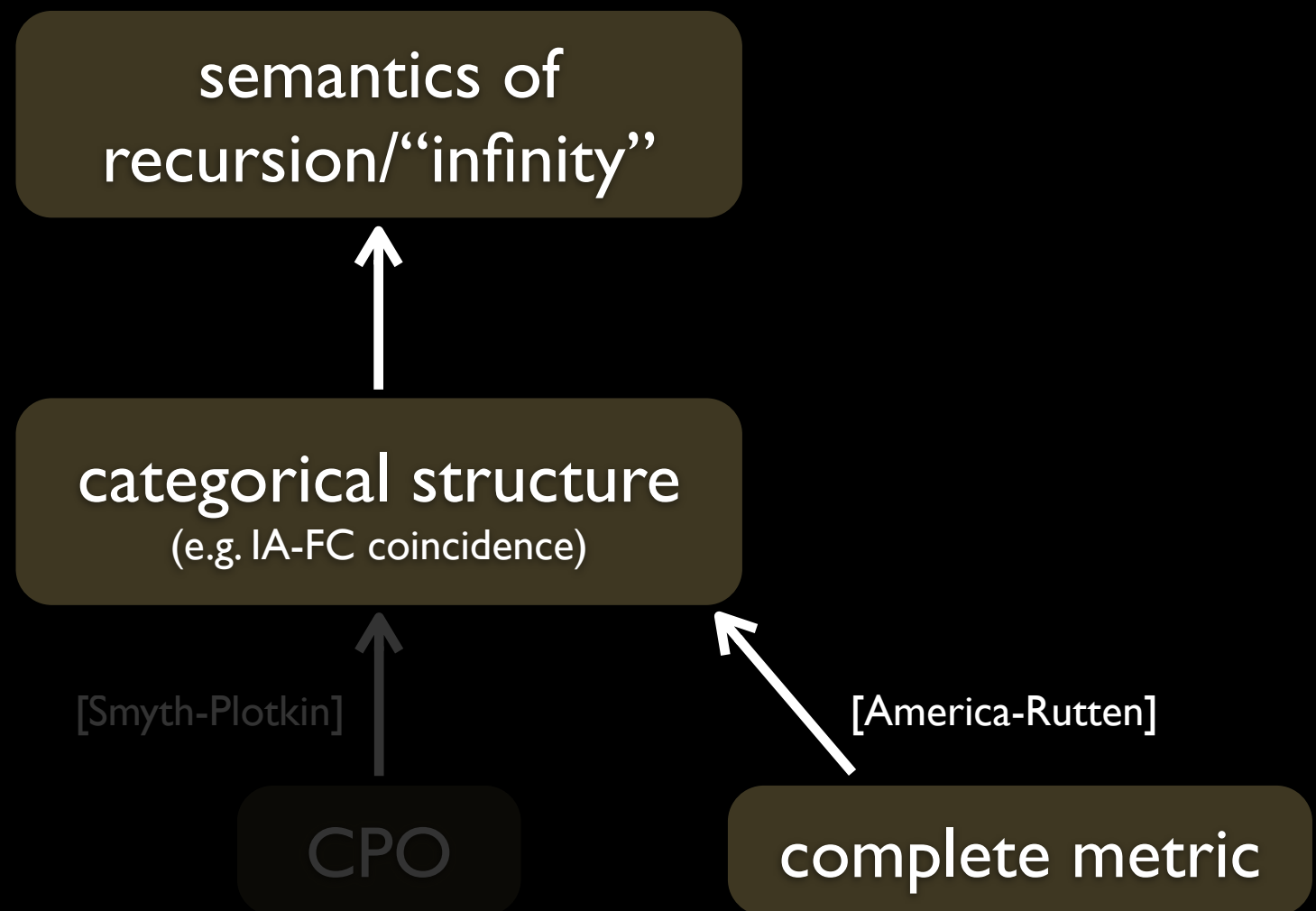
**coalgebraic  
modeling**  
[Rutten, Jacobs, ...]



# Related Work

- Domain-theoretic approach to fractals  
[Hayashi, Edalat, Scriven, Coquand, ...]

current work



# Related Work

- Tom Leinster.  
A general theory of self-similarity I, II.



# The Scenario

1. combinatorial spec.

$$\mathbf{n} \cdot (\_) : \mathbf{Sets} \rightarrow \mathbf{Sets}$$

2. symbolic  
representatives

$$\begin{array}{c} \mathbf{n} \cdot \mathbf{n}^\omega \\ \uparrow \text{final} \cong \\ \mathbf{n}^\omega \end{array}$$

3. IFS

$$\begin{array}{c} \mathbf{n} \cdot \mathbf{X} \\ \downarrow \chi \\ \mathbf{X} \end{array}$$

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$$\begin{array}{ccc} \mathbf{n} \cdot \mathbf{n}^\omega & \xrightarrow{\mathbf{n} \cdot [-]_\chi} & \mathbf{n} \cdot \mathbf{X} \\ \uparrow \cong \text{final} & & \downarrow \chi \\ \mathbf{n}^\omega & \xrightarrow{[-]_\chi} & \mathbf{X} \\ & \searrow \cong & \uparrow \\ & & \mathbf{Im}[-]_\chi \end{array}$$

with gluing

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# Fractal as Presheaf Coalgebra [Leinster]

- **Aim** Turn

$$\begin{aligned}\varphi_0, \varphi_1 &: \mathbb{C} \longrightarrow \mathbb{C}, \\ \varphi_0(x) &= \frac{x}{2}, \\ \varphi_1(x) &= \frac{1+x}{2}.\end{aligned}$$

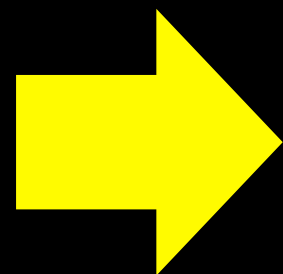
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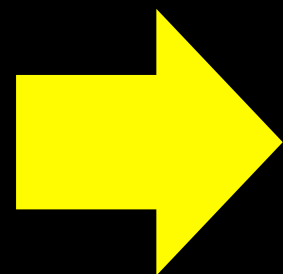
“mod out” via *presheaves* and *modules*

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- **Aim** Turn

$$\begin{aligned}\varphi_0, \varphi_1 &: \mathbb{C} \longrightarrow \mathbb{C}, \\ \varphi_0(x) &= \frac{x}{2}, \\ \varphi_1(x) &= \frac{1+x}{2}.\end{aligned}$$

into an *injective*  $\mathcal{X}$



“mod out” via *presheaves* and *modules*

In this talk:

- presheaf  
 $P : \mathbb{A} \rightarrow \mathbf{Sets}$
- module  
$$\frac{M : \mathbb{A} \dashrightarrow \mathbb{B}}{\hline M : \mathbb{A}^{\text{op}} \times \mathbb{B} \longrightarrow \mathbf{Sets}}$$

# Module

(bimodule/distributor/profunctor)

$M : \mathbb{A} \multimap \mathbb{A}$  , a module

---

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$M : \mathbb{A}^{\text{op}} \times \mathbb{A} \longrightarrow \mathbf{Sets}$  , a functor

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(bimodule/distributor/profunctor)

$$\frac{M : \mathbb{A} \multimap \mathbb{A} , \text{ a module}}{\hline M : \mathbb{A}^{\text{op}} \times \mathbb{A} \longrightarrow \mathbf{Sets} , \text{ a functor}}$$

- $\frac{\text{relation}}{\text{function}} = \frac{\text{module}}{\text{functor}}$

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$$\frac{M : \mathbb{A} \multimap \mathbb{A} \text{ , a module}}{\underline{\underline{M : \mathbb{A}^{\text{op}} \times \mathbb{A} \longrightarrow \mathbf{Sets} \text{ , a functor}}}}$$

- $\frac{\text{relation}}{\text{function}} = \frac{\text{module}}{\text{functor}}$

- *sets with left and right  $\mathbb{A}$ -action*

$$M(a, b) = \{ \overset{a}{\longrightarrow} \boxed{m} \overset{b}{\longrightarrow} \}$$

$$M(f, g)m = g \cdot m \cdot f = \overset{a'}{\longrightarrow} \circledast f \overset{a}{\longrightarrow} \boxed{m} \overset{b}{\longrightarrow} \circledast g \overset{b'}{\longrightarrow}$$



# Module

(bimodule/distributor/profunctor)

$$\frac{M : \mathbb{A} \multimap \mathbb{A} , \text{ a module}}{\underline{\underline{M : \mathbb{A}^{\text{op}} \times \mathbb{A} \longrightarrow \mathbf{Sets} , \text{ a functor}}}}$$

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# The Base Category $\mathbb{A}_{\mathbb{I}}$

$$\mathbb{A}_{\mathbb{I}} = \left( \mathbf{0} \begin{array}{c} \xrightarrow{l} \\ \xrightarrow{r} \end{array} \mathbf{1} \right)$$

# The Base Category $\mathbb{A}_{\mathbb{I}}$

$$\mathbb{A}_{\mathbb{I}} = \left( \mathbf{0} \begin{array}{c} \xrightarrow{l} \\ \xrightarrow{r} \end{array} \mathbf{1} \right)$$

$$P_{\mathbb{I}} : \mathbb{A}_{\mathbb{I}} \longrightarrow \mathbf{Sets}$$

$$\begin{array}{ccc} \mathbf{0} & & \{*\} \\ l \downarrow \downarrow r & \longmapsto & \mathbf{0} \downarrow \downarrow \mathbf{1} \\ \mathbf{1} & & \mathbb{I} \end{array}$$

# The Base Category $\mathbb{A}_{\mathbb{I}}$

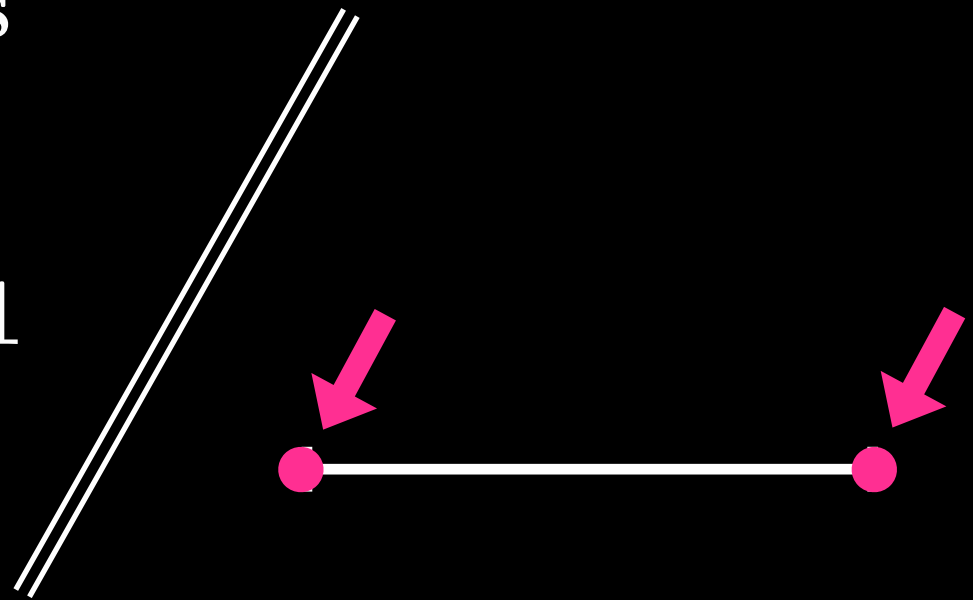
$$\mathbb{A}_{\mathbb{I}} = \left( \mathbf{0} \begin{array}{c} \xrightarrow{l} \\ \xrightarrow{r} \end{array} \mathbf{1} \right)$$

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$$\begin{array}{c} \mathbf{0} \\ l \downarrow \downarrow r \\ \mathbf{1} \end{array}$$

$\longmapsto$

$$\begin{array}{c} \{*\} \\ \mathbf{0} \downarrow \downarrow \mathbf{1} \\ \mathbb{I} \end{array}$$



# Combinatorial Specification as a Module



$$\frac{M_{\mathbb{I}} : \mathbb{A}_{\mathbb{I}} \dashrightarrow \mathbb{A}_{\mathbb{I}}}{\hline M_{\mathbb{I}} : \mathbb{A}_{\mathbb{I}}^{\text{op}} \times \mathbb{A}_{\mathbb{I}} \longrightarrow \text{Sets}}$$

gluing, indeed!

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$$\begin{array}{ccc}
 M_{\mathbb{I}}(\mathbf{0}, \mathbf{0}) & \begin{array}{c} \xrightarrow{l \cdot \_} \\ \xleftarrow{r \cdot \_} \end{array} & M_{\mathbb{I}}(\mathbf{0}, \mathbf{1}) \\
 \begin{array}{c} \_ \cdot l \uparrow \uparrow \_ \cdot r \\ \uparrow \uparrow \end{array} & & \uparrow \uparrow \\
 M_{\mathbb{I}}(\mathbf{1}, \mathbf{0}) & \xrightarrow{\quad} & M_{\mathbb{I}}(\mathbf{1}, \mathbf{1})
 \end{array}
 =
 \begin{array}{ccc}
 \{ \bullet \} & \begin{array}{c} \xrightarrow{0} \\ \xleftarrow{1} \end{array} & \{ \bullet \dashv, \dashv \bullet, \dashv \bullet \} \\
 \uparrow \uparrow & & \text{inf} \uparrow \uparrow \text{sup} \\
 \emptyset & \xrightarrow{\quad} & \{ \mathbf{H} \dashv, \dashv \mathbf{H} \}
 \end{array}$$

gluing, indeed!

# Combinatorial Specification as a Module



$$M_{\mathbb{I}} : \mathbb{A}_{\mathbb{I}} \dashv\vdash \mathbb{A}_{\mathbb{I}}$$

$$M(a, b) = \{ \overset{a}{\rightarrow} \boxed{m} \rightarrow^b \}$$

how many *ingredients* (*a*-shapes) are used in the *outcome* (*b*-shape)

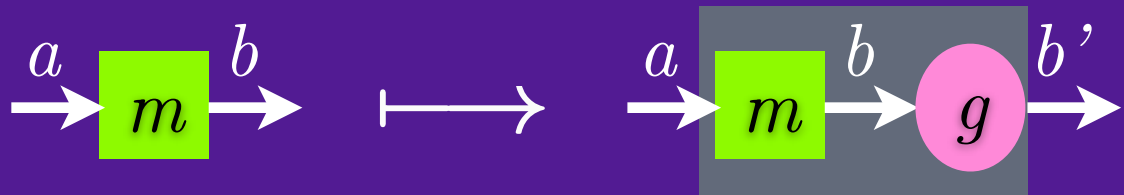
$$\begin{array}{ccc}
 M_{\mathbb{I}}(\mathbf{0}, \mathbf{0}) \begin{array}{c} \xrightarrow{l \cdot \_} \\ \xleftarrow{r \cdot \_} \end{array} M_{\mathbb{I}}(\mathbf{0}, \mathbf{1}) & = & \{ \bullet \} \begin{array}{c} \xrightarrow{0} \\ \xleftarrow{1} \end{array} \{ \bullet \dashv, \dashv \bullet, \dashv \bullet \} \\
 \begin{array}{c} \_ \cdot l \uparrow \uparrow \_ \cdot r \\ \uparrow \uparrow \end{array} & & \begin{array}{c} \uparrow \uparrow \\ \text{inf} \uparrow \uparrow \text{sup} \end{array} \\
 M_{\mathbb{I}}(\mathbf{1}, \mathbf{0}) \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} M_{\mathbb{I}}(\mathbf{1}, \mathbf{1}) & & \emptyset \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} \{ \mathbf{H} \dashv, \dashv \mathbf{H} \}
 \end{array}$$

gluing, indeed!

# Combinatorial Specification as a Module

$$M_{\mathbb{I}} : \mathbb{A}_{\mathbb{I}} \dashrightarrow \mathbb{A}_{\mathbb{I}}$$

$$M(\text{id}, g) = g \cdot \_ :$$



$$M(a, b) = \{ \overset{a}{\rightarrow} \boxed{m} \overset{b}{\rightarrow} \}$$

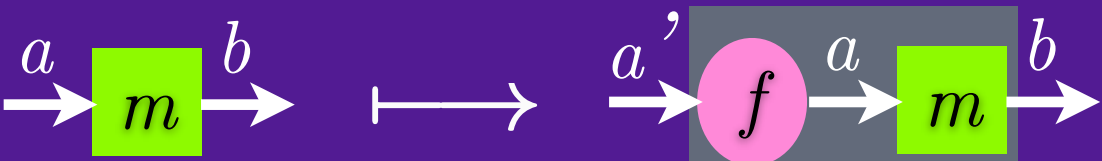
how many *ingredients* (*a*-shapes) are used in the *outcome* (*b*-shape)

$$M_{\mathbb{I}}(\mathbf{0}, \mathbf{0}) \xrightleftharpoons[l \cdot \_]{r \cdot \_} M_{\mathbb{I}}(\mathbf{0}, \mathbf{1})$$

$$= \begin{array}{ccc} \{ \bullet \} & \xrightleftharpoons[1]{0} & \{ \bullet \dashrightarrow, \dashrightarrow \bullet, \dashrightarrow \bullet \} \\ \uparrow \uparrow & & \text{inf} \uparrow \uparrow \text{sup} \\ \emptyset & \xrightleftharpoons{\quad} & \{ \mathbf{H} \dashrightarrow, \dashrightarrow \mathbf{H} \} \end{array}$$

$$M_{\mathbb{I}}(\mathbf{1}, \mathbf{0}) \xrightleftharpoons{\quad} M_{\mathbb{I}}(\mathbf{1}, \mathbf{1})$$

$$M(f, \text{id}) = \_ \cdot f :$$

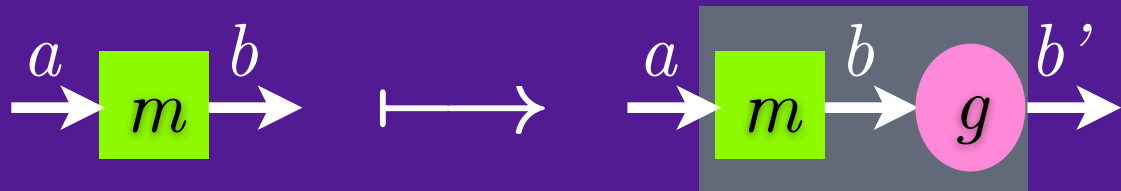


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how many *ingredients* ( $a$ -shapes) are used in the *outcome* ( $b$ -shape)

gluing

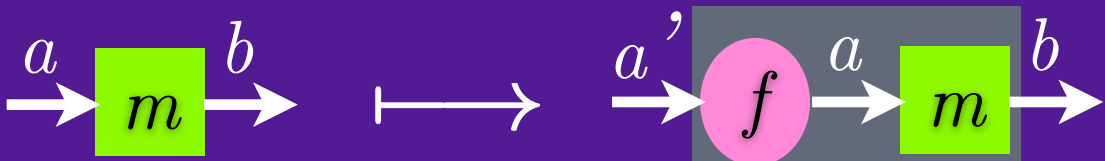
$$M_{\mathbb{I}}(\mathbf{0}, \mathbf{0}) \xrightleftharpoons[l \cdot \_]{r \cdot \_} M_{\mathbb{I}}(\mathbf{0}, \mathbf{1})$$

$$\_ \cdot l \uparrow \uparrow \_ \cdot r$$

$$= \begin{matrix} \{ \bullet \} & \xrightleftharpoons[1]{0} & \{ \bullet \dashrightarrow, \dashrightarrow \bullet, \dashrightarrow \bullet \} \\ \uparrow \uparrow & & \text{inf} \uparrow \uparrow \text{sup} \\ \emptyset & \xrightleftharpoons{\quad} & \{ \mathbf{H} \dashrightarrow, \dashrightarrow \mathbf{H} \} \end{matrix}$$

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$$M(f, \text{id}) = \_ \cdot f :$$



gluing, indeed!

# The Scenario

1. combinatorial spec.

$$\mathbf{n} \cdot (\_) : \mathbf{Sets} \rightarrow \mathbf{Sets}$$

2. symbolic  
representatives

$$\begin{array}{c} \mathbf{n} \cdot \mathbf{n}^\omega \\ \uparrow \text{final} \\ \mathbf{n}^\omega \end{array} \cong$$

3. IFS

$$\begin{array}{c} \mathbf{n} \cdot \mathbf{X} \\ \downarrow \chi \\ \mathbf{X} \end{array}$$

4.

$$\begin{array}{ccc} \mathbf{n} \cdot \mathbf{n}^\omega & \xrightarrow{\mathbf{n} \cdot [-]_\chi} & \mathbf{n} \cdot \mathbf{X} \\ \uparrow \cong \text{final} & & \downarrow \chi \\ \mathbf{n}^\omega & \xrightarrow{[-]_\chi} & \mathbf{X} \\ & \searrow \cong & \uparrow \\ & & \mathbf{Im}[-]_\chi \end{array}$$

1. combinatorial spec.

$$(\mathbb{A}, M : \mathbb{A} \leftrightarrow \mathbb{A})$$

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with gluing

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# Tensor Product

$$M \otimes (\_ ) : \text{Sets}^{\mathbb{A}} \longrightarrow \text{Sets}^{\mathbb{A}}$$

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$$P \longmapsto M \otimes P$$



# Tensor Product

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$$P \longmapsto M \otimes P$$

$$M(a, b) = \left\{ \begin{array}{c} a \longrightarrow \boxed{m} \longrightarrow b \end{array} \right\}$$

$$P(a) = \left\{ \begin{array}{c} \triangleleft x \longrightarrow a \end{array} \right\}$$

# Tensor Product

$$M \otimes (\_): \text{Sets}^{\mathbb{A}} \longrightarrow \text{Sets}^{\mathbb{A}}$$

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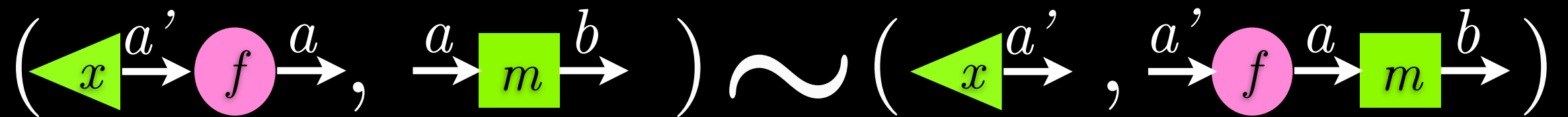
$$P(a) = \left\{ \begin{array}{c} \triangleleft x \longrightarrow a \end{array} \right\}$$

---


$$(M \otimes P)b = \left( \coprod_{a \in \mathbb{A}} P(a) \times M(a, b) \right) / \sim$$

$$= \left\{ \left( \triangleleft x \longrightarrow a, \begin{array}{c} a \longrightarrow \boxed{m} \longrightarrow b \end{array} \right) \mid a \in \mathbb{A} \right\} / \sim$$

# Tensor Product



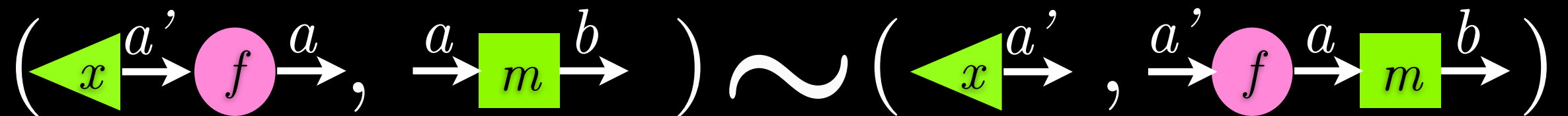
# Tensor Product



- cf. tensor product of vector spaces

$$cx \otimes y = x \otimes cy$$

# Tensor Product

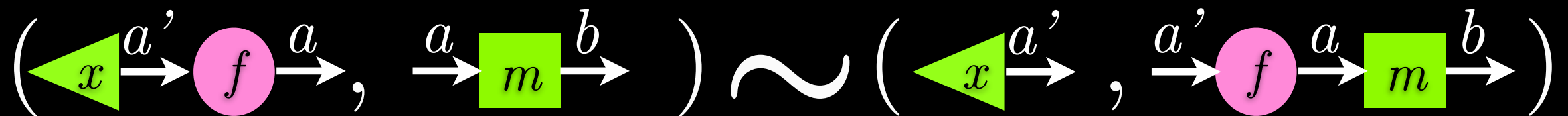


- cf. tensor product of vector spaces

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- def. by coend:  $(M \otimes P)b = \int^{a \in \mathbb{A}} P(a) \times M(a, b)$

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$$cx \otimes y = x \otimes cy$$

- def. by coend:  $(M \otimes P)b = \int^{a \in \mathbb{A}} P(a) \times M(a, b)$

- In the bicategory **Prof**/**Dist** :

$$\frac{\begin{array}{c} 1 \xrightarrow{P} \mathbb{A} \quad \mathbb{A} \xrightarrow{M} \mathbb{A} \\ \hline M \otimes P : 1 \xrightarrow{P} \mathbb{A} \xrightarrow{M} \mathbb{A} \end{array}}$$

# Tensor Product

$$\begin{array}{ccc}
 \{ \bullet \} & \xrightleftharpoons[1]{0} & \{ \bullet \dashv, \dashv \bullet, \dashv \bullet \} \\
 \uparrow \uparrow & & \text{inf} \uparrow \uparrow \text{sup} \\
 \emptyset & \xrightleftharpoons{\quad} & \{ \dashv \dashv, \dashv \dashv \}
 \end{array}
 \qquad
 P_{\mathbb{I}} : \{ * \} \rightrightarrows \mathbb{I}$$

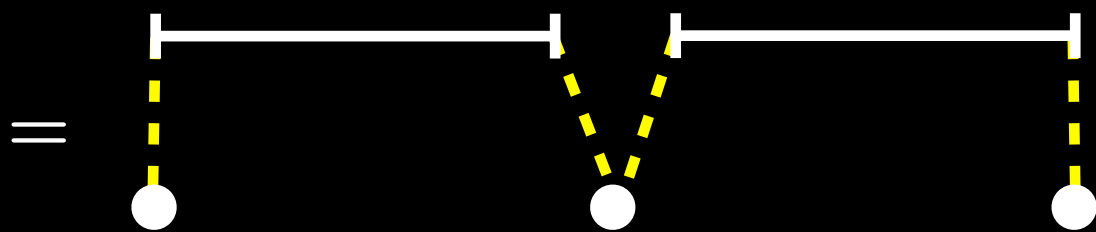
# Tensor Product

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 \end{array}
 \qquad
 P_{\mathbb{I}} : \{ * \} \rightrightarrows \mathbb{I}$$

$$(M_{\mathbb{I}} \otimes P_{\mathbb{I}})_1$$

$$= \left( P_{\mathbb{I}}(\mathbf{0}) \times M_{\mathbb{I}}(\mathbf{0}, \mathbf{1}) + P_{\mathbb{I}}(\mathbf{1}) \times M_{\mathbb{I}}(\mathbf{1}, \mathbf{1}) \right) / \sim$$

$$= \left( \{ * \} \times \{ \bullet \dashv, \dashv \bullet, \dashv \bullet \} + \mathbb{I} \times \{ \dashv \dashv, \dashv \dashv \} \right) / \sim$$



$$\cong \mathbb{I}$$



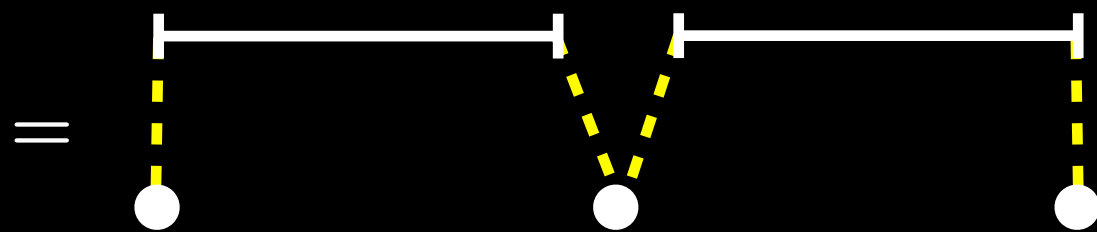
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$$= \left( \{ * \} \times \{ \bullet \dashv, \dashv \bullet, \dashv \bullet \} + \mathbb{I} \times \{ \dashv \dashv, \dashv \dashv \} \right) / \sim$$



$$\therefore M_{\mathbb{I}} \otimes P_{\mathbb{I}} \cong P_{\mathbb{I}}$$

$$\cong \mathbb{I}$$

# The Scenario

1. combinatorial spec.

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with gluing

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1. combinatorial spec.

$$\mathbf{M} \otimes (\_) : \mathbf{Sets}^\Delta \rightarrow \mathbf{Sets}^\Delta$$

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representatives

$$\begin{array}{c} \mathbf{M} \otimes \mathbf{Z} \\ \uparrow \text{final} \cong \\ \mathbf{Z} \end{array}$$

3. IFS

4.



with gluing

# The Scenario

1. combinatorial spec.

$$\mathbf{n} \cdot (\_) : \mathbf{Sets} \rightarrow \mathbf{Sets}$$

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$$\begin{array}{c} \mathbf{n} \cdot \mathbf{n}^\omega \\ \uparrow \text{final} \\ \mathbf{n}^\omega \end{array} \cong$$



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3. IFS

$$\begin{array}{c} \mathbf{n} \cdot \mathbf{X} \\ \downarrow \chi \\ \mathbf{X} \end{array}$$

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4.

# Non-Degeneracy: “Only Forced Equalities”

$$\begin{array}{c}
 M_{\mathbb{I}} \otimes P_{\text{deg}} \\
 \uparrow \cong \\
 \text{final} \\
 P_{\text{deg}}
 \end{array}$$

$P_{\text{deg}} :$	$\mathbb{A}_{\mathbb{I}}$	$\longrightarrow$	Sets
	$0$		$\{*\}$
	$l \downarrow \downarrow r$	$\dashrightarrow$	$!\downarrow\downarrow!$
	$1$		$\{*\}$

# Non-Degeneracy: “Only Forced Equalities”

$$\begin{array}{c}
 M_{\mathbb{I}} \otimes P_{\text{deg}} \\
 \uparrow \text{final} \cong \\
 P_{\text{deg}}
 \end{array}$$

$$\begin{array}{ccc}
 P_{\text{deg}} : & \mathbb{A}_{\mathbb{I}} & \longrightarrow \text{Sets} \\
 & \mathbf{0} & \{*\} \\
 & l \downarrow \downarrow r & \longmapsto \{!\downarrow!\} \\
 & \mathbf{1} & \{*\}
 \end{array}$$

- **Def.**  $P : \mathbb{A} \rightarrow \text{Sets}$  is *non-degenerate* if, in  $\text{el}(P)$

(ND1)

$$\begin{array}{ccc}
 (a, x) & (a', x') & \Longrightarrow \\
 \downarrow f & \downarrow f' & \\
 (b, y) & & \\
 \uparrow f & \uparrow f' & \\
 (a, x) & (a', x') & \\
 \uparrow \exists g & \uparrow \exists g' & \\
 \exists (c, z) & &
 \end{array}$$

(ND2)

$$\begin{array}{ccc}
 (a, x) & \Longrightarrow & \exists (c, z) \\
 \downarrow f' \downarrow f & & \downarrow \exists g \\
 (b, y) & & (a, x) \\
 & & \downarrow f' \downarrow f \\
 & & (b, y)
 \end{array}$$

# Non-Degeneracy: “Only Forced Equalities”

• **Prop.**  $P : \mathbb{A}_{\mathbb{I}} \longrightarrow \text{Sets}$

$$\begin{array}{ccc}
 \mathbf{0} & & P(\mathbf{0}) \\
 l \downarrow \downarrow r & \longmapsto & P(l) \downarrow \downarrow P(r) \\
 \mathbf{1} & & P(\mathbf{1})
 \end{array}$$

is non-degenerate iff

- $P(l)$  and  $P(r)$  are injective
- their images are disjoint



# Final *Non-Degenerate* Coalgebra

- **Prop.** [Freyd]

$$\begin{array}{c} M_{\mathbb{I}} \otimes P_{\mathbb{I}} \\ \uparrow \cong \\ P_{\mathbb{I}} \end{array}$$

is a final *non-degenerate* coalgebra

of symbolic nature

# Final *Non-Degenerate* Coalgebra

- **Prop.** [Freyd]

$$\begin{array}{c}
 M_{\mathbb{I}} \otimes P_{\mathbb{I}} \\
 \uparrow \cong \\
 P_{\mathbb{I}}
 \end{array}$$

is a final *non-degenerate* coalgebra

- **Thm.** (Leinster's construction)

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 \end{array}$$

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of symbolic nature

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1. combinatorial spec.

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with gluing

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# IIFS = Injective IFS

$$\varphi_0, \varphi_1 : \mathbb{C} \longrightarrow \mathbb{C} ,$$

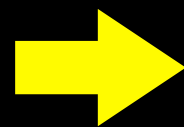
$$\varphi_0(x) = \frac{x}{2} ,$$

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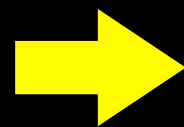


$$\begin{array}{c} 2 \cdot \mathbb{I} \\ \downarrow \chi \\ \mathbb{I} \end{array}$$

not injective!

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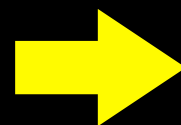
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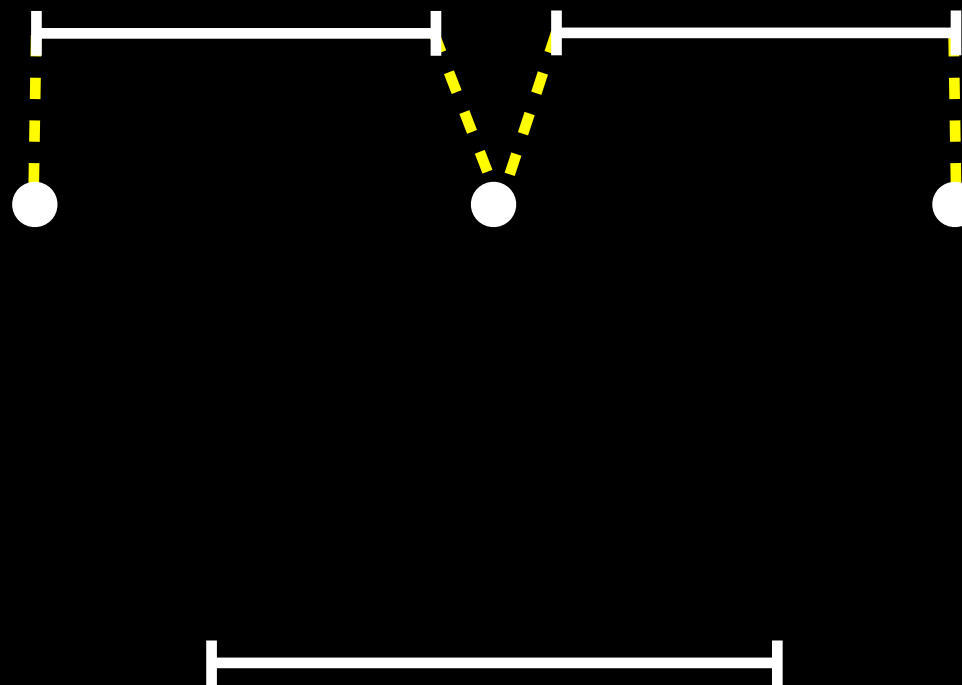
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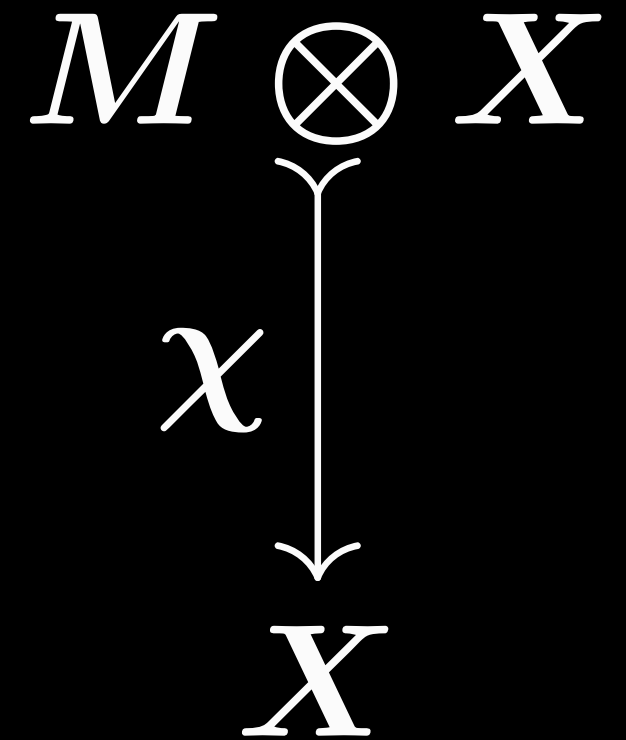


# Injective IFS ?

$$\begin{array}{ccc} M & \otimes & X \\ & \downarrow \chi & \\ & X & \end{array}$$

I. IFS + explicit gluing structure

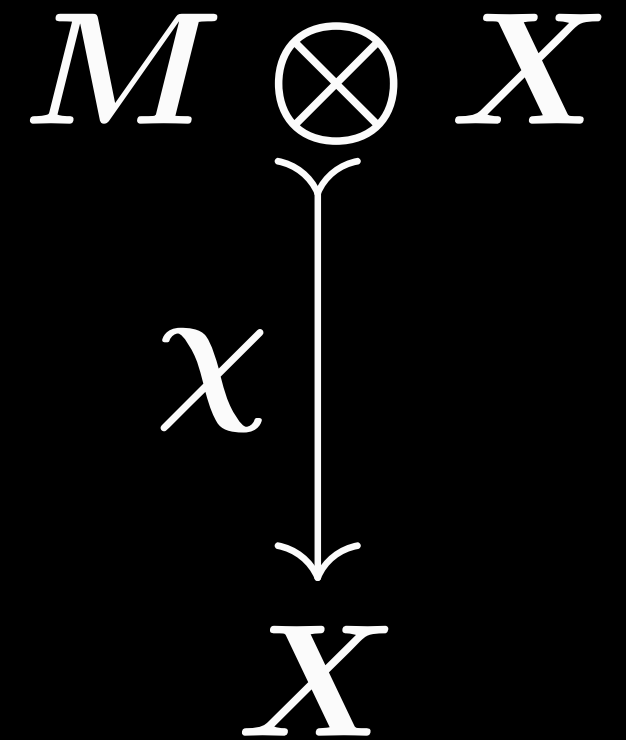
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- currently: IFS  IIFS

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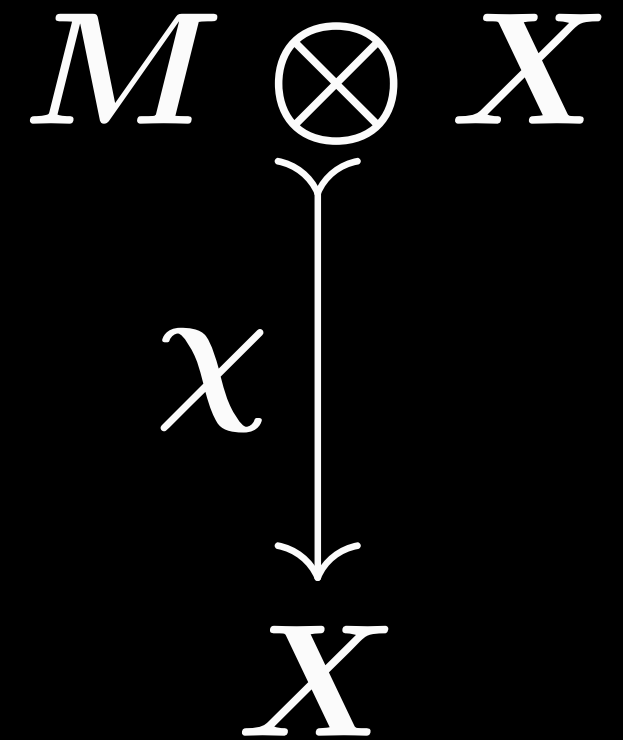
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+ “how to *metrically realize* it in a CMS  $X$ ”

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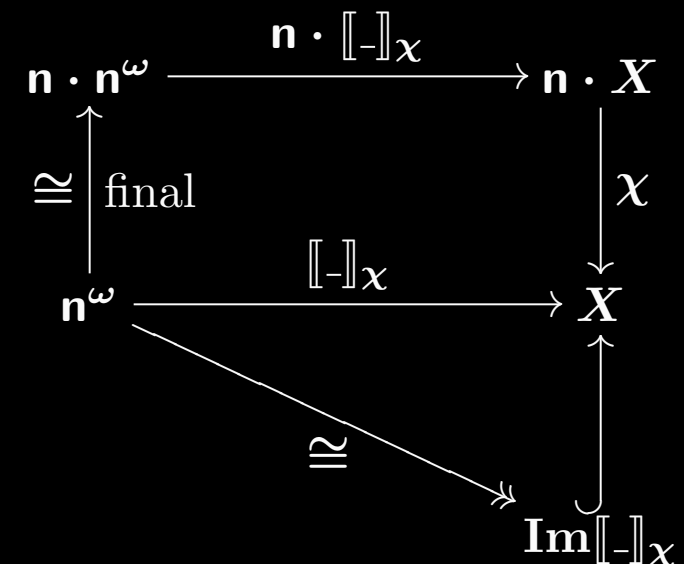
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with gluing

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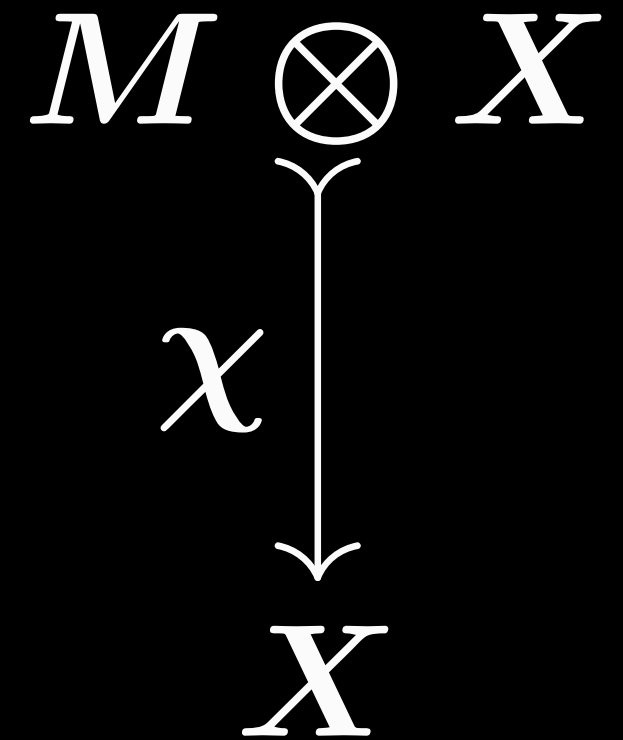
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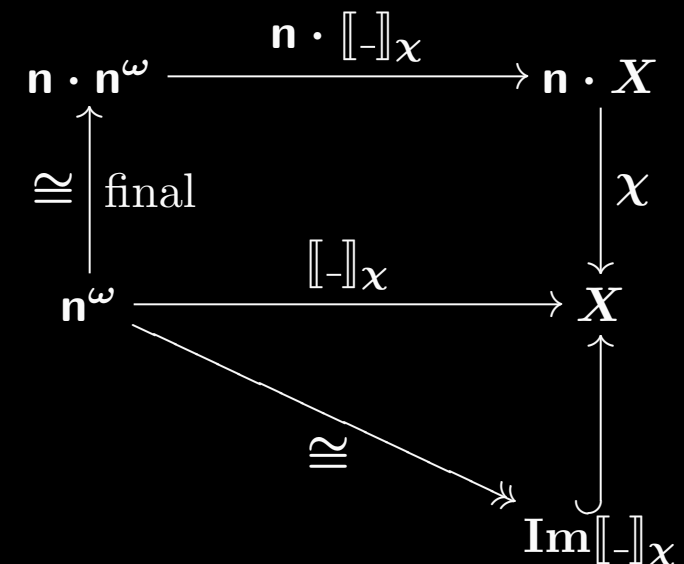
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# Injective IFS

- **Def.** An *injective IFS* over  $(\mathbb{A}, M)$  and  $\delta \in [0, 1)$  is

- $X : \mathbb{A} \longrightarrow \mathbf{CMet}_1^{\text{TB}}$ , non-degenerate

- $$\begin{array}{c} M \otimes \delta X \\ \chi \downarrow \\ X \end{array}$$

subject to

1.  $\chi_a$  is injective,  $\forall a \in \mathbb{A}$
2.  $Xa \neq \emptyset$  and  $Ia \neq \emptyset$ ,  $\forall a \in \mathbb{A}$
3. for  $f : b \rightarrow a$  in  $\mathbb{A}$ ,  
 $(Xf)^{-1}(\text{Im } \chi_a) \subseteq \text{Im } \chi_b$

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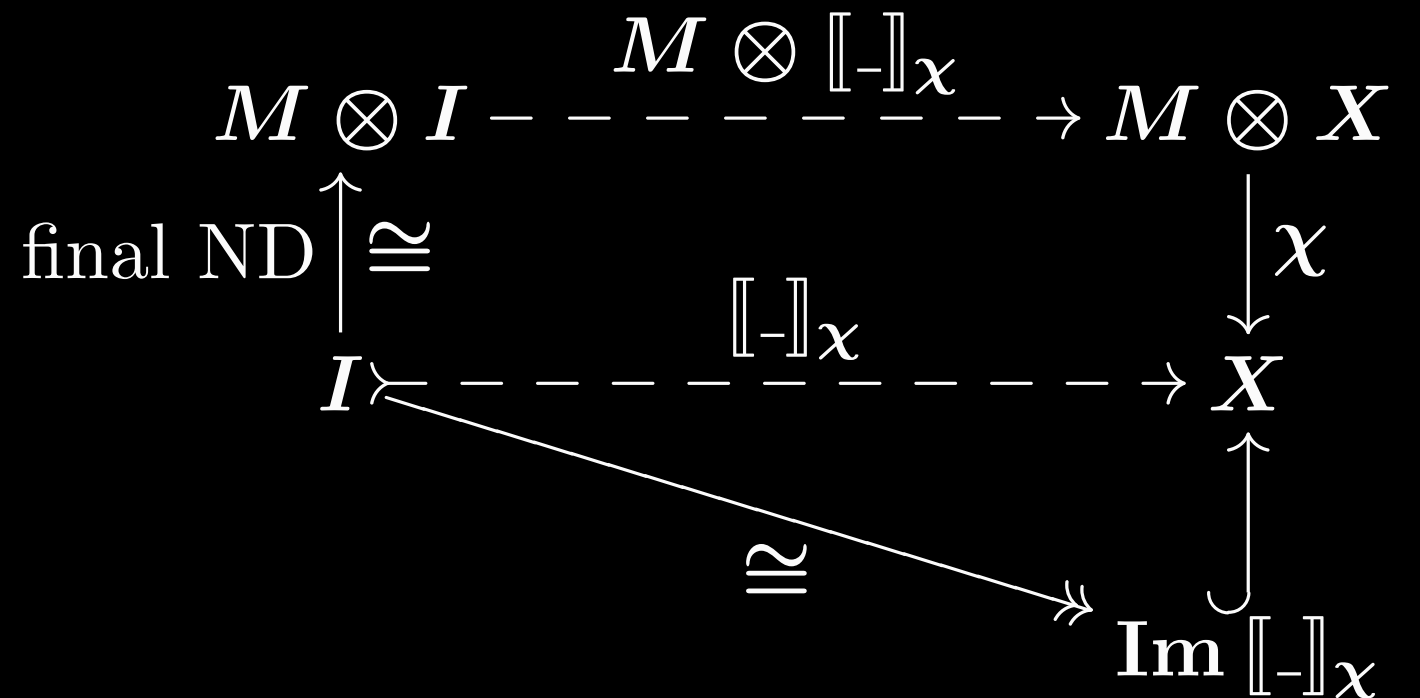
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# Main Result I

- **Thm.**  $M \otimes \delta X \xrightarrow{\chi} X$  : IIFS

1.  $\exists! \llbracket - \rrbracket_\chi$
2.  $\llbracket - \rrbracket_\chi$  is monic



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- Def.**  $M \otimes \delta X$   
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 such that

- $\varepsilon_a$  is injective,  $\forall a \in \mathbb{A}$
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- $\exists \sigma$ , an *isomorphism*, s.t.

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- Thm.**  $M \otimes I$   
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 $I$  is the unique attractor.

# Related Work

- Tom Leinster.  
A general theory of self-similarity I, II.

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- canonical topology over  $I$   
cf. [Barr, Adamek]
- *recognition theorem*: every compact metrizable space is a “fractal”

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metric

- injective IFS
- symbolic representation

current work

# Conclusion

fractal structure  symbolic representation of the shape, by infinite streams

- (co)algebra, IA-FC coincidence (no CPO)
- bijective representation
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**Thank you for your attention!**  
Ichiro Hasuo (RIMS, Kyoto U.)  
<http://www.kurims.kyoto-u.ac.jp/~ichiro>