

# 情報論理・プログラムの数理

## 第5回演習レポート（5月18日出題）解答

1.

For  $\Sigma$ -terms  $\mathbf{t}, \mathbf{s}_1, \dots, \mathbf{s}_n$  and mutually different variables  $\mathbf{x}_1, \dots, \mathbf{x}_n \in \mathbf{Var}$ , the *simultaneous substitution* of  $\mathbf{s}_1, \dots, \mathbf{s}_n$  for  $\mathbf{x}_1, \dots, \mathbf{x}_n$  in  $\mathbf{t}$  is denoted by  $\mathbf{t}[\mathbf{s}_1, \dots, \mathbf{s}_n/\mathbf{x}_1, \dots, \mathbf{x}_n]$ . It is defined as the  $\Sigma$ -term obtained by replacing free occurrences of  $\mathbf{x}_1, \dots, \mathbf{x}_n$  in  $\mathbf{t}$  with  $\mathbf{s}_1, \dots, \mathbf{s}_n$  respectively. Note here that free occurrences of  $\mathbf{x}_1, \dots, \mathbf{x}_n$  in  $\mathbf{s}_i$  are not replaced.

Then by induction on the construction of the term  $\mathbf{t}$ , we can prove the following proposition that is a generalization of Proposition 2.5.7.

**Proposition 1.**  $\llbracket \mathbf{t}[\mathbf{s}_1, \dots, \mathbf{s}_n/\mathbf{x}_1, \dots, \mathbf{x}_n] \rrbracket_{\mathbb{X}, J} = \llbracket \mathbf{t} \rrbracket_{\mathbb{X}, J[\mathbf{x}_1 \mapsto \llbracket \mathbf{s}_1 \rrbracket_{\mathbb{X}, J}, \dots, \mathbf{x}_n \mapsto \llbracket \mathbf{s}_n \rrbracket_{\mathbb{X}, J}]} \cdot \quad \square$

*Proof for Sublemma 2.6.9.* By Definition 2.5.11, in order to prove that  $\mathbb{X}$  is a  $(\Sigma, E)$ -algebra, it suffices to show that for all  $(\mathbf{s} = \mathbf{t}) \in E$ , we have  $\mathbb{X} \models \mathbf{s} = \mathbf{t}$ , i.e.  $\llbracket \mathbf{s} \rrbracket_{\mathbb{X}, J} = \llbracket \mathbf{t} \rrbracket_{\mathbb{X}, J}$  for each valuation  $J : \mathbf{Var} \rightarrow X$  on  $\mathbb{X}$ .

Let  $(\mathbf{s} = \mathbf{t}) \in E$  and  $J : \mathbf{Var} \rightarrow X$  be a valuation on  $\mathbb{X}$ . Moreover, as the numbers of variables in  $\mathbf{s}$  and  $\mathbf{t}$  are finite, let  $\text{FV}(\mathbf{s}) \cup \text{FV}(\mathbf{t}) = \{\mathbf{x}_1, \dots, \mathbf{x}_n\}$ . For each  $i \in \{1, \dots, n\}$ , let  $J(\mathbf{x}_i) = [\mathbf{u}_i]_{\sim_E}$ .

Then we have:

$$\begin{aligned}
 \llbracket \mathbf{s} \rrbracket_{\mathbb{X}, J} &= \llbracket \mathbf{s} \rrbracket_{\mathbb{X}, J_c[\mathbf{x}_1 \mapsto J(\mathbf{x}_1), \dots, \mathbf{x}_n \mapsto J(\mathbf{x}_n)]} && \text{(by Lemma 2.5.5)} \\
 &= \llbracket \mathbf{s} \rrbracket_{\mathbb{X}, J_c[\mathbf{x}_1 \mapsto [\mathbf{u}_1]_{\sim_E}, \dots, \mathbf{x}_n \mapsto [\mathbf{u}_n]_{\sim_E}]} && \text{(by definition)} \\
 &= \llbracket \mathbf{s} \rrbracket_{\mathbb{X}, J_c[\mathbf{x}_1 \mapsto \llbracket \mathbf{u}_1 \rrbracket_{\mathbb{X}, J_c}, \dots, \mathbf{x}_n \mapsto \llbracket \mathbf{u}_n \rrbracket_{\mathbb{X}, J_c}]} && \text{(by Sublemma 2.6.8)} \\
 &= \llbracket \mathbf{s}[\mathbf{u}_1, \dots, \mathbf{u}_n/\mathbf{x}_1, \dots, \mathbf{x}_n] \rrbracket_{\mathbb{X}, J_c} && \text{(by Proposition 1 above)} \\
 &= \llbracket \mathbf{s}[\mathbf{u}_1, \dots, \mathbf{u}_n/\mathbf{x}_1, \dots, \mathbf{x}_n] \rrbracket_{\sim_E} && \text{(by Sublemma 2.6.8)}
 \end{aligned}$$

Here  $J_c : \mathbf{Var} \rightarrow X$  denotes the canonical valuation introduced in the proof of Sublemma 2.6.7. Similarly, we have:

$$\llbracket \mathbf{t} \rrbracket_{\mathbb{X}, J} = \llbracket \mathbf{t}[\mathbf{u}_1, \dots, \mathbf{u}_n/\mathbf{x}_1, \dots, \mathbf{x}_n] \rrbracket_{\sim_E}$$

Let  $\mathbf{y}_1, \dots, \mathbf{y}_n \in \mathbf{Var}$  be mutually different variables such that  $\mathbf{y}_1, \dots, \mathbf{y}_n \notin \text{FV}(\mathbf{s}) \cup \text{FV}(\mathbf{t}) \cup \text{FV}(\mathbf{u}_1) \cup \dots \cup \text{FV}(\mathbf{u}_n)$  (as  $\mathbf{Var}$  is an infinite set, such a family indeed exists). Then we have the following proof tree.

$$\begin{array}{c}
\frac{}{\mathbf{s} = \mathbf{t}} \text{ (AXIOM)}, (\mathbf{s} = \mathbf{t}) \in E \\
\frac{}{\mathbf{s}[\mathbf{y}_1/\mathbf{x}_1] = \mathbf{t}[\mathbf{y}_1/\mathbf{x}_1]} \text{ (SUBST)} \\
\frac{}{\mathbf{s}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n] = \mathbf{t}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n]} \text{ (SUBST)} \\
\vdots \\
\frac{}{\mathbf{s}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n][\mathbf{u}_1/\mathbf{y}_1] = \mathbf{t}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n][\mathbf{u}_1/\mathbf{y}_1]} \text{ (SUBST)} \\
\vdots \\
\frac{}{\mathbf{s}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n][\mathbf{u}_1/\mathbf{y}_1] \dots [\mathbf{u}_n/\mathbf{y}_n] = \mathbf{t}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n][\mathbf{u}_1/\mathbf{y}_1] \dots [\mathbf{u}_n/\mathbf{y}_n]} \text{ (SUBST)}
\end{array}$$

Note here that

$$\mathbf{s}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n][\mathbf{u}_1/\mathbf{y}_1] \dots [\mathbf{u}_n/\mathbf{y}_n] \equiv \mathbf{s}[\mathbf{u}_1, \dots, \mathbf{u}_n/\mathbf{x}_1, \dots, \mathbf{x}_n]$$

and

$$\mathbf{t}[\mathbf{y}_1/\mathbf{x}_1] \dots [\mathbf{y}_n/\mathbf{x}_n][\mathbf{u}_1/\mathbf{y}_1] \dots [\mathbf{u}_n/\mathbf{y}_n] \equiv \mathbf{t}[\mathbf{u}_1, \dots, \mathbf{u}_n/\mathbf{x}_1, \dots, \mathbf{x}_n].$$

Hence by the definition of  $\sim_E$  (equation (2.18)), we have

$$[\mathbf{s}[\mathbf{u}_1, \dots, \mathbf{u}_n/\mathbf{x}_1, \dots, \mathbf{x}_n]]_{\sim_E} = [\mathbf{t}[\mathbf{u}_1, \dots, \mathbf{u}_n/\mathbf{x}_1, \dots, \mathbf{x}_n]]_{\sim_E}.$$

Hence we have  $\llbracket \mathbf{s} \rrbracket_{\mathbf{x}, J} = \llbracket \mathbf{t} \rrbracket_{\mathbf{x}, J}$  and this concludes the proof.  $\square$

## 2 (Exercise 3.2).

### 3.

1.

$$\frac{\frac{}{A \Rightarrow A} \text{ (INIT)}}{\neg A, A \Rightarrow} \text{ (\neg-L)} \\
\frac{}{A \Rightarrow \neg \neg A} \text{ (\neg-R)}$$

2.

$$\frac{\frac{}{A \Rightarrow A} \text{ (INIT)}}{\Rightarrow A, \neg A} \text{ (\neg-R)} \\
\frac{}{\neg \neg A \Rightarrow A} \text{ (\neg-L)}$$

$$\frac{\frac{}{A \Rightarrow A} \text{ (INIT)}}{B, A \Rightarrow A} \text{ (WEAKENING-L)} \\
\frac{}{A \Rightarrow B \supset A} \text{ (\supset-R)} \\
\frac{}{\Rightarrow A \supset (B \supset A)} \text{ (\supset-R)}$$

4.

$$\frac{\frac{\frac{}{B \Rightarrow B} \text{ (INIT)}}{A, B \Rightarrow B} \text{ (WEAKENING-L)}}{B \Rightarrow A \supset B} \text{ (\supset-R)} \quad \frac{\frac{}{B \Rightarrow B} \text{ (INIT)} \quad \frac{}{C \Rightarrow C} \text{ (INIT)}}{B \supset C, B \Rightarrow C} \text{ (\supset-L)} \\
\frac{}{(A \supset B) \supset (B \supset C), B, B \Rightarrow C} \text{ (EXCHANGE-L)} \\
\frac{}{B, (A \supset B) \supset (B \supset C), B \Rightarrow C} \text{ (EXCHANGE-L)} \\
\frac{}{B, B, (A \supset B) \supset (B \supset C) \Rightarrow C} \text{ (CONTRACTION-L)} \\
\frac{}{B, (A \supset B) \supset (B \supset C) \Rightarrow C} \text{ (\supset-R)} \\
\frac{}{(A \supset B) \supset (B \supset C) \Rightarrow B \supset C} \text{ (\supset-R)} \\
\frac{}{A, (A \supset B) \supset (B \supset C) \Rightarrow B \supset C} \text{ (WEAKENING-L)} \\
\frac{}{(A \supset B) \supset (B \supset C) \Rightarrow A \supset (B \supset C)} \text{ (\supset-R)}$$

5.

$$\begin{array}{c}
\frac{\overline{A \Rightarrow A} \text{ (INIT)}}{\frac{A \wedge B \Rightarrow A}{\neg A, A \wedge B \Rightarrow} (\wedge\text{-L1})} \quad \frac{\overline{B \Rightarrow B} \text{ (INIT)}}{\frac{A \wedge B \Rightarrow B}{\neg B, A \wedge B \Rightarrow} (\wedge\text{-L2})} \\
\frac{\frac{A \wedge B \Rightarrow A}{\neg A, A \wedge B \Rightarrow} (\neg\text{-L}) \quad \frac{A \wedge B \Rightarrow B}{\neg B, A \wedge B \Rightarrow} (\neg\text{-L})}{\frac{\neg A \vee \neg B, A \wedge B \Rightarrow}{A \wedge B \Rightarrow \neg(\neg A \vee \neg B)} (\vee\text{-L})} \\
\frac{A \wedge B \Rightarrow \neg(\neg A \vee \neg B)}{\Rightarrow (A \wedge B) \supset \neg(\neg A \vee \neg B)} (\supset\text{-R})
\end{array}$$

### 3 (Exercise 3.9).

An expression  $\models A$  denotes that the formula  $A$  is a tautology.

For  $\not\models A$ , we have:

$$\begin{aligned}
\not\models A &\Leftrightarrow A \text{ is not a tautology} && \text{(by definition)} \\
&\Leftrightarrow (\forall J : \mathbf{PVar} \rightarrow \{\text{tt}, \text{ff}\}. \llbracket A \rrbracket_J = \text{tt}) \text{ does not hold} && \text{(by Definition 3.3.3)} \\
&\Leftrightarrow \exists J : \mathbf{PVar} \rightarrow \{\text{tt}, \text{ff}\}. \llbracket A \rrbracket_J = \text{ff}.
\end{aligned}$$

In contrast, for  $\models \neg A$ , we have:

$$\begin{aligned}
\models \neg A &\Leftrightarrow \neg A \text{ is a tautology} && \text{(by definition)} \\
&\Leftrightarrow \forall J : \mathbf{PVar} \rightarrow \{\text{tt}, \text{ff}\}. \llbracket \neg A \rrbracket_J = \text{tt} && \text{(by Definition 3.3.3)} \\
&\Leftrightarrow \forall J : \mathbf{PVar} \rightarrow \{\text{tt}, \text{ff}\}. \llbracket A \rrbracket_J = \text{ff} && \text{(by Definition 3.3.2)}.
\end{aligned}$$

Therefore  $\not\models A$  and  $\models \neg A$  are different. Indeed, if  $A \equiv P \supset Q$  where  $P, Q \in \mathbf{PVar}$  and  $P \neq Q$ , then  $\not\models A$  holds as  $\llbracket A \rrbracket_{J[P \mapsto \text{tt}, Q \mapsto \text{ff}]} = \text{ff}$  while  $\models \neg A$  does not hold as  $\llbracket A \rrbracket_{J[P \mapsto \text{tt}, Q \mapsto \text{tt}]} = \text{tt}$ .

Moreover, by Definition 3.2.3,  $\not\models A$  iff there does not exist a proof tree whose root is  $\Rightarrow A$ . In contrast,  $\vdash \neg A$  iff there exists a proof tree whose root is  $\Rightarrow \neg A$ . Indeed, if  $A \equiv P \supset Q$  where  $P, Q \in \mathbf{PVar}$  and  $P \neq Q$ , then  $\not\models A$  holds while  $\vdash \neg A$  does not hold.