

A Compositional Approach to Parity Games

Kazuki Watanabe^{a,b,1,2} Clovis Eberhart^{b,c,1,3} Kazuyuki Asada^{d,1,4} Ichiro Hasuo^{a,b,e,1,5}

^a *The Graduate University for Advanced Studies (SOKENDAI), Hayama, Japan*

^b *National Institute of Informatics, Tokyo, Japan*

^c *Japanese-French Laboratory of Informatics, IRL 3527, CNRS, Tokyo, Japan*

^d *Tohoku University, Sendai, Japan*

^e *Ritsumeikan University, Kusatsu, Japan*

Abstract

In this paper, we introduce open parity games, which is a compositional approach to parity games. This is achieved by adding open ends to the usual notion of parity games. We introduce the category of open parity games, which is defined using standard definitions for graph games. We also define a graphical language for open parity games as a prop, which have recently been used in many applications as graphical languages. We introduce a suitable semantic category inspired by the work by Grellois and Melliès on the semantics of higher-order model checking. Computing the set of winning positions in open parity games yields a functor to the semantic category. Finally, by interpreting the graphical language in the semantic category, we show that this computation can be carried out compositionally.

Keywords: parity games, compositionality, graphical language, compact closed categories.

1 Introduction

Parity game is a major tool in theoretical computer science. Many formal verification problems such as model checking, satisfiability, etc.—can be reduced to solving parity games [39], where alternation of least and greatest fixed point operators in a specification is modeled by the parity winning condition. Efficient solutions of parity games, therefore, benefit many problems; recent algorithmic works include [9].

In this paper, we are interested in *compositionality* in formal verification in general, and in parity games in particular. It means that the property of a big system can be deduced from those of its constituent parts. One benefit is *efficiency*: compositionality can yield an efficient divide-and-conquer algorithm. Another is *maintainability*: compositional verification explicates an assumption that each subsystem must satisfy for the safety of the whole system; a subsystem can then be replaced freely as long as the local assumption is satisfied.

Compositional methods in model checking have been pursued in the literature, such as [8,31]. Many of those methods require a user to provide interfaces between subsystems, either as systems [8] or as specifications [31]. The role of compositionality is stressed in *higher-order model checking (HOMC)* [17,37], too, where intermediate results are combined along typing rules.

In this paper, influenced by the semantical constructs from [17], we introduce a categorical framework in which parity games are both presented and solved in a compositional manner. The presentation is by a

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² Email: kazukiwatanabe@nii.ac.jp

³ Email: eberhart@nii.ac.jp

⁴ Email: asada@riec.tohoku.ac.jp

⁵ Email: i.hasuo@acm.org

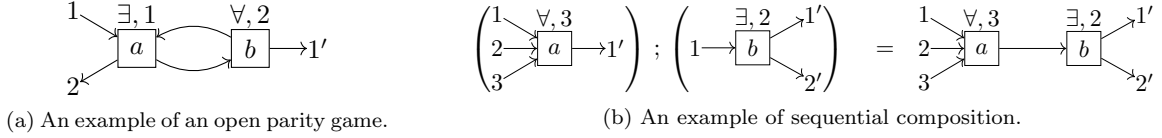


Fig. 1. Examples of open parity games.

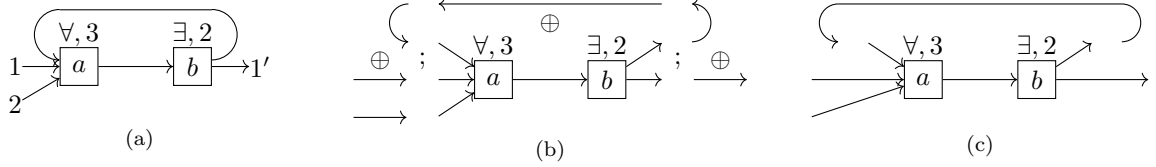


Fig. 2. An example of a cycle and its decomposition using the compact closed structure and the traced monoidal structure.

$$\begin{array}{ccc}
 (\text{Int}(\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}})) \simeq) \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) & \xrightarrow{\mathcal{R}_M} & \mathbf{OPG}_M (= \text{Int}(\mathbf{OPG}_M^r)) \\
 \llbracket - \rrbracket_M \searrow & & \swarrow \mathcal{W}_M \\
 & \text{Int}(\mathbf{FinScottL}_{!M}^{\text{op}}) &
 \end{array}$$

 Fig. 3. An outline. \mathcal{R}_M is the *realization functor* that maps a string diagram to an open parity game; \mathcal{W}_M is the *winning position functor* which extends the usual definition of winning positions in parity games; and $\llbracket - \rrbracket_M$ is the *interpretation functor*.

prop [33] (products and permutations category), a categorical notion of “monoidal” algebraic structure. This categorical presentation enables us to formulate compositionality as the preservation of suitable structures of certain functors. It also enables us to exploit general categorical structures (traced, compact closed, etc.) and properties (such as freeness). The use of props as graphical languages for various mathematical structures has been actively pursued recently (such as signal flow diagrams, matrices, and network games) [4, 5, 11]; the current work adds a new item to the list, namely parity games.

Contribution. The outline of our paper is Fig. 3. We extend parity games with so-called *open ends* so that we can compose them. The resulting notion (*open parity game*) is organized in a compact closed category denoted by \mathbf{OPG}_M . As a graphical language for open parity games, we use the prop $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ freely generated by a suitable monoidal (algebraic) theory $(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$. The other category $\text{Int}(\mathbf{FinScottL}_{!M}^{\text{op}})$ in Fig. 3 originates from [17]—it is our *semantic category* that tells which player is winning for (closed) parity games; for open parity games, it provides intermediate results of a suitable granularity to decide winners later.

Our main theorem (Thm. 5.10) is the commutativity of Fig. 3; it says that the semantics of parity games \mathcal{W}_M —defined as usual in terms of plays, strategies, and the parity acceptance condition—can be computed compositionally by a compact closed functor $\llbracket - \rrbracket_M$. The last compositional computation is illustrated in Ex. 5.11. After all, in the framework in Fig. 3, one writes down a parity game as a composition of smaller ones, in the graphical language of the prop $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$; when it comes to solving games, the winning positions for larger games are computed from those of the smaller ones, using that $\llbracket - \rrbracket_M$ preserves composition.

We illustrate our notion of *open parity games* that populates the category \mathbf{OPG}_M . Open parity games are parity games that come additionally with interfaces called *open ends*, along which they can be composed. An example is in Fig. 1(a). The domain interface consists of two *open ends*, 1 and 2, and the codomain interface simply of $1'$, while the internal positions are a and b , each equipped with a role and a priority, as usual in parity games. We give an example of sequential composition in Fig. 1(b). There, arrows are composed through the intermediate interfaces $1', 1$ between the two games. We also have a parallel composition \oplus , and Fig. 2(b) gives an example of how a cycle (in Fig. 2(a)) can be defined using sequential and parallel composition.

The technical key in Fig. 3 is the identification of *compact closed* structures. All the three categories are compact closed; moreover, we identify the prop $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ to be a *free* compact closed category in a suitable sense. The functors \mathcal{R}_M and $\llbracket - \rrbracket_M$ arise by the freeness; the commutativity is proved by the freeness, too.

In this paper, we find a new application of props as graphical languages in parity games. It allows one to solve parity games in a compositional manner (Ex. 5.11), thanks also to the identification of the right semantical domain (namely $\text{Int}(\mathbf{FinScottL}_{!M}^{\text{op}})$) that retains the right level of information in intermediate results. Such compositional solution has multiple potential applications. Firstly, the categorical structure we identify has a lot in common with those used for HOMC [17, 37]. Therefore we expect we can streamline known HOMC algorithms and reveal their categorical essences. Secondly, we will pursue algorithmic applications, such as

efficient divide-and-conquer algorithms and those which accommodate blackbox components as part of a game.

Organization. In §2, we introduce open parity games. In §3, we define the graphical language $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ and the realization functor \mathcal{R}_M . In §4, we define the semantic category $\text{Int}(\mathbf{FinScottL}_{!M}^{\text{op}})$ and the interpretation functor $\llbracket - \rrbracket_M$. In §5, we define the winning position functor \mathcal{W}_M and establish the triangle in Fig. 3. We also exhibit an example of compositional solution of a parity game in Ex. 5.11. We conclude in §6.

Related Work. We use $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ as the graphical language for open parity games. The use of monoidal categories as graphical languages dates back to [34]. There have been numerous such languages; see [36] for a survey. Languages that compositionally describe graph-like structures are of particular interest to us: [12] describes the algebra of directed acyclic graphs but does not consider cyclic structures; [2] describes open Petri nets, and compositionality is achieved “externally” by the use of cospans.

In particular, props have been used extensively as graphical languages. They define graphical languages as models for some mathematical structures (signal flow diagrams [5], networks [1], Petri nets [3], automata [35], and the ZX-calculus [6, 7] respectively) and prove that the graphical language is equivalent to the category that they are studying. They can therefore transfer properties of the graphical language (for example, decidability of equivalence of diagrams) to the category they are studying. In our work, however, we use the graphical language for expressing open parity games compositionally, and we are not necessarily interested in equivalence between $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ and \mathbf{OPG}_M (see also Rem. 3.16).

Note that our work uses a 2-colored prop for modeling the two possible directions for edges in an open parity game, while [3, 5, 11] only have a single type of edges, which are undirected. In [35], the authors use a colored prop to model different kinds of edges, and in particular, they use two colors to model directed edges.

Kissinger gives a general construction of the free traced symmetric monoidal categories $\mathbb{F}_{\text{tr}}(\Sigma)$, which are also props [29]. This is related to the free compact closed category $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ in the present paper, as explained in §3. Free traced monoidal categories are also given in [26] in the study of attribute grammars, where many-to-one signatures are treated while Kissinger’s paper treats many-to-many signatures.

Another related work is [15], which introduces the concept of composing games. Their approach is mainly applied to economic models, and they use a symmetric monoidal category for compositional game theories. However, their framework is different from ours in the sense that the objects along which games are composed have different meanings: in our framework, games are composed along graph edges, while in theirs, games are composed along interfaces describing player choices, game utility, etc.

Notation. We use the following notations. (i) $[m] := \{1, \dots, m\}$. (ii) We write the unit and the multiplication of a free monoid C^* as ε and \cdot . (iii) For $w \in \{\mathfrak{r}, \mathfrak{l}\}^*$, \bar{w} is the number of \mathfrak{r} in w , and \overleftarrow{w} is the number of \mathfrak{l} in w . (iv) $\mathbb{N}_M := \{0, \dots, M\}$ where $M \in \mathbb{N}$. (v) $\mathbb{N}_{\geq 1} := \{i \in \mathbb{N} \mid i \geq 1\}$. (vi) $\mathfrak{r}^\perp := \mathfrak{l}$, $\mathfrak{l}^\perp := \mathfrak{r}$, and $(w_1 \cdots w_n)^\perp := w_n^\perp \cdots w_1^\perp$ where $w_1 \cdots w_n \in \{\mathfrak{r}, \mathfrak{l}\}^*$. (vii) $!_X^\emptyset$ is the unique function from \emptyset to X . (viii) For $f : A \rightarrow C$ and $g : B \rightarrow C$, we write $[f, g] : A + B \rightarrow C$ for the copairing function, i.e., $[f, g](a) = f(a)$ and $[f, g](b) = g(b)$. (ix) We often omit the injections $\text{in}_i : A_i \rightarrow A_1 + A_2$, i.e., for example, we may write $a \in A + B$ for $\text{in}_1(a) \in A + B$, if no confusion happens.

2 Categories of Open Parity Games

We introduce the notion of open parity games. It extends parity games by adding *open ends* to the game, which are used to define composition of parity games; specifically, we obtain a compact closed category \mathbf{OPG}_M of open parity games (Def. 2.11).

In this paper, we often encounter situations where the structure of interest can be organized both as a *traced symmetric monoidal category (TSMC)* or as a *compact closed category (CpCC)*. (Specifically, we have three such classes of structures, yielding three TSMCs and CpCCs. See Fig. 3.) While our applicational interests lie in the CpCC structures, we work mostly with the TSMC structures for technical convenience, and use the Int construction [25] to define the CpCC structures from them (\mathbf{OPG}_M is defined as $\text{Int}(\mathbf{OPG}_M^{\text{tr}})$, and we show that $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ is equivalent to $\text{Int}(\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}}))$ for some signature Σ_M^{int}).

2.1 Open Parity Games

Recall that a *parity game* is a tuple $\mathcal{A} = (Q, E, \rho, \omega)$ where (Q, E) is a finite directed graph of *positions* and *edges*, $\rho : Q \rightarrow \{\exists, \forall\}$ is the *role function*, and $\omega : Q \rightarrow \mathbb{N}$ is the *priority function*. An *infinite play* on \mathcal{A} is an infinite sequence $q_0 q_1 \cdots \in Q^\mathbb{N}$ such that, for all $i \geq 0$, $(q_i, q_{i+1}) \in E$. A *finite play* is defined similarly. Let Play_\exists and Play_\forall be the sets of finite plays $q_0 \dots q_n$ such that $\rho(q_n) = \exists$ or $\rho(q_n) = \forall$, respectively. An infinite play $q_0 q_1 \dots$ is *winning* for \exists if the maximum priority appearing infinitely often in $\omega(q_0)\omega(q_1) \dots$ is even. A finite play $q_0 \dots q_n$ is winning for \exists if $\rho(q_n) = \forall$. A *strategy* of \exists is a partial function $\sigma_\exists : \text{Play}_\exists \rightarrow Q$ such that $(q_n, \sigma_\exists(q_0 \dots q_n)) \in E$ if $\sigma_\exists(q_0 \dots q_n)$ is defined. For any position q and strategies σ_\exists and σ_\forall , we denote by $\text{play}_q^{\sigma_\exists, \sigma_\forall}$ the unique play starting from q and consistent with both σ_\exists and σ_\forall . A strategy σ_\exists is *winning* for

\exists from $q \in Q$ if for all strategies σ_{\forall} , $\text{play}_q^{\sigma_{\exists}, \sigma_{\forall}}$ is winning for \exists . A position $q \in Q$ is *winning* for \exists if there is a strategy σ_{\exists} winning for \exists from q . We define *open parity games* by extending parity games with open ends.

Definition 2.1 (open parity game) An *open parity game* from \bar{m} to \bar{n} is a tuple $(\bar{m}, \bar{n}, Q, E, \rho, M, \omega)$ such that the following conditions are satisfied:

- (i) $\bar{m} = (m_r, m_l)$ and $\bar{n} = (n_r, n_l)$ are pairs of natural numbers, where \bar{m} represents the domain interface of the game and \bar{n} the codomain interface.
- (ii) Q is a finite set, whose elements are called *internal positions*.
- (iii) E is a relation $E \subseteq ([m_r + n_l] + Q) \times ([n_r + m_l] + Q)$, whose element is called an *edge*. Moreover, for any $s \in [m_r + n_l]$, there is a unique $s' \in [n_r + m_l] + Q$ such that $(s, s') \in E$; and similarly for any $t \in [n_r + m_l]$, there is a unique $t' \in [m_r + n_l] + Q$ such that $(t', t) \in E$.
- (iv) ρ is a function $\rho : Q \rightarrow \{\exists, \forall\}$, which assigns a *role* to each internal position.
- (v) $M \in \mathbb{N}$ is called the *maximal rank* and $\omega : Q \rightarrow \mathbb{N}_M$ is called the *priority function*.

We call an element of $([m_r + n_l] + [n_r + m_l]) + Q$ a *position*, one of $[m_r + n_l] + [n_r + m_l]$ an *open end*, one of $[m_r + n_l]$ an *entry position*, and one of $[n_r + m_l]$ an *exit position*.

We extend the priority function ω to $\omega : ([m_r + n_l] + [n_r + m_l]) + Q \rightarrow \mathbb{N}_M$ by $\omega(i) = 0$ for $i \in [m_r + n_l] + [n_r + m_l]$, i.e., we define the priority of each open end to be 0.

Example 2.2 The open parity game in Fig. 1(a) is the tuple $(\bar{m}, \bar{n}, Q, E, \rho, M, \omega)$ where

$$\begin{aligned} \bar{m} &= (1, 1), & \bar{n} &= (1, 0), & Q &= \{a, b\}, & \rho(a) &= \exists, & \rho(b) &= \forall, \\ 1 E a, & a E 2, & a E b, & b E a, & b E 1', & \omega(a) &= 1, & \omega(b) &= 2, & \text{and } M &= 2. \end{aligned}$$

In Fig. 1(a), the open end 1 is the entry position in $[m_r + n_l] = [1 + 0]$. The open ends 1' and 2 are the exit positions in $[n_r + m_l] = [1 + 1]$. The two boxes are internal positions in Q , with annotations on roles ρ and priorities ω . As usual, E is depicted by arrows.

Condition (iii) of Def. 2.1 requires that a unique outgoing/incoming edge from/to an entry/exit position, respectively. This condition can be enforced by adding some dummy positions.

The following definition is a first step towards introducing a trace operator.

Definition 2.3 (rightward open parity game) An open parity game $\mathcal{A} = (\bar{m}, \bar{n}, Q, E, \rho, M, \omega)$ is *rightward* if $\bar{m} = (m_r, 0_l)$ and $\bar{n} = (n_r, 0_l)$ for some m_r and n_r .

In the last definition, we require each open end in \bar{m} and \bar{n} to be headed in the right. Note that we do not impose the same requirement on (internal) edges in E —a rightward open parity game may contain cycles.

2.2 A Traced Symmetric Monoidal Category of Rightward Open Parity Games

We shall first define the traced symmetric monoidal category \mathbf{OPG}_M^r of rightward open parity games. It yields the compact closed category \mathbf{OPG}_M of open parity games by the Int construction (see Fig. 3).

In fact, we do so restricting the priorities to be below a certain natural number M , talking about \mathbf{OPG}_M^r and \mathbf{OPG}_M . The reason for doing so is discussed in Rem. 4.4.

In what follows, we assume that a given rightward open parity game \mathcal{A} is of the form $\mathcal{A} = ((m^A, 0), (n^A, 0), Q^A, E^A, \rho^A, M, \omega^A)$. The convention also applies to \mathcal{B} .

We need an equivalence relation on the set of rightward open parity games to define \mathbf{OPG}_M^r . For our purpose here, we define the equivalence in terms of structure-preserving bijections. It is easy to define an equivalence relation on open parity games in the same way.

Definition 2.4 (equivalence relation \sim on rightward open parity games) We define an equivalence relation \sim on the set of rightward open parity games as follows: $\mathcal{A} \sim \mathcal{B}$ if $m^A = m^B$, $n^A = n^B$, and there is a bijection $\eta : Q^A \rightarrow Q^B$ such that the following conditions are satisfied: (i) for $(s, t) \in ([m^A] + Q^A) \times ([n^A] + Q^A)$, $(s, t) \in E^A \iff (\bar{\eta}(s), \bar{\eta}(t)) \in E^B$, (ii) for $s \in Q^A$, $\rho^A(s) = \rho^B(\eta(s))$, and (iii) for $s \in Q^A$, $\omega^A(s) = \omega^B(\eta(s))$. Here we extend η to $\bar{\eta} : (\mathbb{N} + Q^A) \rightarrow (\mathbb{N} + Q^B)$ by $\bar{\eta}(n) = n$ for $n \in \mathbb{N}$.

We define the category \mathbf{OPG}_M^r as follows. Objects are natural numbers, and a morphism from m to n is an equivalence class $[\mathcal{A}]_{\sim}$ of rightward open parity games from $(m, 0)$ to $(n, 0)$. The identity and composition of morphisms are given by $\text{id}_n := [\mathcal{I}_n]_{\sim}$ and $[\mathcal{A}]_{\sim}; [\mathcal{B}]_{\sim} := [\mathcal{A}; \mathcal{B}]_{\sim}$, where \mathcal{I}_n and $\mathcal{A}; \mathcal{B}$ are given in Def. 2.5 and Def. 2.6 below, respectively.

Definition 2.5 (identity) For $n \in \mathbb{N}$, we define the *identity game* \mathcal{I}_n as $((n, 0), (n, 0), \emptyset, E, !_{\{\exists, \forall\}}^\emptyset, M, !_{\mathbb{N}_M}^\emptyset)$ where $E = \{(a, a) \mid a \in [n]\}$.

Fig. 4 shows the identity id_3 .

Next, we define the sequential composition $\mathcal{A}; \mathcal{B}$ of rightward open parity games. The intuition is to connect each exit position of \mathcal{A} with the corresponding entry position of \mathcal{B} , and then to hide those interface open ends. Fig. 1(b) in the introduction illustrates this construction.

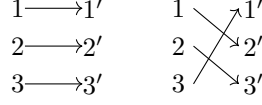


Fig. 4: id_3 .

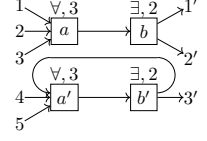


Fig. 5: $\sigma_{2,1}$. Fig. 6: Parallel composition of Fig. 1(b) & Fig. 2(a).

Definition 2.6 (sequential composition) Let \mathcal{A} and \mathcal{B} be rightward open parity games and $n^{\mathcal{A}} = m^{\mathcal{B}}$. We define the sequential composition $\mathcal{A}; \mathcal{B}$ as follows: $\mathcal{A}; \mathcal{B} = ((m^{\mathcal{A}}, 0), (n^{\mathcal{B}}, 0), Q^{\mathcal{A}} + Q^{\mathcal{B}}, E^{\mathcal{A}; \mathcal{B}}, [\rho^{\mathcal{A}}, \rho^{\mathcal{B}}], M, [\omega^{\mathcal{A}}, \omega^{\mathcal{B}}])$, where $E^{\mathcal{A}; \mathcal{B}} = E^{\mathcal{A}} \setminus (([m^{\mathcal{A}}] + Q^{\mathcal{A}}) \times [n^{\mathcal{A}}]) + E^{\mathcal{B}} \setminus ([m^{\mathcal{B}}] \times ([n^{\mathcal{B}}] + Q^{\mathcal{B}})) + \{(s, s') \mid \text{there exists } a \in [n^{\mathcal{A}}] = [m^{\mathcal{B}}] \text{ such that } (s, a) \in E^{\mathcal{A}} \text{ and } (a, s') \in E^{\mathcal{B}}\}$.

We can show associativity and unitality up to structure-preserving bijection, which entails that \mathbf{OPG}_M^r is a category by Def. 2.4.

We also define a parallel (or vertical) composition \oplus of rightward open parity games, which gives a monoidal product structure of \mathbf{OPG}_M^r by $[\mathcal{A}]_{\sim} \oplus [\mathcal{B}]_{\sim} = [\mathcal{A} \oplus \mathcal{B}]_{\sim}$. Fig. 6 gives an example, notice that the open ends in the second game need to be shifted, for which we need the following definition: for $l \in \mathbb{N}$ and $s \in [m] + Q$, let $s^{\downarrow l} \in [l + m] + Q$ be defined by $s^{\downarrow l} = l + s$ if $s \in [m]$, and $s^{\downarrow l} = s$ if $s \in Q$.

Definition 2.7 (parallel composition) Let \mathcal{A} and \mathcal{B} be rightward open parity games. The parallel composition $\mathcal{A} \oplus \mathcal{B}$ is defined as follows: $\mathcal{A} \oplus \mathcal{B} = ((m^{\mathcal{A}} + m^{\mathcal{B}}, 0), (n^{\mathcal{A}} + n^{\mathcal{B}}, 0), Q^{\mathcal{A}} + Q^{\mathcal{B}}, E^{\mathcal{A} \oplus \mathcal{B}}, [\rho^{\mathcal{A}}, \rho^{\mathcal{B}}], M, [\omega^{\mathcal{A}}, \omega^{\mathcal{B}}])$, where $E^{\mathcal{A} \oplus \mathcal{B}}$ is given by $E^{\mathcal{A} \oplus \mathcal{B}} = E^{\mathcal{A}} + \{(s^{\downarrow m^{\mathcal{A}}}, t^{\downarrow n^{\mathcal{A}}}) \mid (s, t) \in E^{\mathcal{B}}\}$.

The following game swaps the order of entry positions and that of exit positions. This makes \mathbf{OPG}_M^r a *symmetric* monoidal category. Fig. 5 shows the swap game $\sigma_{2,1}$.

Definition 2.8 (swap) For any $m, n \in \mathbb{N}_{\geq 1}$, we define the *swap game* $\sigma_{m,n}$ as follows: $\sigma_{m,n} = ((m+n, 0), (n+m, 0), \emptyset, E^{\sigma_{m,n}}, !_{\{\exists, \forall\}}^\emptyset, M, !_{\mathbb{N}_M}^\emptyset)$, where $E^{\sigma_{m,n}} = \{(a, n+a) \mid a \in [m]\} \cup \{(m+a, a) \mid a \in [n]\}$.

Cycles are essential to parity games: without them there would not be any infinite play. To introduce cycles in rightward open parity games, we use a trace operator on \mathbf{OPG}_M^r , as illustrated in Fig. 2(c).

Definition 2.9 (trace operator of \mathbf{OPG}_M^r) Let l, m , and n be objects in \mathbf{OPG}_M^r . We define the trace operator $\text{tr}_{l,m,n} : \mathbf{OPG}_M^r(l+m, l+n) \rightarrow \mathbf{OPG}_M^r(m, n)$ as follows. Let $\mathcal{A} \in \mathbf{OPG}_M^r(l+m, l+n)$, i.e., let $m^{\mathcal{A}} = l+m$ and $n^{\mathcal{A}} = l+n$. Then, $\text{tr}_{l,m,n}([\mathcal{A}]_{\sim}) := [\text{tr}_{l,m,n}(\mathcal{A})]_{\sim}$ where

$$\begin{aligned} \text{tr}_{l,m,n}(\mathcal{A}) &= ((m, 0), (n, 0), Q^{\mathcal{A}}, E^{\text{tr}_{l,m,n}(\mathcal{A})}, \rho^{\mathcal{A}}, M, \omega^{\mathcal{A}}), \text{ where} \\ E^{\text{tr}_{l,m,n}(\mathcal{A})} &= \{(s, s') \in ([m] + Q^{\mathcal{A}}) \times ([n] + Q^{\mathcal{A}}) \mid s^{\downarrow l} E^{\mathcal{A}} a_1 E^{\mathcal{A}} \dots E^{\mathcal{A}} a_k E^{\mathcal{A}} s'^{\downarrow l} \text{ for some } k \in \mathbb{N}, (a_i)_i \in [l]^k\}. \end{aligned}$$

Here is the main result of this section. With the given definitions, the proof is (lengthy but) routine work.

Theorem 2.10 (\mathbf{OPG}_M^r) *The data $(\mathbf{OPG}_M^r, \oplus, \emptyset, \sigma, \text{tr})$ defined so far constitutes a strict traced symmetric monoidal category, where \emptyset denotes the obvious empty game.* \square

2.3 A Compact Closed Category of Open Parity Games

To obtain the category \mathbf{OPG}_M of open parity games, we use the *Int construction* [25] (see also [21] for some correction). It is a *free* construction from a traced symmetric monoidal category \mathbb{C} to a compact closed category $\text{Int}(\mathbb{C})$. We briefly explain how it is defined here, but see Appendix A or [25] for more details.

Let \mathbf{CpCC} be the 2-category of (locally small) CpCCs, compact closed functors, and monoidal natural transformations. Note that its 2-cells automatically respect compact closed structures and are monoidal natural *isomorphisms* [23, Proposition 7.1]. Also, let \mathbf{TrSMC}_g be the 2-category of (locally small) TSMCs, traced symmetric strong monoidal functors, and monoidal natural *isomorphisms*.

Then the Int construction is a left biadjoint to the embedding $\mathbf{TrSMC}_g \rightarrow \mathbf{CpCC}$. Specifically, given a traced symmetric monoidal category $(\mathbb{C}, \otimes, \mathbf{I}, \sigma, \text{tr})$, the category $\text{Int}(\mathbb{C})$ is defined as follows: An object of

$\text{Int}(\mathbb{C})$ is a pair (X_+, X_-) of objects of \mathbb{C} . Then $\text{Int}(\mathbb{C})((X_+, X_-), (Y_+, Y_-)) := \mathbb{C}(X_+ \otimes Y_-, Y_+ \otimes X_-)$, and $\text{id}_{(X_+, X_-)} := \text{id}_{X_+ \otimes X_-}$. Notably, for $f \in \text{Int}(\mathbb{C})((X_+, X_-), (Y_+, Y_-))$ and $g \in \text{Int}(\mathbb{C})((Y_+, Y_-), (Z_+, Z_-))$, the composite of f and g is defined using the trace operator, namely by $\text{tr}_{Y_-, X_+ \otimes Z_-, Z_+ \otimes X_-}^{\mathbb{C}}((\sigma_{Z_+, Y_-} \otimes \text{id}_{X_-}) \circ (g \otimes \text{id}_{X_-}) \circ (\text{id}_{Y_+} \otimes \sigma_{X_-, Z_-}) \circ (f \otimes \text{id}_{Z_-}) \circ (\sigma_{Y_-, X_+} \otimes \text{id}_{Z_-}))$. See [25] for details, including diagrammatic illustration.

Definition 2.11 (\mathbf{OPG}_M) Let \mathbf{OPG}_M be the compact closed category $\text{Int}(\mathbf{OPG}_M^r)$ of open parity games.

The following proposition is trivial from the definition.

Proposition 2.12 A morphism of \mathbf{OPG}_M is an \sim -equivalence class of open parity games (in Def. 2.1). \square

The compact closed structure is the basis of our compositional approach to parity games. The structure serves as *game constructors*, and the compact closed functor $\mathcal{W}_M : \mathbf{OPG}_M \rightarrow \text{Int}(\mathbf{FinScottL}_{!M}^{op})$ (see Fig. 3) given in §5 computes compositionally if an entry position in a composed game wins.

3 A Graphical Language of Open Parity Games

In this section, we introduce the category $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ as a graphical language for open parity games. The category is a *prop*, a symmetric monoidal category version of the notion of Lawvere theory whose use has been actively pursued recently [1, 5, 7]. It gives to open parity games introduced in §2 a language of string diagrams generated by certain generators and equations. Moreover, we find that the category $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ is free in two senses: (i) as the prop induced by a theory $(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ for open parity games; and (ii) as a compact closed category $\text{Int}(\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}}))$ (see Fig. 3). The second freeness is exploited in the compositional definition of the interpretation functor $\llbracket - \rrbracket_M$ in Fig. 3.

3.1 The Graphical Language $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$

We define $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ as a colored prop constructed from a symmetric monoidal theory $(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$. For the detail of this prop construction, the reader can consult, e.g., [7].

Definition 3.1 (*C-prop, morphism, C-Prop*) Let C be a set (of colors). A *C-prop* is a small strict symmetric monoidal category where the monoid of all the objects is the free monoid C^* of C . A *C-prop morphism* between C -props is a strict symmetric monoidal functor that is the identity on objects. We write **C-Prop** for the category of C -props and C -prop morphisms.

In this paper, we consider $\{\mathfrak{r}, \mathfrak{l}\}$ -props. The colors \mathfrak{r} and \mathfrak{l} represent “rightward” and “leftward”, respectively. They intuitively correspond to $m_{\mathfrak{r}}$ and $m_{\mathfrak{l}}$ of $\bar{m} = (m_{\mathfrak{r}}, m_{\mathfrak{l}})$ in an open parity game.

We want to define $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ as a free $\{\mathfrak{r}, \mathfrak{l}\}$ -prop. A free C -prop is generated from a *C-symmetric monoidal theory* (*C-SMT* for short), i.e., a pair of a *C-signature* and a set of *C-equations*. Intuitively, given a *C-SMT*, morphisms of the corresponding free C -prop are terms built freely from the signature (as well as sequential and parallel composition) and quotiented by the equations.

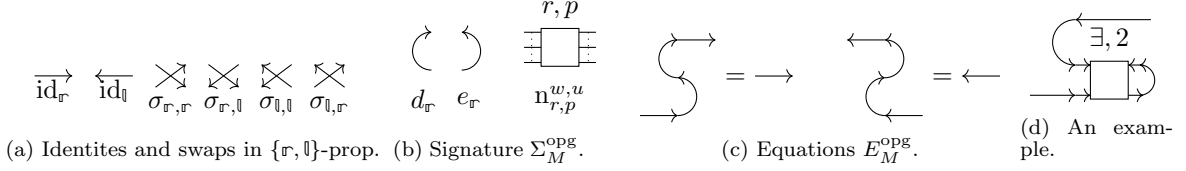
Definition 3.2 (*C-signature, morphism, C-Sig*) A *C-signature* is a functor $\Sigma : C^* \times C^* \rightarrow \mathbf{Set}$ where the free monoid C^* is thought of as a discrete category. A *C-signature morphism* $f : \Sigma \rightarrow \Sigma'$ is a natural transformation from Σ to Σ' . We write **C-Sig** for the category of C -signatures and C -signature morphisms.

Thus, $\mathbf{C-Sig} = \mathbf{Set}^{C^* \times C^*}$. We define an $\{\mathfrak{r}, \mathfrak{l}\}$ -signature Σ_M^{opg} , which is used to define the graphical language $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$. In the signature Σ_M^{opg} , for each domain $w \in \{\mathfrak{r}, \mathfrak{l}\}^*$, codomain $u \in \{\mathfrak{r}, \mathfrak{l}\}^*$, role $r \in \{\forall, \exists\}$, and priority $p \in \mathbb{N}_M$, there is a single generator $n_{r,p}^{w,u}$ that represents the type of nodes in open parity games with these specific domain, codomain, role, and priority.

Definition 3.3 ($\{\mathfrak{r}, \mathfrak{l}\}$ -signature Σ_M^{opg}) For $w, u \in \{\mathfrak{r}, \mathfrak{l}\}^*$, we define $\Sigma_M^{\text{opg}}(w, u)$ as follows: Let $N_{w,u} = \{n_{r,p}^{w,u} \mid r \in \{\forall, \exists\}, p \in \mathbb{N}_M\}$. Then $\Sigma_M^{\text{opg}}(\varepsilon, \mathfrak{r} \cdot \mathfrak{l}) = N_{\varepsilon, \mathfrak{r} \cdot \mathfrak{l}} \cup \{d_{\mathfrak{r}}\}$, $\Sigma_M^{\text{opg}}(\mathfrak{l} \cdot \mathfrak{r}, \varepsilon) = N_{\mathfrak{l} \cdot \mathfrak{r}, \varepsilon} \cup \{e_{\mathfrak{r}}\}$, and $\Sigma_M^{\text{opg}}(w, u) = N_{w,u}$ otherwise.

The generators $d_{\mathfrak{r}}$ and $e_{\mathfrak{r}}$ intuitively represent a unit and counit over \mathfrak{r} , respectively. We now turn to equations, for which we first need to define terms of a *C-SMT*. They are given by the following free construction. Let $\mathbb{U}_{\text{sig}}^C : \mathbf{C-Prop} \rightarrow \mathbf{C-Sig}$ be the obvious forgetful functor.

Theorem 3.4 ([13, 18]) The forgetful functor $\mathbb{U}_{\text{sig}}^C$ has a left adjoint $\mathbb{F}_{\text{sig}}^C : \mathbf{C-Sig} \rightarrow \mathbf{C-Prop}$. \square


 Fig. 7. Illustration of $(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$.

For the unit $\eta^{\text{sig}} : \text{Id}_{C\text{-sig}} \rightarrow \mathbb{U}_{\text{sig}}^C \circ \mathbb{F}_{\text{sig}}^C$, we identify $(\eta_{\Sigma}^{\text{sig}})_{w,u}(f) \in \mathbb{U}_{\text{sig}}^C(\mathbb{F}_{\text{sig}}^C(\Sigma))(w, u) = \mathbb{F}_{\text{sig}}^C(\Sigma)(w, u)$ with $f \in \Sigma(w, u)$ for simplicity of presentation.

Definition 3.5 (C -SMT) A C -colored symmetric monoidal theory (C -SMT for short) is a tuple (Σ, E, l, r) where Σ and E are C -signatures and $l, r : E \rightarrow \mathbb{U}_{\text{sig}}^C(\mathbb{F}_{\text{sig}}^C(\Sigma))$ are C -signature morphisms.

We often write simply (Σ, E) for (Σ, E, l, r) . We call $\mathbb{F}_{\text{sig}}^C(\Sigma)$ the set of *terms* generated by Σ , and E the set of (C -)equations in Σ , where each $e \in E$ represents the equation $l(e) = r(e)$.

Definition 3.6 (SMT $(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$) We complete the definition of the $\{\mathfrak{r}, \mathfrak{l}\}$ -SMT $(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}}, l_M^{\text{oppg}}, r_M^{\text{oppg}})$ by giving the equations: $(e_{\mathfrak{r}} \oplus \text{id}_{\mathfrak{l}}) \circ (\text{id}_{\mathfrak{l}} \oplus d_{\mathfrak{r}}) = \text{id}_{\mathfrak{l}}$ and $(\text{id}_{\mathfrak{r}} \oplus e_{\mathfrak{r}}) \circ (d_{\mathfrak{r}} \oplus \text{id}_{\mathfrak{r}}) = \text{id}_{\mathfrak{r}}$.

The $\{\mathfrak{r}, \mathfrak{l}\}$ -SMT $(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ describes open parity games, and the equations in Fig. 7(c) represent the coherence conditions of compact closed categories.

In general, a C -SMT induces a free C -prop $\mathbb{F}(\Sigma, E)$, whose arrows give a graphical language.

Definition 3.7 (free prop $\mathbb{F}(T)$ [1, 7]) Let $T = (\Sigma, E, l, r)$ be a C -SMT. We define a C -prop $\mathbb{F}(T)$ as the coequalizer of $l^{\dagger}, r^{\dagger} : \mathbb{F}_{\text{sig}}^C(E) \rightarrow \mathbb{F}_{\text{sig}}^C(\Sigma)$ in $C\text{-Prop}$ where l^{\dagger} and r^{\dagger} are, respectively, the transposition of $l, r : E \rightarrow \mathbb{U}_{\text{sig}}^C(\mathbb{F}_{\text{sig}}^C(\Sigma))$ in $C\text{-Sig}$ by $\mathbb{F}_{\text{sig}}^C \dashv \mathbb{U}_{\text{sig}}^C$.

Definition 3.8 (graphical language $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$) We define $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ by Def. 3.6 and Def. 3.7.

By definition, an object of $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ is an element of $\{\mathfrak{r}, \mathfrak{l}\}^*$, and a morphism is the class of terms generated by sequential and parallel composition applied to constants in Σ_M^{oppg} , and quotiented by the congruence produced by the equations in E_M^{oppg} (again under sequential and parallel composition). The prop $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ is illustrated in Fig. 7(a), Fig. 7(b), and Fig. 7(c).

Example 3.9 Fig. 7(d) is the morphism $(\text{id}_{\mathfrak{l}} \oplus e_{\mathfrak{r}}) \circ (\text{id}_{\mathfrak{l}} \oplus n_{\exists, 2}^{\mathfrak{r}, \mathfrak{l}, \mathfrak{r}}) \circ (d_{\mathfrak{l}} \oplus \text{id}_{\mathfrak{r}})$ from \mathfrak{r} to \mathfrak{l} in $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$.

3.2 Free Compact Closedness of $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ and the Full Functor \mathcal{R}_M

We show that the graphical language $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ is a free compact closed category, so that we can freely define a compact closed functor from $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ to any compact closed category \mathbb{C} with additional structure. This way we obtain the realization functor $\mathcal{R}_M : \mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}}) \rightarrow \mathbf{OPG}_M$ (see Fig. 3); we show that \mathcal{R}_M is full, meaning that every open parity game has a presentation in $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$. Due to the space limitation, we put the detailed information that is written here in Appendix B.

We need some definitions for proving that $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ is a free compact closed category. We call an object of the category $\mathbf{Set}^{\{\mathfrak{r}, \mathfrak{l}\}^* \times \{\mathfrak{r}, \mathfrak{l}\}^*}$ a *compact closed signature* (CCS, for short). Recall that the signature Σ_M^{oppg} consists of nodes $n_{r,p}^{w,u}$ of open parity games and the unit $d_{\mathfrak{r}}$ and counit $e_{\mathfrak{r}}$ of compact closed structure. We define a CCS Σ_M that is a signature of open parity games without the compact closed structure $d_{\mathfrak{r}}$ or $e_{\mathfrak{r}}$.

Definition 3.10 (CCS Σ_M) We define a CCS Σ_M by $\Sigma_M(w, u) := \{n_{r,p}^{w,u} \mid r \in \{\exists, \forall\} \text{ and } p \in \mathbb{N}_M\}$.

To state the free compact closedness of $\mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$, we define a *valuation*, which defines a way to interpret elements of a signature into a compact closed category.

Definition 3.11 (valuation) For a CCS Σ and a compact closed category \mathbb{C} , a *valuation* of Σ into \mathbb{C} is a pair $(V_{\mathfrak{r}}, (V_{w,u})_{w,u})$ such that (i) $V_{\mathfrak{r}} \in \text{ob}(\mathbb{C})$ and (ii) $V_{w,u} : \Sigma(w, u) \rightarrow \mathbb{C}(V_w^*, V_u^*)$ for $w, u \in \{\mathfrak{r}, \mathfrak{l}\}^*$ where V_w^* ($w \in \{\mathfrak{r}, \mathfrak{l}\}^*$) is defined as follows: $V_{d_1 \dots d_n}^* := V_{d_1} \otimes \dots \otimes V_{d_n}$ where $V_{\mathfrak{l}} := V_{\mathfrak{r}}^{\perp}$.

Definition 3.12 (action on valuations) Given a compact closed functor $F : \mathbb{C} \rightarrow \mathbb{D}$ and a valuation V of Σ into \mathbb{C} , the *action* $(F \circ V)$ on V by F is defined by (i) $(F \circ V)_{\mathfrak{r}} := F(V_{\mathfrak{r}})$ and (ii) $(F \circ V)_{w,u}(f) :=$

$(\phi_u^F)^{-1} \circ F(V_{w,u}(f)) \circ \phi_w^F$, where ϕ_w^F (and ϕ_u^F) are defined as follows: for any $w = r^{i_1} \cdots r^{i_n} \in \{\ulcorner, \llcorner\}^*$ where each r^{i_j} is either \ulcorner or \ulcorner^\perp , the morphism $\phi_w^F : F(V_r)^{i_1} \otimes \cdots \otimes F(V_r)^{i_n} \rightarrow F(V_r^{i_1} \otimes \cdots \otimes V_r^{i_n})$ is the isomorphism given by the fact that F respects the compact closed structures.

Finally, we prove the free compact closedness of $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ by using the above definitions.

Theorem 3.13 (free compact closedness of $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$) *The prop $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ is a strict compact closed category. Furthermore, $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ is a free compact closed category, i.e., there exists a valuation η_{Σ_M} of Σ_M into $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ such that, for any compact closed category \mathbb{C} and any valuation V of Σ_M into \mathbb{C} , there exists a unique (up to iso) compact closed functor $F : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \mathbb{C}$ such that $(F \circ \eta_{\Sigma_M}) = V$. \square*

For a 2-categorical and generalized version of Thm. 3.13, see Appendix B. By the general result (Thm. 3.13), we define the realization functor $\mathcal{R}_M : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \mathbf{OPG}_M$ (see Fig. 3):

Definition 3.14 (realization functor \mathcal{R}_M) We let the realization functor $\mathcal{R}_M : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \mathbf{OPG}_M$ be the functor determined by Thm. 3.13 with $\mathcal{R}_M(\ulcorner) = (1, 0)$ and $\mathcal{R}_M(\ulcorner_{r,p}^{w,u}) = ((\vec{w}, \overleftarrow{w}), (\vec{u}, \overleftarrow{u}), \{\star\}, E, \{\star \mapsto r\}, M, \{\star \mapsto p\})$ where $E = \{(a, \star) \mid a \in [\vec{w} + \overleftarrow{w}]\} \cup \{(\star, a) \mid a \in [\vec{u} + \overleftarrow{u}]\}$.

The following theorem says that every open parity game in \mathbf{OPG}_M can be represented as a graphical one, i.e., a morphism in $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$. The proof is given in a standard way and can be found in Appendix B.3.

Theorem 3.15 (fullness) *The functor $\mathcal{R}_M : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \mathbf{OPG}_M$ is full. \square*

Remark 3.16 The functor \mathcal{R}_M is not faithful (a one-node example can be easily given). In recent works on props such as [5, 35], the main interest is in the faithfulness of a semantics functor whose codomain is a well-known semantic category (that of linear relations [5], automata [35], etc.). In this case, faithfulness amounts to the *completeness* of equational axioms. We do not share this interest: in Fig. 3, the codomain of $\mathcal{R}_M : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \mathbf{OPG}_M$ is *not* a well-known category, and the value of a corresponding complete equational axiomatization is not clear. Faithfulness of the interpretation functor $\llbracket - \rrbracket_M : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \text{Int}(\mathbf{FinScottL}_{!_M}^{\text{op}})$ (introduced in §4) seems more interesting, since it amounts to an equational characterization of the equivalence of parity games in terms of who is winning. The problem seems challenging, however, given the complexity of solving parity games, and we leave it as future work. We note that, for our purpose of compositional solution of parity games (see e.g. Ex. 5.11), faithfulness of \mathcal{R}_M or $\llbracket - \rrbracket_M$ is not needed.

Kissinger gives a construction for free traced symmetric monoidal categories $\mathbb{F}_{\text{tr}}(\Sigma)$ [29]. We show an equivalence $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \simeq \text{Int}(\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}}))$ in \mathbf{CpCC} for some 1-signature Σ_M^{int} (see Appendix B.2).

4 The Semantic Category of Open Parity Games

In this section, we define a semantic category of open parity games $\text{Int}(\mathbf{FinScottL}_{!_M}^{\text{op}})$. Grellois and Melliès restricted \mathbf{ScottL} to the full subcategory $\mathbf{FinScottL}$ [16, 17] of *finite* preordered sets, in order to introduce a fixpoint operator on $\mathbf{FinScottL}_{!_M}$ for some suitable comonad $!_M$ so that it forms a model of higher-order model checking. We use this fixpoint operator for the Int construction of $\text{Int}(\mathbf{FinScottL}_{!_M}^{\text{op}})$. We then define the *interpretation functor* $\llbracket - \rrbracket_M : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \text{Int}(\mathbf{FinScottL}_{!_M}^{\text{op}})$. (The reader may look at Ex. 5.11 for an example of computation of $\llbracket \mathcal{A} \rrbracket_M$ on a concrete open parity game \mathcal{A} .)

Definition 4.1 ($\mathbf{FinScottL}$ [16, 17]) The category $\mathbf{FinScottL}$ has as objects finite preordered sets $A = (|A|, \leq_A)$ and as morphisms $R : A \multimap B$ downward-closed binary relations R between A^{op} and B : i.e., a binary relation $R \subseteq |A| \times |B|$ such that if $a' \geq_A a$, $a R b$, $b \geq_B b'$, then $a' R b'$. Composition is defined as usual: $a(S \circ R)c$ iff $a R b$ and $b S c$ for some $b \in |B|$. The identity on A is $\text{id}_A = \{(a', a) \mid a \leq_A a'\}$.

Note that $\mathbf{FinScottL}$ has a symmetric monoidal structure induced by the existence of finite cartesian products $1 = (\{\star\}, =)$ and $(|A|, \leq_A) \times (|B|, \leq_B) = (|A| + |B|, \leq_{A+B})$.

In open parity games, we have two roles: \exists and \forall . For player \exists , a choice by \forall is not predictable. This nondeterminism is represented by the finite powerset comonad \mathcal{P} .

Definition 4.2 (finite powerset comonad) The *finite powerset comonad* $(\mathcal{P}, \epsilon^{\mathcal{P}}, \delta^{\mathcal{P}})$ on $\mathbf{FinScottL}$ is defined by $\mathcal{P}(|A|, \leq_A) := (P(|A|), \leq)$, where $X \leq Y$ iff for any $x \in X$ there exists $y \in Y$ such that $x \leq_A y$. For $R : A \multimap B$, $\mathcal{P}(R) := \{(X, Y) \in P(|A|) \times P(|B|) \mid \forall y \in Y, \exists x \in X, (x, y) \in R\}$. Then $\epsilon_A^{\mathcal{P}} := \{(X, a) \in P(|A|) \times |A| \mid \exists x \in X, a \leq_A x\}$ and $\delta_A^{\mathcal{P}} := \{(X, \{Y_1, \dots, Y_n\}) \in P(|A|) \times P(P(|A|)) \mid Y_1 \cup \dots \cup Y_n \leq_{\mathcal{P}(A)} X\}$.

Priorities that are not greater than M in open parity games are represented by the coloring comonad \square_M .

Definition 4.3 (coloring comonad [16, 17]) The *coloring comonad* $(\square_M, \epsilon^{\square_M}, \delta^{\square_M})$ on $\mathbf{FinScottL}$ is defined as follows: $\square_M(|A|, \leq_A) = (\mathbb{N}_M \times |A|, \leq_{\square_M A})$ where $(p, a) \leq_{\square_M A} (q, b)$ iff $p = q$ and $a \leq_A b$. For $R : A \dashrightarrow B$, $\square_M(R) := \{((p, a), (p, b)) \in (\mathbb{N}_M \times |A|) \times (\mathbb{N}_M \times |B|) \mid p \in \mathbb{N}_M \text{ and } (a, b) \in R\}$. Then $\epsilon_A^{\square_M} := \{((0, a), a') \in (\mathbb{N}_M \times |A|) \times |A| \mid a' \leq_A a\}$ and $\delta_A^{\square_M} := \{((\max(p, q), a), (p, (q, a')))) \in (\mathbb{N}_M \times |A|) \times (\mathbb{N}_M \times (\mathbb{N}_M \times |A|)) \mid a' \leq_A a\}$.

Remark 4.4 We expect that this comonad can be extended to a *graded comonad* [14, 27] so that its Kleisli category interprets all open parity games, without fixing the parity bound M . However, we do not take this approach because it is reasonably harmless to fix the maximal parity, while the use of the complex notion of graded comonad makes it hard to see the essential idea.

Combining the above notions, we define the comonad $!_M$.

Definition 4.5 (comonad $!_M$ [16, 17]) We define a distributive law $\lambda : \mathcal{P} \circ \square_M \Rightarrow \square_M \circ \mathcal{P}$ on $\mathbf{FinScottL}$ by $\lambda_{(A, \leq_A)} := \{(X, (p, Y)) \in P(\mathbb{N}_M \times |A|) \times (\mathbb{N}_M \times P(|A|)) \mid \forall y \in Y, \exists a \in A, (p, a) \in X \text{ and } y \leq_A a\}$, and we define a comonad $!_M = (!_M, \epsilon^{!_M}, \delta^{!_M})$ on $\mathbf{FinScottL}$ by: (i) $!_M := \mathcal{P} \circ \square_M$, (ii) $\epsilon^{!_M} := \epsilon^{\mathcal{P}} \circ (\mathcal{P} * \epsilon^{\square_M})$, and (iii) $\delta^{!_M} := (\mathcal{P} * \lambda * \square_M) \circ (\delta^{\mathcal{P}} * \delta^{\square_M})$ where $*$ is the horizontal composition of natural transformations.

In general, the Kleisli category of a comonad inherits the cartesian product from the original category, and so the Kleisli category $\mathbf{FinScottL}_{!_M}$ has a cartesian products given by $(|A| + |B|, \leq_A + \leq_B)$. (Furthermore, $\mathbf{FinScottL}_{!_M}$ is cartesian closed [16, 17], though we do not use this fact.)

In order to give a model of higher-order model checking, Grellois and Melliès introduced a fixpoint operator $\text{fix}^{\mathbf{GM}}$ on $\mathbf{FinScottL}_{!_M}$ to deal with infinite plays [16, 17]. Its definition is based on the notion of *semantic run-tree*. From $\text{fix}^{\mathbf{GM}}$, we get a trace operator on $\mathbf{FinScottL}_{!_M}$, because having a fixpoint operator is equivalent to having a trace operator for a cartesian category [19]. We then get a trace operator $\text{tr}^{\mathbf{GM}}$ on $\mathbf{FinScottL}_{!_M}^{op}$, since if \mathbb{C} is a traced symmetric monoidal category, then so is \mathbb{C}^{op} canonically.

Now we give the definition of $\text{tr}^{\mathbf{GM}}$ (the definition of $\text{fix}^{\mathbf{GM}}$ can be found in Appendix C). In order to do this, we first adapt the notion of semantic run-tree by Grellois and Melliès (which we also call *semantic run-tree*) through the correspondance above between fixpoint operators and trace operators.

Definition 4.6 (semantic run-tree for $\text{tr}^{\mathbf{GM}}$) Let $R \in \mathbf{FinScottL}_{!_M}^{op}(D + A, D + B)$ and $a \in |A|$; then especially, $R \subseteq P(\mathbb{N}_M \times (|D| + |B|)) \times (|D| + |A|)$. A *semantic run-tree* ψ for R and a (for the trace operator) is a possibly infinite $(\mathbb{N}_M \times (|D| + |A| + |B|))$ -labeled tree ψ that satisfies the following conditions:

- (i) The label of the root of ψ is $(0, a) \in \mathbb{N}_M \times |A|$.
- (ii) Any node of ψ that is neither a leaf nor the root has its label in $\mathbb{N}_M \times |D|$.
- (iii) For any non-leaf node (possibly being the root that is not a leaf) of ψ with label $(p, x) \in \mathbb{N}_M \times (|D| + |A|)$, let $X \subseteq \mathbb{N}_M \times (|D| + |A| + |B|)$ be the set of the labels of all the children of the node. Then $(X, x) \in R$.
- (iv) For any leaf node (possibly being the root that is a leaf) of ψ such that its label belongs to $\mathbb{N}_M \times (|D| + |A|)$ (rather than $\mathbb{N}_M \times |B|$) and is (p, x) , we have $(\emptyset, x) \in R$.

We write $\text{SRT}(A, B, D, R, a)$ for the set of semantic run-trees with respect to A, B, D, R and a . For a semantic run-tree $\psi \in \text{SRT}(A, B, D, R, a)$, we define $\text{leaves}(\psi) \in |\mathcal{P}(\square_M(B))| = P(\mathbb{N}_M \times |B|)$ as the set of elements $(p, b) \in \mathbb{N}_M \times |B|$ such that there exists a leaf ℓ of ψ such that: (i) the label of leaf ℓ is (p', b) for some $p' \in \mathbb{N}_M$ and (ii) p is the maximal priority encountered on the path from the leaf ℓ to the root of ψ .

A semantic run-tree is similar to a (usual) run for a parity game, except that (i) its branching models \forall 's choices, and (ii) it is induced by a suitable semantic construct R instead of a graph-theoretic notion of game. In our use of the notion (§5), R will be a “summary” of an open parity game, which retains the necessary data to decide who is winning yet is much smaller than the original open parity game.

Definition 4.7 (trace operator $\text{tr}^{\mathbf{GM}}$ [16, 17]) For every $A, B, D \in \mathbf{FinScottL}_{!_M}^{op}$, we define a trace operator $\text{tr}_{D;A,B}^{\mathbf{GM}} : \mathbf{FinScottL}_{!_M}^{op}(D \otimes A, D \otimes B) \rightarrow \mathbf{FinScottL}_{!_M}^{op}(A, B)$ as follows:

$$\text{tr}_{D;A,B}^{\mathbf{GM}}(R) := \{(\text{leaves}(\psi), a) \mid \psi \in \text{SRT}(A, B, D, R, a) \text{ that meets the parity condition}\}$$

where a semantic run-tree meets the *parity condition* if for every infinite path $((p_i, x_i))_{i \in \mathbb{N}}$, the maximum priority met infinitely along the path is even (i.e., $\max\{q \mid \#\{i \mid p_i = q\} = \infty\}$ is even).

Thus $\mathbf{FinScottL}_M^{op}$ is a traced symmetric monoidal category. The trace operator above is used in the (sequential) composition of $\text{Int}(\mathbf{FinScottL}_M^{op})$ given below.

Now we define the semantic category $\text{Int}(\mathbf{FinScottL}_M^{op})$ for open parity games. In §5, we explain how $\mathbf{FinScottL}_M^{op}$ and $\text{Int}(\mathbf{FinScottL}_M^{op})$ serve as the semantic categories in the traced and compact closed structures, respectively, by giving a suitable *winning-position* functor \mathcal{W}_M^r from \mathbf{OPG}_M^r to $\mathbf{FinScottL}_M^{op}$ and then by inducing \mathcal{W}_M (see Fig. 3).

Definition 4.8 (semantic category $\text{Int}(\mathbf{FinScottL}_M^{op})$) By applying the Int construction to the traced symmetric monoidal category $\mathbf{FinScottL}_M^{op}$, we obtain the compact closed category $\text{Int}(\mathbf{FinScottL}_M^{op})$.

Remark 4.9 We have $\mathbf{FinScottL}_M^{op} \cong \mathbf{FinPreord}_T$, where $\mathbf{FinPreord}$ is the category of finite pre-ordered sets and monotonic functions, and (a Kleisli morphism of) the monad T is of the following form: $\mathbf{FinScottL}_M^{op}(A, B) \cong \mathbf{FinPreord}(A, (P^\uparrow(\mathcal{P}(\mathbb{N}_M \times |B|), \leq_{\square_{MB}}), \supseteq))$, where P^\uparrow is the upward-closed powerset. This description is closed to the double-powerset style semantics for 2-player games, e.g., in [22].

We want to define an interpretation functor $\llbracket - \rrbracket_M : \mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}}) \rightarrow \text{Int}(\mathbf{FinScottL}_M^{op})$ that reflects the winning condition on open parity games. The idea is that, if $((j_k, p_k)_{k \in n}, i) \in \llbracket G \rrbracket_M$, then player \exists can force any play that starts from the entry position corresponding to i in G to end in one of the exit positions corresponding to the j_k 's while encountering a maximum priority of p_k . By Thm. 3.13, we obtain this functor as:

Definition 4.10 (interpretation functor $\llbracket - \rrbracket_M$) We define the *interpretation functor* $\llbracket - \rrbracket_M : \mathbb{F}(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}}) \rightarrow \text{Int}(\mathbf{FinScottL}_M^{op})$ to be the compact closed functor whose action on objects is generated by $\llbracket [r] \rrbracket_M = (([1], =), (\emptyset, =)) \in \text{Int}(\mathbf{FinScottL}_M^{op})$ and whose action on morphisms is generated by:

$$\begin{aligned} \llbracket n_{\exists, p}^{a, b} \rrbracket_M &= \begin{cases} \emptyset & (\vec{b} + \overleftarrow{a} = 0) \\ \{(T, i) \in P(\mathbb{N}_M \times [\vec{b} + \overleftarrow{a}]) \times [\vec{a} + \overleftarrow{b}] \mid (j, p) \in T \text{ for some } j \in [\vec{b} + \overleftarrow{a}]\} & (\vec{b} + \overleftarrow{a} \neq 0) \end{cases} \\ \llbracket n_{\forall, p}^{a, b} \rrbracket_M &= \begin{cases} P(\emptyset) \times [\vec{a} + \overleftarrow{b}] & (\vec{b} + \overleftarrow{a} = 0) \\ \{(T, i) \in P(\mathbb{N}_M \times [\vec{b} + \overleftarrow{a}]) \times [\vec{a} + \overleftarrow{b}] \mid \{(j, p) \mid j \in [\vec{b} + \overleftarrow{a}]\} \subseteq T\} & (\vec{b} + \overleftarrow{a} \neq 0) \end{cases} \end{aligned}$$

Both morphisms above are from $\llbracket a \rrbracket_M = ([\vec{a}], [\overleftarrow{a}])$ to $\llbracket b \rrbracket_M = ([\vec{b}], [\overleftarrow{b}])$ in $\text{Int}(\mathbf{FinScottL}_M^{op})$, i.e., $P(\mathbb{N}_M \times [\vec{b} + \overleftarrow{a}]) \leftrightarrow [\vec{a} + \overleftarrow{b}]$ in $\mathbf{FinScottL}$, where powerset is ordered by inclusion.

5 Strategies and the Winning Position Functor

In §5.1, we define notions of play and strategy for open parity games (in the traditional style of graph game), as well as winning, losing, and *pending* strategies and positions. We use these definitions in §5.2 to define the *winning-position* functor $\mathcal{W}_M^r : \mathbf{OPG}_M^r \rightarrow \mathbf{FinScottL}_M^{op}$, which gives information that allows compositional computation of winning positions. We show that the diagram in Figure 3 commutes, which gives a justification of our compositional approach to parity games. In this section, we assume that a given rightward open parity game \mathcal{A} is of the form $\mathcal{A} = (m, n, Q, E, \rho, M, \omega)$.

5.1 Winning Strategies and Winning Positions for Open Parity Games

Here we give the notions of strategy and play. We also define the *denotation* of a strategy/position, which is how they win, lose, or are pending. These definitions are given only for *rightward* open parity games, but we can readily extend them to general open parity games, because any open parity game is a rightward open parity game by definition: $\mathbf{OPG}_M((m_r, m_l), (n_r, n_l)) = \mathbf{OPG}_M^r(m_r + n_l, n_r + m_l)$.

First we define the notion of strategy. For an open parity game $\mathcal{A} : [m] \rightarrow [n]$, a family $(s_i)_{i \in I}$ of positions in $[m] + [n] + Q$ is called a *position sequence* if $I = \mathbb{N}_{\geq 1}$ or $I = \{1, \dots, k\}$ for some $k \in \mathbb{N}_{\geq 1}$, (in that case, we also write $(s_i)_{i \in I}$ as $s_1 \cdots s_k$).

Definition 5.1 (\exists -strategy and \forall -strategy) Let \mathcal{A} be a rightward open parity game from m to n . We define $\text{Play}_{\exists} := \{s_1 \cdots s_k \mid k \geq 1, s_i \in Q (i \in [k]), (s_i, s_{i+1}) \in E (i \in [k-1]), \rho(s_k) = \exists\}$. Then, an \exists -strategy on \mathcal{A} is a partial function $\tau : \text{Play}_{\exists} \rightarrow [n] + Q$ where for any $s_1 \cdots s_k \in \text{Play}_{\exists}$, (i) if $\tau(s_1 \cdots s_k) = s$, then $(s_k, s) \in E$, and (ii) if $\tau(s_1 \cdots s_k)$ is undefined, then for all $s \in [n] + Q$, $(s_k, s) \notin E$. A \forall -strategy on an open parity game \mathcal{A} is defined in the same way, by replacing the occurrence of \exists with \forall in the above definition. The sets of \exists -strategies and \forall -strategies on \mathcal{A} are $\text{Str}_{\exists}(\mathcal{A})$ and $\text{Str}_{\forall}(\mathcal{A})$, respectively.

A pair of an \exists -strategy and a \forall -strategy resolves the non-determinism in a game to induce a unique play:

Definition 5.2 (induced play $\text{play}_a^{\tau_\exists, \tau_\forall}$) Let \mathcal{A} be a rightward open parity game from m to n . The *induced play* $\text{play}_a^{\tau_\exists, \tau_\forall}$ from an entry position $a \in [m]$ by an \exists -strategy τ_\exists and a \forall -strategy τ_\forall is the (necessarily unique) maximal position-sequence $(s_i)_{i \in I}$ (for the prefix order) such that: (i) $a E s_1$, (ii) for any $i \in I$, if $\rho(s_i) = \exists$ and $\tau_\exists(s_1 \cdots s_i)$ is defined, then $\tau_\exists(s_1 \cdots s_i) = s_{i+1}$, and similarly (iii) for any $i \in I$, if $\rho(s_i) = \forall$ and $\tau_\forall(s_1 \cdots s_i)$ is defined, then $\tau_\forall(s_1 \cdots s_i) = s_{i+1}$.

The following notion for a play corresponds to the *winning condition* in (traditional) game theory, where the condition is two-valued, “win” or “lose”. Below \exists and \forall correspond to “win” and “lose”, but we have other intermediate results $(m, s_{|I})$ due to the openness, which we call *pending states*. In this paper, we call the following many-valued winning/losing/pending condition $\llbracket - \rrbracket_{\mathcal{A}}$ on plays simply *winning condition*. An infinite position sequence $(s_i)_{i \in \mathbb{N}}$ satisfies the *parity condition* if the maximum of priorities that occur infinitely in the play is even. We apply the following notion $\llbracket - \rrbracket_{\mathcal{A}}$ only to induced plays.

Definition 5.3 (winning condition $\llbracket - \rrbracket_{\mathcal{A}}$ on plays) Let \mathcal{A} be a rightward open parity game. The *denotation* $\llbracket (s_i)_{i \in I} \rrbracket_{\mathcal{A}}$ of a position sequence $(s_i)_{i \in I}$ is defined as

$$\begin{aligned} (m, s_{|I}) & \text{ if } I \text{ is finite, } m = \max\{\omega(s_i) : i \in I\}, \text{ and } s_{|I} \text{ is an open end,} \\ \exists & \quad \text{if } (I \text{ is finite and } \rho(s_{|I}) = \forall) \text{ or } (I \text{ is infinite and } (s_i)_{i \in I} \text{ satisfies the parity condition),} \\ \forall & \quad \text{if } (I \text{ is finite and } \rho(s_{|I}) = \exists) \text{ or } (I \text{ is infinite and } (s_i)_{i \in I} \text{ does not satisfy the parity condition).} \end{aligned}$$

We call the function $\llbracket - \rrbracket_{\mathcal{A}}$ the *winning condition* of \mathcal{A} .

Next, we define the denotation of an \exists -strategy; note that an \exists -strategy is a strategy for the “player” while \forall -strategies are those for the “opponent”. The denotation is “lose” if there is a losing play, and otherwise is the collection of all the pending states; if the collection is the empty set, then the denotation is “win”.

Definition 5.4 (denotation of positions and \exists -strategies) Let \mathcal{A} be a rightward open parity game from m to n . The *denotation* $\llbracket (a, \tau_\exists) \rrbracket$ of an entry position $a \in [m]$ and an \exists -strategy τ_\exists is defined by

$$\llbracket (a, \tau_\exists) \rrbracket := \begin{cases} \text{lose} & \text{if there is } \tau_\forall \text{ such that } \llbracket \text{play}_a^{\tau_\exists, \tau_\forall} \rrbracket_{\mathcal{A}} = \forall, \\ \{ \llbracket \text{play}_a^{\tau_\exists, \tau_\forall} \rrbracket_{\mathcal{A}} \in \mathbb{N}_M \times [n] \mid \tau_\forall \in \text{Str}_\forall(\mathcal{A}) \text{ and } \llbracket \text{play}_a^{\tau_\exists, \tau_\forall} \rrbracket_{\mathcal{A}} \neq \exists \} & \text{otherwise.} \end{cases}$$

5.2 The Winning Position Functor \mathcal{W}_M

Now we give the central notion of this section, the *winning position* functor \mathcal{W}_M , which is a compact closed functor constructed by the Int-construction of a traced symmetric strong monoidal functor \mathcal{W}_M^r . In the definition of \mathcal{W}_M^r , if we fix an entry position a , then $\mathcal{W}_M^r(\mathcal{A})$ (or precisely $\{T \mid (T, a) \in \mathcal{W}_M^r(\mathcal{A})\}$) is the upward-closed set generated by the denotations $\llbracket (a, \tau_\exists) \rrbracket$ of a and all \exists -strategies τ_\exists that does not lose from a :

Definition 5.5 (the functor \mathcal{W}_M^r) We define a functor $\mathcal{W}_M^r : \mathbf{OPG}_M^r \rightarrow \mathbf{FinScottL}_{!M}^{op}$. The mapping on objects is given by $\mathcal{W}_M^r(m) := ([m], =)$, and for a morphism $\mathcal{A} \in \mathbf{OPG}_M^r(m, n)$,

$$\mathcal{W}_M^r(\mathcal{A}) := \{ (T, a) \in P(\mathbb{N}_M \times [n]) \times [m] \mid \llbracket (a, \tau_\exists) \rrbracket \neq \text{lose} \text{ and } \llbracket (a, \tau_\exists) \rrbracket \subseteq T \text{ for some } \exists\text{-strategy } \tau_\exists \}.$$

The functor \mathcal{W}_M^r determines whether an entry position wins, but the precise perspective is as follows. As mentioned in the introduction, in the traditional notion of (non-open) parity games, a position is just either winning or losing, two-valued. With the new notion of open ends, however, we have the intermediate result of pending states. The following definition reflects this idea.

Definition 5.6 (winning/losing/pending positions) Let \mathcal{A} and a be a rightward open parity game and an entry position, respectively. (i) a is *winning* if $(\emptyset, a) \in \mathcal{W}_M^r(\mathcal{A})$, (ii) a is *losing* if $(T, a) \notin \mathcal{W}_M^r(\mathcal{A})$ for any T , and (iii) a is *pending* otherwise (i.e., if $(\emptyset, a) \notin \mathcal{W}_M^r(\mathcal{A})$ and $(T, a) \in \mathcal{W}_M^r(\mathcal{A})$ for some $T \neq \emptyset$).

There is an obvious transformation that maps a traditional parity game G and position x into an open parity game $\mathcal{A}_x^G : 1 \rightarrow 0 \in \mathbf{OPG}_M^r$, where $1 \in [1]$ points to the internal position x . The notion of winning/losing defined above agrees with the traditional one in the following sense (n.b. there is no pending case):

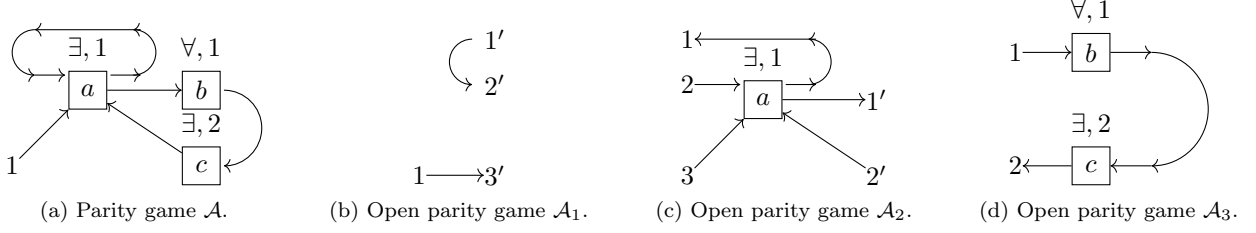


Fig. 8. An extended example.

Proposition 5.7 *Given a (traditional) parity game G and a position x in G , x is winning (resp. losing) in G iff x is winning (resp. losing) in \mathcal{A}_x^G .* \square

The main technical result of this section is stated below, and allows us to define a compact closed functor $\mathcal{W}_M : \mathbf{OPG}_M \rightarrow \mathbf{Int}(\mathbf{FinScottL}_{!M}^{op})$.

Theorem 5.8 *The functor $\mathcal{W}_M^r : \mathbf{OPG}_M^r \rightarrow \mathbf{FinScottL}_{!M}^{op}$ is a traced symmetric strict monoidal functor.* \square

Definition 5.9 (winning position functor \mathcal{W}_M) We define the *winning position functor* \mathcal{W}_M by $\mathbf{Int}(\mathcal{W}_M^r)$.

Summarizing all the main results in this paper, we obtain the following theorem:

Theorem 5.10 *The triangle in Fig. 3 commutes: $\llbracket - \rrbracket_M \simeq \mathcal{W}_M \circ \mathcal{R}_M$.* \square

We remark that we can obtain a similar result to the above in the TSMC setting by the freeness of $\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}})$.

Given any open parity game, which can be represented also by a morphism in $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ by the fullness of \mathcal{R}_M (Thm. 3.15), the above Thm. 5.10 says that we can calculate whether an entry position is winning, losing, or pending, either (i) by calculating strategies (i.e., by \mathcal{W}_M), or *equivalently* (ii) by induction (i.e., by $\llbracket - \rrbracket_M$) without calculating strategies. An elaborated example on how we can compute the denotation of an entry position of an open parity game by the induction $\llbracket - \rrbracket_M$ can be found below.

Finally, note that the notion of winning/losing/pending *position* is defined for $\mathbf{Int}(\mathbf{FinScottL}_{!M}^{op})$, and hence is defined also for $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ and \mathbf{OPG}_M , by using $\llbracket - \rrbracket_M$ and \mathcal{W}_M , respectively. On the other hand, the notion of winning/losing/pending *strategy* is defined for morphisms of \mathbf{OPG}_M (and hence of $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$, too) but not of $\mathbf{Int}(\mathbf{FinScottL}_{!M}^{op})$. In particular, we can conclude that we have given an *abstract* (or *extensional*) semantics for open parity games, by eliminating the information of strategies.

Example 5.11 Let \mathcal{A} be the open parity game in Fig. 8(a). We want to check whether the position 1 is winning by composing the interpretations of \mathcal{A} 's subgames. For more details, see Appendix D.

Concretely, \mathcal{A} is divided as $\mathcal{A}_3 \circ \mathcal{A}_2 \circ \mathcal{A}_1$ with \mathcal{A}_1 , \mathcal{A}_2 , and \mathcal{A}_3 shown in Fig. 8(b), 8(c), and 8(d), respectively (note that open ends are labelled using prop-style ordering, and while the ordering in $\mathbf{Int}(\mathbf{FinScottL}_{!M}^{op})$ is different, we keep the same notations for readability). It follows directly by unfolding definitions and by compact closedness of $\llbracket - \rrbracket_M$ that $\llbracket \mathcal{A}_1 \rrbracket_M = \{(T, 1) \mid (0, 3') \in T\} \cup \{(T, 1') \mid (0, 2') \in T\}$, $\llbracket \mathcal{A}_2 \rrbracket_M = \{(T, i) \mid i \in \{2, 3, 2'\}, \exists j \in \{1, 1'\}, (1, j) \in T\}$, and $\llbracket \mathcal{A}_3 \rrbracket_M = \{(T, 1) \mid (2, 2) \in T\}$, which are indeed the expected results. For example, to compute $\llbracket \mathcal{A}_1 \rrbracket_M$, we can decompose \mathcal{A}_1 as $d_{\mathbb{1}} \oplus \text{id}_{\mathbb{r}}$, so $\llbracket \mathcal{A}_1 \rrbracket_M = d_{\llbracket \mathbb{1} \rrbracket_M} \oplus \text{id}_{\llbracket \mathbb{r} \rrbracket_M}$, which can easily be computed from the definition of the identity in $\mathbf{FinScottL}_{!M}^{op}$.

In order to compute the composition of two interpretations in $\mathbf{Int}(\mathbf{FinScottL}_{!M}^{op})$, we need to compute a trace, and therefore semantic run-trees (we can avoid it in \mathcal{A}_2 and \mathcal{A}_3 above because they can be reorganized so that composition involves a trivial trace). In semantic run-trees corresponding to $\llbracket \mathcal{A}_2 \rrbracket_M \circ \llbracket \mathcal{A}_1 \rrbracket_M$, there must be no infinite path (corresponding to the intuition that the only infinite path is losing), and the only possible leaf is $1'$ (the only exit position), while conditions in Def. 4.6 (involving $\llbracket \mathcal{A}_2 \rrbracket_M$ and $\llbracket \mathcal{A}_1 \rrbracket_M$ above) ensure that $(1, 1')$ must be one of the leaves, which gives $\llbracket \mathcal{A}_2 \circ \mathcal{A}_1 \rrbracket_M = \{(T, i) \mid i \in \{1, 2'\}, (1, 1') \in T\}$.

Similarly, to compose $\llbracket \mathcal{A}_3 \rrbracket_M$ with $\llbracket \mathcal{A}_2 \circ \mathcal{A}_1 \rrbracket_M$, we also have to compute the corresponding semantic run-trees. Here, there can be no leaves (no exit positions), and the run-tree corresponding to taking the loop infinitely meets the parity condition (because all its nodes are $(2, 1')$, except for the root), so $(\emptyset, 1)$ is in the interpretation, whence $\llbracket \mathcal{A} \rrbracket_M = \{(T, 1) \mid \text{true}\} = \{(\emptyset, 1)\}$. Therefore, 1 is a winning position in G .

6 Conclusions and Future Work

We have given a compositional approach to parity games by exhibiting their underlying compact closed structure. Parity games can be composed by considering open ends, and we defined a prop that gives a graphical language to describe such open parity games. At the semantic level, we have given a notion of winning/losing positions that takes open ends into account. It retains enough information to be compositional, but is still extensional, as it can be computed without referring to starategies.

The current semantic category is a *strategy-insensitive* model, in that it only keeps track of the (non)existence of a winning strategy, while *strategy-sensitive* models should keep track of all strategies (perhaps up to some suitable equivalence). Strategy-sensitive models can easily be obtained when restricting to *history-free* strategies. One future work is to find some history-dependent strategy-sensitive model.

It could be fruitful to deepen the link between the existing body of work on props and our use of props in this work. For example, by showing equivalence between the graphical language and the category of open parity games, we could get decidability results on parity games from the syntax, as in [5].

Another possible future work is the coalgebraic treatment of open parity games. There seems a bijective correspondence between open parity games up to some notion of bisimilarity and pairs of functions $[m] \rightarrow Q$ and $Q \rightarrow P(Q + [n]) \times \{\exists, \forall\} \times \mathbb{N}_M$ up to the bisimilarity of coalgebras. Then there might exist a compact closed structure in the category of coalgebraic open parity games up to the bisimilarity. However, it seems subtle to give the same categorical structure in the category of, say, the above form of coalgebras up to *isomorphism*. Also, it might be interesting to give a functor (similar to \mathcal{R}_M) from some graphical category to the category of coalgebras up to bisimilarity by using bialgebraic methods [30, 38]. We did not take the bisimilarity approach in this paper because in game theory we basically consider the level up to isomorphism, say, for complexity.

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A The Int Construction

In this section, we recall the notions of traced symmetric monoidal category, compact closed category, and the Int construction. We assume that readers are familiar with *symmetric monoidal categories* [32]. A symmetric monoidal category $(\mathbb{C}, \otimes, \mathbf{I}, \mathbf{a}, \mathbf{l}, \mathbf{r}, \sigma)$ is a *strict symmetric monoidal category* if the associator \mathbf{a} , the left unitor \mathbf{l} , and the right unitor \mathbf{r} are the identity natural transformations. For readability, we deal with symmetric monoidal categories as if they were strict symmetric monoidal categories.

We begin with the definition of a *traced symmetric monoidal category* [25]. For another equivalent axiomatization, see [20].

Definition A.1 (traced symmetric monoidal category [25]) A *traced symmetric monoidal category* is a symmetric monoidal category $(\mathbb{C}, \otimes, \mathbf{I}, \sigma)$ equipped with a *trace operator* $\text{tr}_{Z;X,Y}^{\mathbb{C}} : \mathbb{C}(Z \otimes X, Z \otimes Y) \rightarrow \mathbb{C}(X, Y)$ that satisfies the following conditions:

$$\begin{aligned}
 \text{tr}_{\mathbf{I};X,Y}^{\mathbb{C}}(f) &= f & (f : X \rightarrow Y) & \quad \text{(Vanishing1)} \\
 \text{tr}_{U \otimes W;X,Y}^{\mathbb{C}}(f) &= \text{tr}_{W;X,Y}^{\mathbb{C}}(\text{tr}_{U;W \otimes X, W \otimes Y}^{\mathbb{C}}(f)) & (f : U \otimes W \otimes X \rightarrow U \otimes W \otimes Y) & \quad \text{(Vanishing2)} \\
 \text{tr}_{U;X,Y}^{\mathbb{C}}(f) \otimes g &= \text{tr}_{U;X \otimes W, Y \otimes Z}^{\mathbb{C}}(f \otimes g) & (f : U \otimes X \rightarrow U \otimes Y \text{ and } g : W \rightarrow Z) & \quad \text{(Superposing)} \\
 \text{tr}_{X;X,X}^{\mathbb{C}}(\sigma_{X,X}) &= \text{id}_X & & \quad \text{(Yanking)} \\
 \text{tr}_{U;X',Y}^{\mathbb{C}}(f \circ (\text{id}_U \otimes g)) &= \text{tr}_{U;X,Y}^{\mathbb{C}}(f) \circ g & (f : U \otimes X \rightarrow U \otimes Y \text{ and } g : X' \rightarrow X) & \quad \text{(Naturality in } X) \\
 \text{tr}_{U;X,Y'}^{\mathbb{C}}((\text{id}_U \otimes g) \circ f) &= g \circ \text{tr}_{U;X,Y}^{\mathbb{C}}(f) & (f : U \otimes X \rightarrow U \otimes Y \text{ and } g : Y \rightarrow Y') & \quad \text{(Naturality in } Y) \\
 \text{tr}_{U;X,Y}^{\mathbb{C}}((g \otimes \text{id}_Y) \circ f) &= \text{tr}_{U';X,Y}^{\mathbb{C}}(f \circ (g \otimes \text{id}_X)) & (f : U \otimes X \rightarrow U' \otimes Y \text{ and } g : U' \rightarrow U) & \quad \text{(Dinaturality in } U).
 \end{aligned}$$

Definition A.2 (traced symmetric strong monoidal functor [25]) Let $(\mathbb{C}, \otimes^{\mathbb{C}}, \mathbf{I}^{\mathbb{C}}, \sigma^{\mathbb{C}}, \text{tr}^{\mathbb{C}})$ and $(\mathbb{D}, \otimes^{\mathbb{D}}, \mathbf{I}^{\mathbb{D}}, \sigma^{\mathbb{D}}, \text{tr}^{\mathbb{D}})$ be traced symmetric monoidal categories. A *traced symmetric strong monoidal functor* $F : (\mathbb{C}, \otimes^{\mathbb{C}}, \mathbf{I}^{\mathbb{C}}, \sigma^{\mathbb{C}}, \text{tr}^{\mathbb{C}}) \rightarrow (\mathbb{D}, \otimes^{\mathbb{D}}, \mathbf{I}^{\mathbb{D}}, \sigma^{\mathbb{D}}, \text{tr}^{\mathbb{D}})$ is a symmetric strong monoidal functor $(F : \mathbb{C} \rightarrow \mathbb{D}, (\phi_{X,Y} : F(X) \otimes^{\mathbb{D}} F(Y) \rightarrow F(X \otimes^{\mathbb{C}} Y))_{X,Y}, \phi^0 : \mathbf{I}^{\mathbb{D}} \rightarrow F(\mathbf{I}^{\mathbb{C}}))$ that satisfies the following condition:

$$\text{tr}_{F(Z);F(X),F(Y)}^{\mathbb{D}}(\phi_{F(Z),F(Y)}^{-1} \circ F(f) \circ \phi_{F(Z),F(X)}) = F(\text{tr}_{Z;X,Y}^{\mathbb{C}}(f)) \quad (f : Z \otimes^{\mathbb{C}} X \rightarrow Z \otimes^{\mathbb{C}} Y).$$

A *compact closed category* [10] is a symmetric monoidal category where each object A has a dual object A^\perp ; one typical example is the category of finite dimensional vector spaces V , which have the dual spaces V^* . For the notion of *strict compact closed category*, see [28, §9] (or Rem. B.2).

Definition A.3 (compact closed category) A *compact closed category* is a symmetric monoidal category $(\mathbb{C}, \otimes, \mathbf{I}, \sigma)$ that has, for every object A , a chosen triple of a *(left) dual object* A^\perp , an unit $d_A : \mathbf{I} \rightarrow A \otimes A^\perp$, and a counit $e_A : A^\perp \otimes A \rightarrow \mathbf{I}$, and satisfies the following conditions:

$$\begin{aligned}
 (e_A \otimes \text{id}_{A^\perp}) \circ (\text{id}_{A^\perp} \otimes d_A) &= \text{id}_{A^\perp}, \\
 (\text{id}_A \otimes e_A) \circ (d_A \otimes \text{id}_A) &= \text{id}_A.
 \end{aligned}$$

A *compact closed functor* is defined to be just a symmetric strong monoidal functor. We note that every symmetric strong monoidal functor between compact closed functors preserves the dual objects up to canonical isomorphism.

Every compact closed category has the *canonical trace operator*. For example, the application of the trace operator in Fig. 2(c) is obtained by the unit and the counit as in Fig. 2(b); see [25] for the formal definition. Thus, there is an embedding from compact closed categories to traced symmetric monoidal categories. Conversely, there is a free construction of a compact closed category $\text{Int}(\mathbb{C})$ from a traced symmetric monoidal category \mathbb{C} . This free construction is called the *Int construction* [25].

Definition A.4 (the Int construction of 0-cell [25]) Let $(\mathbb{C}, \otimes, \mathbf{I}, \sigma, \text{tr})$ be a traced symmetric monoidal category. We define a compact closed category $(\text{Int}(\mathbb{C}), \otimes^{\text{Int}(\mathbb{C})}, \mathbf{I}^{\text{Int}(\mathbb{C})}, \sigma^{\text{Int}(\mathbb{C})})$ by the following:

- (i) An object of $\text{Int}(\mathbb{C})$ is a pair (X_+, X_-) of objects X_+, X_- in \mathbb{C} .
- (ii) A morphism f in $\text{Int}(\mathbb{C})((X_+, X_-), (Y_+, Y_-))$ is a morphism $f \in \mathbb{C}(X_+ \otimes Y_-, Y_+ \otimes X_-)$.
- (iii) Let $f \in \text{Int}(\mathbb{C})((X_+, X_-), (Y_+, Y_-))$ and $g \in \text{Int}(\mathbb{C})((Y_+, Y_-), (Z_+, Z_-))$. The composite of f and g is $\text{tr}_{Y_-;X_+ \otimes Z_-, Z_+ \otimes X_-}^{\mathbb{C}}((\sigma_{Z_+, Y_-} \otimes \text{id}_{X_-}) \circ (g \otimes \text{id}_{X_-}) \circ (\text{id}_{Y_+} \otimes \sigma_{X_-, Z_-}) \circ (f \otimes \text{id}_{Z_-}) \circ (\sigma_{Y_-, X_+} \otimes \text{id}_{Z_-}))$.

- (iv) The identity over (X_+, X_-) is defined by $\text{id}_{(X_+, X_-)} := \text{id}_{X_+ \otimes X_-}$.
- (v) The tensor product $(X_+, X_-) \otimes^{\text{Int}(\mathbb{C})} (Y_+, Y_-)$ is $(X_+ \otimes Y_+, Y_- \otimes X_-)$, and $(f : (X_+, X_-) \rightarrow (Y_+, Y_-)) \otimes^{\text{Int}(\mathbb{C})} (g : (X'_+, X'_-) \rightarrow (Y'_+, Y'_-))$ is $(\text{id}_{Y_+} \otimes g \otimes \text{id}_{X_-}) \circ (\sigma_{X'_+, Y_+} \otimes \sigma_{X_-, Y'_-}) \circ (\text{id}_{X'_+} \otimes f \otimes \text{id}_{Y'_-}) \circ (\sigma_{X_+, X'_+} \otimes \sigma_{Y'_-, Y_-})$.
- (vi) The monoidal unit $\mathbf{I}^{\text{Int}(\mathbb{C})}$ is (\mathbf{I}, \mathbf{I}) .
- (vii) The swap $\sigma_{(X_+, X_-), (Y_+, Y_-)}^{\text{Int}(\mathbb{C})}$ is $\sigma_{(X_+, Y_+)} \otimes \sigma_{(Y_-, X_-)}$.
- (viii) (X_+, X_-) has as dual object (X_-, X_+) , with $\text{id}_{X_+ \otimes X_-}$ as a unit $d_{(X_+, X_-)}^{\text{Int}(\mathbb{C})}$, and $\text{id}_{X_- \otimes X_+}$ as a counit $e_{(X_+, X_-)}^{\text{Int}(\mathbb{C})}$.

Definition A.5 (the Int construction of 1-cell [25]) Let \mathbb{C} and \mathbb{D} be traced symmetric monoidal categories, and $F : \mathbb{C} \rightarrow \mathbb{D}$ be a traced symmetric strong monoidal functor. We define a compact closed functor $\text{Int}(F) : \text{Int}(\mathbb{C}) \rightarrow \text{Int}(\mathbb{D})$ by the following:

$$\begin{aligned} \text{Int}(F)((X_+, X_-)) &:= (F(X_+), F(X_-)), \\ \text{Int}(F)(f : (X_+, X_-) \rightarrow (Y_+, Y_-)) &:= (\phi_{Y_+, X_-}^F)^{-1} \circ F(f) \circ \phi_{X_+, Y_-}^F, \\ \phi_{(X_+, X_-), (Y_+, Y_-)}^{\text{Int}(F)} &:= \phi_{X_+, Y_+}^F \otimes (\phi_{Y_-, X_-}^F)^{-1}, \\ \phi^{\text{Int}(F), 0} &:= \text{id}_{F(\mathbf{I}^{\mathbb{C}})}. \end{aligned}$$

Definition A.6 Let \mathbb{C} be a traced symmetric monoidal category. We define a traced symmetric strong monoidal functor $\eta_{\mathbb{C}}^{\text{Int}} : \mathbb{C} \rightarrow \text{Int}(\mathbb{C})$ by the following:

$$\begin{aligned} \eta_{\mathbb{C}}^{\text{Int}}(X) &:= (X, \mathbf{I}^{\mathbb{C}}), \\ \eta_{\mathbb{C}}^{\text{Int}}(f) &:= f, \\ \phi_{X, Y}^{\eta_{\mathbb{C}}^{\text{Int}}} &:= \text{id}_{X \otimes Y}, \\ \phi^{\eta_{\mathbb{C}}^{\text{Int}}, \mathbf{I}^{\mathbb{C}}} &:= \text{id}_{\mathbf{I}^{\mathbb{C}}}. \end{aligned}$$

The following proposition states that the Int construction is a free construction.

Theorem A.7 ([25] and [21]) *The embedding of the 2-category \mathbf{CpCC} in the 2-category \mathbf{TrSMC}_g has a left biadjoint Int whose unit is given by the functors $(\eta_{\mathbb{C}}^{\text{Int}} : \mathbb{C} \rightarrow \text{Int}(\mathbb{C}))_{\mathbb{C} \in \mathbf{TrSMC}_g}$.* \square

B Free Compact Closed Categories and Fullness of the Functor \mathcal{R}_M

Here we explain the free compact closedness of $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ with a 2-categorical and generalized statement (§B.1), the relationship between the free compact closed categories and the free traced symmetric monoidal categories by Kissinger (§B.2), and the fullness of \mathcal{R}_M (§B.3).

Given categories \mathbb{C} and \mathbb{D} , we write $[\mathbb{C}, \mathbb{D}]$ for the functor category, and we write $[\mathbb{C}, \mathbb{D}]_g$ for the subcategory of functors and natural isomorphisms.

B.1 Free Compact Closedness of $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$

Here we give a 2-categorical and generalized version of Thm. 3.13, and suppose the familiarity to 2-category theory. We first give a generalization of the notion of CCS to arbitrary color sets C :

Definition B.1 We call an object of the category $\mathbf{Set}^{(C+C)^* \times (C+C)^*}$ a *(C-colored) compact closed signature* (*C-CCS*, for short).

We define an involution operator $(-)^{\perp} : (C+C)^* \rightarrow (C+C)^*$ by $(\text{in}_i(c))^{\perp} := \text{in}_{3-i}(c)$ (for $i = 1, 2$), $\varepsilon^{\perp} := \varepsilon$, and $(x_1 \cdots x_n)^{\perp} := x_n^{\perp} \cdots x_1^{\perp}$. Then $(C+C)^*$ with embedding $\text{in}_1(-) : C \rightarrow (C+C)^*$ is the free involutive monoid generated by C . (An *involutive monoid* is a monoid with unary operator $(-)^{\perp}$ and the following axioms: $(x^{\perp})^{\perp} = x$, $(xy)^{\perp} = y^{\perp}x^{\perp}$, and redundantly $e^{\perp} = e$.) We write $\text{in}_1(c) \in (C+C)^*$ simply as c . We identify $\{\mathfrak{r}, \mathfrak{l}\}^*$ with $(\{\mathfrak{r}\} + \{\mathfrak{l}\})^*$; thus the notion of CCS defined in §3.2 agrees with the notion of $\{\mathfrak{r}\}$ -colored CCS, i.e., *single-colored* CCS.

Remark B.2 Def. B.1 is inspired by the notion of *strict compact closed category* [28]. A compact closed category \mathbb{C} is *strict* if \mathbb{C} is strict as a monoidal category and moreover the following canonical isomorphisms are the identities: $(A^\perp)^\perp \cong A$, $(A \otimes B)^\perp \cong B^\perp \otimes A^\perp$, and $I^\perp \cong I$. Thus we are led to use the notion of involutive monoid in the definitions of the compact closed version of the notions of prop and symmetric monoidal signature, though we do not pursue the detail in this paper.

Next we generalize $\{\mathfrak{r}, \mathfrak{l}\}$ -SMT $(\Sigma_M^{\text{de}}, E_M^{\text{de}})$ given in Def. 3.6 as follows:

Definition B.3 (SMT $T_\Sigma^{\text{de}} = (\Sigma^{\text{de}}, E_\Sigma^{\text{de}})$ of CCS Σ) Given a C -CCS Σ , we define a $(C + C)$ -SMT $T_\Sigma^{\text{de}} = (\Sigma^{\text{de}}, E_\Sigma^{\text{de}}, l, r)$ as follows:

$$\begin{aligned} \Sigma^{\text{de}}(\varepsilon, c \cdot c^\perp) &:= \Sigma(\varepsilon, c \cdot c^\perp) + \{d_c\} && \text{(for } c \in C), \\ \Sigma^{\text{de}}(c^\perp \cdot c, \varepsilon) &:= \Sigma(c^\perp \cdot c, \varepsilon) + \{e_c\} && \text{(for } c \in C), \\ \Sigma^{\text{de}}(w, u) &:= \Sigma(w, u) && \text{(for other } w, u \in (C + C)^*), \end{aligned}$$

and $E_\Sigma^{\text{de}}(c, c) := \{(\text{id}_c \oplus e_c) \circ (d_c \oplus \text{id}_c) = \text{id}_c\}$, $E_\Sigma^{\text{de}}(c^\perp, c^\perp) := \{(e_c \oplus \text{id}_{c^\perp}) \circ (\text{id}_{c^\perp} \oplus d_c) = \text{id}_{c^\perp}\}$, and $E_\Sigma^{\text{de}}(w, u) := \emptyset$ otherwise, where the monoidal category structures id, \circ, \oplus are those of $\mathbb{F}_{\text{sig}}^{C+C}(\Sigma^{\text{de}})$. Then l and r are the left hand side and right hand side, respectively.

Note that the $\{\mathfrak{r}, \mathfrak{l}\}$ -SMT $(\Sigma_M^{\text{oppg}}, E_M^{\text{oppg}})$ is nothing but $T_{\Sigma_M}^{\text{de}} = (\Sigma_M^{\text{de}}, E_{\Sigma_M}^{\text{de}})$, where Σ_M is given in Def. 3.10.

Above, the unit d_c and counit e_c are given just for a color $c \in C$, but we can extend them to all the elements of $(C + C)^*$ in an obvious way so that the following coherence conditions for *strict* compact closed category [28, Sections 6 and 9] hold:

$$\begin{aligned} d_\varepsilon &= \text{id}_\varepsilon, & d_{w \cdot u} &= (\text{id}_w \oplus d_u \oplus \text{id}_{d_{w^\perp}}) \circ d_w, & d_{w^\perp} &= \sigma_{w, w^\perp} \circ d_w, \\ e_\varepsilon &= \text{id}_\varepsilon, & e_{w \cdot u} &= e_u \circ (\text{id}_{u^\perp} \oplus e_w \oplus \text{id}_u), \quad \text{and} & e_{w^\perp} &= e_w \circ \sigma_{w, w^\perp}. \end{aligned}$$

where $\sigma_{w, u} : w \cdot u \rightarrow u \cdot w$ is the swap. Then, we also have the triangular identities for d_w and e_w . (To construct a prop isomorphic to $\mathbb{F}(T_\Sigma^{\text{de}})$, we can alternatively use the $(C + C)$ -SMT that is obtained by adding to T_Σ^{de} all d_w, e_w , and the triangular identities for them as well as the above coherence equations, among which the former (or latter) three are sufficient.)

We generalize Def. 3.11 and 3.12 with arbitrary C -CCSs, and also extend them by dealing with also morphisms.

Definition B.4 (valuation, $[\Sigma, \mathbb{C}]_g^{\text{cpt}}$) For a C -CCS Σ and a compact closed category \mathbb{C} , we define the category $[\Sigma, \mathbb{C}]_g^{\text{cpt}}$ as follows. An object V , called a *valuation*, is a pair $(V, (V_{w, u})_{w, u})$ such that

$$V : C \rightarrow \text{ob}(\mathbb{C}) \quad V_{w, u} : \Sigma(w, u) \rightarrow \mathbb{C}(V^*(w), V^*(u)) \quad (w, u \in (C + C)^*)$$

where $V^*(w)$ for $w \in (C + C)^*$ is defined by the following: $V^*(\varepsilon) := I$, and for $w \in (C + C)^*$, $V^*(w \cdot c) := V^*(w) \otimes V(c)$ and $V^*(w \cdot c^\perp) := V^*(w) \otimes V(c)^\perp$. A morphism $\theta : V \rightarrow W$, called a *valuation morphism*, is a family $(\theta_c : V(c) \rightarrow W(c))_{c \in C}$ of isomorphisms in \mathbb{C} such that for all $f \in \Sigma(w, u)$,

$$\theta_u^* \circ V_{w, u}(f) = W_{w, u}(f) \circ \theta_w^*$$

holds where $\theta_w^* : V^*(w) \rightarrow W^*(w)$ for $w \in (C + C)^*$ is defined by the following: $\theta_\varepsilon^* := \text{id}_I$, $\theta_{w \cdot c}^* := \theta_w^* \otimes \theta_c$, and $\theta_{w \cdot c^\perp}^* := \theta_w^* \otimes (\theta_c^\perp)^{-1}$.

Definition B.5 (action on valuations) For a C -CCS Σ and compact closed categories \mathbb{C} and \mathbb{D} , the *action* $(-\circ -)_{\mathbb{C}, \mathbb{D}, \Sigma}$, written also $-\circ-$ for short, is a functor $(-\circ -)_{\mathbb{C}, \mathbb{D}, \Sigma} : \mathbf{CpCC}(\mathbb{C}, \mathbb{D}) \times [\Sigma, \mathbb{C}]_g^{\text{cpt}} \rightarrow [\Sigma, \mathbb{D}]_g^{\text{cpt}}$ defined by the following:

$$F \circ V(c) := F(V(c)) \quad \text{and} \quad (F \circ V)_{w, u}(f) := (\phi_u^F)^{-1} \circ F(V_{w, u}(f)) \circ \phi_w^F$$

where the isomorphisms ϕ_w^F (and ϕ_u^F) are defined as follows: For $w = c_1^{i_1} \cdots c_n^{i_n} \in (C + C)^*$, where each $c_j^{i_j}$ is either c_j or c_j^\perp , we let $\phi_{c_1^{i_1} \cdots c_n^{i_n}}^F : F(V(c_1))^{i_1} \otimes \cdots \otimes F(V(c_n))^{i_n} \rightarrow F(V(c_1)^{i_1}) \otimes \cdots \otimes F(V(c_n)^{i_n})$ be the

morphism obtained by the fact that F is a compact closed functor. Given an $\alpha \in \mathbf{CpCC}(\mathbb{C}, \mathbb{D})(F, G)$ and $\theta \in [\Sigma, \mathbb{C}]_g^{\text{cpt}}(V, W)$, we define $(\alpha \circ \theta)_c := \alpha_{W(c)} \circ F(\theta_c) : F(V(c)) \rightarrow G(W(c))$.

Now we give the theorem of free compact closed categories. This statement is in a similar style to that in [24] for free (symmetric) monoidal categories and to that in [29] for free traced symmetric monoidal categories.

Theorem B.6 (free compact closed categories) *For a C -CCS Σ , $\mathbb{F}(T_\Sigma)$ is a strict compact closed category. Furthermore, $\mathbb{F}(T_\Sigma)$ is a free compact closed category, i.e., there exists a valuation $\eta_\Sigma^{\text{cpt}} \in \text{ob}([\Sigma, \mathbb{F}(T_\Sigma)]_g^{\text{cpt}})$ such that, for any compact closed category \mathbb{C} ,*

$$- \circ \eta_\Sigma^{\text{cpt}} : \mathbf{CpCC}(\mathbb{F}(T_\Sigma), \mathbb{C}) \rightarrow [\Sigma, \mathbb{C}]_g^{\text{cpt}}$$

yeilds an equivalence of categories. □

For proving Thm. B.6, we need some preparations. We will use 2-coequalizer in \mathbf{SMC} , therefore we define functors $\widehat{\cdot}$ and $\widetilde{\cdot}$ between $[\Sigma, \mathbb{C}]_g^{\text{cpt}}$ and $[\Sigma^{\text{de}}, \mathbb{C}]^{\text{smc}}$, where the category $[\Sigma^{\text{de}}, \mathbb{C}]^{\text{smc}}$ is defined in [24].

Lemma B.7 *There are functors $\widehat{\cdot} : [\Sigma^{\text{de}}, \mathbb{C}]^{\text{smc}} \rightarrow [\Sigma, \mathbb{C}]_g^{\text{cpt}}$ and $\widetilde{\cdot} : [\Sigma, \mathbb{C}]_g^{\text{cpt}} \rightarrow [\Sigma^{\text{de}}, \mathbb{C}]^{\text{smc}}$:*

$$\begin{aligned} \widehat{V}(c) &:= V(c) \\ \widehat{V}_{w,u}(f) &:= V(f) \\ \widetilde{V}(c) &:= V(c) \\ \widetilde{V}_{w,u}(f) &:= \begin{cases} d & \text{if } f = d \\ e & \text{if } f = e \\ V(f) & \text{otherwise} \end{cases} \end{aligned}$$

Then $\widehat{\cdot} \circ \widetilde{\cdot} = \text{id}$.

Proof. Trivial. □

Obviously, actions, and $\widehat{\cdot}$ and $\widetilde{\cdot}$ are commute.

Lemma B.8 *Let $F \in \mathbf{CpCC}(\mathbb{C}, \mathbb{D})$, $V \in [\Sigma, \mathbb{C}]_g^{\text{cpt}}$, and $W \in [\Sigma, \mathbb{C}]^{\text{smc}}$. Then, $\widehat{F \circ V} \simeq F \circ \widehat{V}$ and $\widetilde{F \circ W} \simeq \widetilde{F} \circ \widetilde{W}$.*

Proof. Trivial. □

The freeness of \mathbb{F}_{sig} as a prop in Thm. 3.4 extends to some 2-dimensional freeness as a symmetric monoidal category. First we recall the 2-dimensional freeness, postponing its relation to \mathbb{F}_{sig} :

Theorem B.9 ([24]) *Let Σ be a C -signature. Then there exist a prop $\mathbb{F}_{\text{smc}}(\Sigma)$ and a valuation $\eta_\Sigma^{\text{smc}} \in \text{ob}([\Sigma, \mathbb{F}_{\text{smc}}(\Sigma)]^{\text{smc}})$ with $\eta_\Sigma^{\text{smc}}(c) = c$ ($c \in C$) such that for any symmetric monoidal category \mathbb{C} ,*

$$- \circ \eta_\Sigma^{\text{smc}} : \mathbf{SMC}(\mathbb{F}_{\text{smc}}(\Sigma), \mathbb{C}) \rightarrow [\Sigma, \mathbb{C}]^{\text{smc}}$$

yeilds an equivalence of categories.

Moreover, if \mathbb{C} is a strict symmetric monoidal category, then for any $V \in [\Sigma, \mathbb{C}]^{\text{smc}}$, there exists unique (up to equality) strict symmetric monoidal functor $F : \mathbb{F}_{\text{smc}}(\Sigma) \rightarrow \mathbb{C}$ such that $F = V \circ \eta_\Sigma^{\text{smc}}$.

Proof. The former statement is just Thm. 2.3 in [24], where the fact that $\mathbb{F}_{\text{smc}}(\Sigma)$ is a prop and $\eta_\Sigma^{\text{smc}}(c) = c$ is shown in the proof. The latter one is not explicitly found in [24]; but its (not necessarily symmetric) monoidal version is written in the proof of Thm. 1.2 in [24] and the part of the proof is easily adapted to the symmetric monoidal version, i.e., the latter one in the current theorem. □

Then \mathbb{F}_{sig} (given in Thm. 3.4) is nothing but \mathbb{F}_{smc} :

Corollary B.10 *Let Σ be a C -signature. We have an isomorphism $\mathbb{F}_{\text{sig}}(\Sigma) \cong \mathbb{F}_{\text{smc}}(\Sigma)$ in C -Prop. This induces an obvious isomorphism $C\text{-Sig}(\Sigma, (\mathbb{U}_{\text{sig}} \circ \mathbb{F}_{\text{sig}})(\Sigma)) \cong \{V \in [\Sigma, \mathbb{F}_{\text{smc}}(\Sigma)]^{\text{smc}} \mid V(c) = c \ (c \in C)\}$, which maps η_Σ^{sig} to η_Σ^{smc} , where η_Σ^{sig} is the unit of $\mathbb{F}_{\text{sig}} \dashv \mathbb{U}_{\text{sig}}$.*

Proof. Let \mathbb{C} be an arbitrary C -prop. We have an obvious isomorphism

$$C\text{-Sig}(\Sigma, \mathbb{U}_{\text{sig}}(\mathbb{C})) \cong \{V \in [\Sigma, \mathbb{C}]^{\text{smc}} \mid V(c) = c (c \in C)\}.$$

By the latter statement of Thm. B.9, we have a bijection $- \circ \eta_{\Sigma}^{\text{smc}} : C\text{-Prop}(\mathbb{F}_{\text{smc}}(\Sigma), \mathbb{C}) \rightarrow \{V \in [\Sigma, \mathbb{C}]^{\text{smc}} \mid V(c) = c (c \in C)\}$. Then, we have

$$\begin{aligned} C\text{-Prop}(\mathbb{F}_{\text{sig}}(\Sigma), \mathbb{C}) &\cong C\text{-Sig}(\Sigma, \mathbb{U}_{\text{sig}}(\mathbb{C})) \\ &\cong \{V \in [\Sigma, \mathbb{C}]^{\text{smc}} \mid V(c) = c (c \in C)\} \\ &\cong C\text{-Prop}(\mathbb{F}_{\text{smc}}(\Sigma), \mathbb{C}). \end{aligned}$$

By the Yoneda lemma, we have $\mathbb{F}_{\text{sig}}(\Sigma) \cong \mathbb{F}_{\text{smc}}(\Sigma)$ in $C\text{-Prop}$. The remaining parts follow immediately. \square

By this corollary, we identify \mathbb{F}_{sig} with \mathbb{F}_{smc} and η^{sig} with η^{smc} . We denote $H : \mathbf{SMC}(\mathbb{F}_{\text{sig}}(\Sigma), \mathbb{C}) \rightarrow [\Sigma, \mathbb{C}]^{\text{smc}}$ as $- \circ \eta_{\Sigma}^{\text{smc}}$ and $G : [\Sigma, \mathbb{C}]^{\text{smc}} \rightarrow \mathbf{SMC}(\mathbb{F}_{\text{sig}}(\Sigma), \mathbb{C})$ is a corresponding inverse functor.

By using Thm. B.9 and Cor. B.10, we can prove a converse statement of Lem. B.7 with an additional condition.

Lemma B.11 *Let $\pi : \mathbb{F}_{\text{sig}}(\Sigma^{\text{de}}) \rightarrow \mathbb{F}(T_{\Sigma})$ is a canonical epic functor of 2-coequalizer. Then $\pi \circ \widetilde{\eta_{\Sigma^{\text{de}}}^{\text{smc}}} \simeq \pi \circ \eta_{\Sigma^{\text{de}}}^{\text{smc}}$.*

Proof. By Cor. B.10, we can identify $\widetilde{\pi \circ \eta_{\Sigma^{\text{de}}}^{\text{smc}}}$ with $\pi \circ \eta_{\Sigma^{\text{de}}}^{\text{sig}}$. By Cor. B.10, $\pi \circ \eta_{\Sigma^{\text{de}}}^{\text{sig}} \simeq \pi \circ \eta_{\Sigma^{\text{de}}}^{\text{smc}}$. \square

By the following lemma, we can use a universality of 2-coequalizer in \mathbf{SMC} .

Lemma B.12 *Let $V \in [\Sigma, \mathbb{C}]_g^{\text{cpt}}$. Then $G(\widetilde{V}) \circ \mathbb{F}_{\text{sig}}(l) = G(\widetilde{V}) \circ \mathbb{F}_{\text{sig}}(r)$.*

Proof. This is proved by the naturality of G . \square

Proof. (Proof of Thm. B.6) First, we define a functor $R : [\Sigma, \mathbb{C}]_g^{\text{cpt}} \rightarrow \mathbf{CpCC}(\mathbb{F}(T_{\Sigma}), \mathbb{C})$. Let $V \in [\Sigma, \mathbb{C}]_g^{\text{cpt}}$. Then we can get $\widetilde{V} \in [\Sigma^{\text{de}}, \mathbb{C}]^{\text{smc}}$ by Lem. B.7. By Thm. B.9, $G(\widetilde{V}) \in \mathbf{SMC}(\mathbb{F}_{\text{sig}}(\Sigma^{\text{de}}), \mathbb{C})$. By Lem. B.12, there is a unique functor $G(\widetilde{V})' \in \mathbf{SMC}(\mathbb{F}(T_{\Sigma}), \mathbb{C})$ such that $G(\widetilde{V})' \circ \pi = G(\widetilde{V})$ where $\pi : \mathbb{F}_{\text{sig}}(\Sigma^{\text{de}}) \rightarrow \mathbb{F}(T_{\Sigma})$ is a canonical epic functor of 2-coequalizer. Then, we define $R(V) := G(\widetilde{V})'$. Let $\alpha \in [\Sigma, \mathbb{C}]_g^{\text{cpt}}(V, W)$. Then we can get $\widetilde{\alpha} \in [\Sigma^{\text{de}}, \mathbb{C}]^{\text{smc}}(\widetilde{V}, \widetilde{W})$, and $G(\widetilde{\alpha}) \in \mathbf{SMC}(\mathbb{F}_{\text{sig}}(\Sigma^{\text{de}}), \mathbb{C})(G(\widetilde{V}), G(\widetilde{W}))$. By 2-coequalizer, there is a unique $G(\widetilde{\alpha})' \in \mathbf{SMC}(\mathbb{F}(T_{\Sigma}), \mathbb{C})(G(\widetilde{V})', G(\widetilde{W})')$. Then we define $R(\alpha) := G(\widetilde{\alpha})'$.

Next we define $\eta_{\Sigma}^{\text{cpt}}$ by the following:

$$\eta_{\Sigma}^{\text{cpt}} := \widehat{H(\pi)}$$

Next we define a natural isomorphism $\alpha : \text{Id} \Rightarrow R(-) \circ \eta_{\Sigma}^{\text{cpt}}$ by the following:

$$\begin{aligned} &R(V) \circ \eta_{\Sigma}^{\text{cpt}} \\ &= G(\widetilde{V})' \circ \widehat{H(\pi)} && \text{(by definition)} \\ &\simeq G(\widetilde{V})' \circ H(\pi) && \text{(Lem. B.8)} \\ &= \widehat{H(G(\widetilde{V}))} && \text{(by naturality of } H \text{ and definition of } \pi) \\ &\simeq V && \text{(Lem. B.7 and Thm. B.9)} \end{aligned}$$

Next we define a natural isomorphism $\beta : R(- \circ \eta_{\Sigma}^{\text{cpt}}) \Rightarrow \text{id}$ and we finish our proof:

$$\begin{aligned} &R(F \circ \eta_{\Sigma}^{\text{cpt}}) \\ &= G(F \circ \widehat{H(\pi)})' && \text{(by definition)} \\ &\simeq G(F \circ \widetilde{\widetilde{H(\pi)}})' && \text{(Lem. B.8)} \end{aligned}$$

$$\begin{aligned}
 &\simeq G(F \circ H(\pi))' && \text{(Lem. B.11)} \\
 &\simeq (F \circ \pi)' && \text{(by naturality of } H \text{ and Thm. B.9)} \\
 &= F && \text{(by definition of } \pi)
 \end{aligned}$$

□

B.2 Relation to the Free Traced Symmetric Monoidal Category

Here we show that we have an equivalence $\mathbb{F}(T_\Sigma^{\text{de}}) \simeq \text{Int}(\mathbb{F}_{\text{tr}}(\Sigma'))$ if Σ and Σ' have a certain relationship.

First we recall the result by Kissinger [29]. Differently from the original definition, in the next definition we restrict morphisms to isomorphisms, to adjust it to the compact closed setting (as in [21]).

Definition B.13 (valuation into traced symmetric monoidal category $[\Sigma, \mathbb{C}]_g^{\text{tr}}$) For a C -signature Σ and a traced symmetric monoidal category \mathbb{C} , we define the category $[\Sigma, \mathbb{C}]_g^{\text{tr}}$ as follows. An object V is a pair $(V, (V_{w,u})_{w,u})$ such that

$$V : C \rightarrow \text{ob}(\mathbb{C}) \quad V_{w,u} : \Sigma(w, u) \rightarrow \mathbb{C}(V^*(w), V^*(u)) \quad (w, u \in C^*)$$

where $V^*(c_1 \cdots c_n) := V(c_1) \otimes \cdots \otimes V(c_n)$. A morphism $\theta : V \rightarrow W$ is a family $(\theta_c : V(c) \rightarrow W(c))_{c \in C}$ of isomorphisms in \mathbb{C} such that for all $f \in \Sigma(w, u)$,

$$\theta_u^* \circ V_{w,u}(f) = W_{w,u}(f) \circ \theta_w^*$$

holds where $\theta_{c_1 \cdots c_n}^* := \theta_{c_1} \otimes \cdots \otimes \theta_{c_n} : V^*(c_1 \cdots c_n) \rightarrow W^*(c_1 \cdots c_n)$.

We omit the definition of the traced version of the action $(- \circ -)_{\mathbb{C}, \mathbb{D}, \Sigma}$ (which is defined similarly to the compact closed case), because it is not important in the proof of $\mathbb{F}(T_\Sigma^{\text{de}}) \simeq \text{Int}(\mathbb{F}_{\text{tr}}(\Sigma'))$ given below. The next theorem is changed from the original one in [29]: morphisms of $\mathbf{CpCC}(\mathbb{F}(T_\Sigma), \mathbb{C})$ and $[\Sigma, \mathbb{C}]_g^{\text{cpt}}$ are restricted to isomorphisms.

Theorem B.14 (free traced symmetric monoidal categories [29]) For a C -signature Σ , $\mathbb{F}_{\text{tr}}(\Sigma)$ is a strict traced symmetric monoidal category. Furthermore, $\mathbb{F}_{\text{tr}}(\Sigma)$ is a free traced symmetric monoidal category, i.e., there exists a valuation $\eta_\Sigma^{\text{tr}} \in \text{ob}([\Sigma, \mathbb{F}(T_\Sigma)]_g^{\text{tr}})$ such that, for any traced symmetric monoidal category \mathbb{C} ,

$$(- \circ \eta_\Sigma^{\text{tr}}) : \mathbf{TrSMC}_g(\mathbb{F}_{\text{tr}}(\Sigma), \mathbb{C}) \rightarrow [\Sigma, \mathbb{C}]_g^{\text{tr}}$$

yeilds an equivalence of categories. □

In order to show $\mathbb{F}(T_\Sigma^{\text{de}}) \simeq \text{Int}(\mathbb{F}_{\text{tr}}(\Sigma'))$, we next give the relationship between the different notions of C -CCS Σ and C -signature Σ' . Specifically, given a C -CCS $\Sigma \in \mathbf{Set}^{(C+C)^* \times (C+C)^*}$, we shall define the desired C -signature $\Sigma^{\text{int}} \in \mathbf{Set}^{C^* \times C^*}$ to be the left Kan extension of Σ along a function $\text{int}_C : (C+C)^* \times (C+C)^* \rightarrow C^* \times C^*$. The function int_C —inspired by Int —is defined as follows.

First let $[f, g] : C+C \rightarrow C^*$ be given by $f(c) = c$ and $g(c) = \varepsilon$, and then we have $[f, g]^\dagger : (C+C)^* \rightarrow C^*$ as the monoid homomorphism induced by $[f, g]$, and similarly, we have $[g, f]^\dagger : (C+C)^* \rightarrow C^*$. Then let $\nu := \langle [f, g]^\dagger, [g, f]^\dagger \rangle : (C+C)^* \rightarrow C^* \times C^*$, which, intuitively, “divides” colors in $w \in (C+C)^*$ into ones from the left C of $C+C$ and ones from the right C of $C+C$. Then we define $\text{int}_C : (C+C)^* \times (C+C)^* \rightarrow C^* \times C^*$ by

$$\begin{aligned}
 \text{int}_C(w, u) &:= (w_+ \cdot u_-, u_+ \cdot w_-) \\
 \text{where } (w_+, w_-) &= \nu(w) \\
 (u_+, u_-) &= \nu(u).
 \end{aligned}$$

Let L_{int_C} be the left adjoint to $\mathbf{Set}^{C^* \times C^*} : \mathbf{Set}^{(C+C)^* \times (C+C)^*} \rightarrow \mathbf{Set}^{(C+C)^* \times (C+C)^*}$ (where the left adjoint can be obtained by the left Kan extension). Then we define $\Sigma^{\text{int}_C} := L_{\text{int}_C}(\Sigma)$. We may write Σ^{int} for Σ^{int_C} . Concretely, $\Sigma^{\text{int}}(w, u) := \coprod_{(w', u') \in U_{(w, u)}} \Sigma(w', u')$ where $U_{(w, u)} := \{(w', u') \in (C+C)^* \times (C+C)^* : \text{int}_C(w', u') = (w, u)\}$.

With the left adjointness of Int (i.e. Thm. A.7) and of L_{int_C} , we can show the following as a corollary to Thm. B.6 (and Thm. B.14).

Corollary B.15 *For a C -CCS Σ , We have an equivalence $\mathbb{F}(T_\Sigma) \simeq \text{Int}(\mathbb{F}_{\text{tr}}(\Sigma^{\text{int}}))$ in \mathbf{CpCC} . Hence the full subcategory of $\mathbb{F}(T_\Sigma)$ whose objects are those in $(\text{in}_1(C))^* (\subseteq (C + C)^*)$ is isomorphic to $\mathbb{F}_{\text{tr}}(\Sigma^{\text{int}})$ as traced symmetric monoidal categories.*

Epecially, specializing to the $\{\mathfrak{r}\}$ -CCS Σ_M given in Def. 3.10, we have an equivalence $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \simeq \text{Int}(\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}}))$ in \mathbf{CpCC} , and the full subcategory of $\mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}})$ whose objects are those in $\{\mathfrak{r}\}^ (\subseteq \{\mathfrak{r}, \mathbb{1}\}^*)$ is isomorphic to $\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}})$ as traced symmetric monoidal categories.*

Thus, $\text{Int}(\mathbb{F}_{\text{tr}}(\Sigma^{\text{int}}))$ induces an alternative (up-to isomorphic) construction of $\mathbb{F}(T_\Sigma)$, by using the strictification of compact closed categories [28] (and then by adjusting objects to ones in $(C + C)^*$). Conversely, as in the corollary, the full subcategory of $\mathbb{F}(T_\Sigma)$ gives an alternative (up-to isomorphic) construction of $\mathbb{F}_{\text{tr}}(\Sigma^{\text{int}})$.

To show the above corollary easily, we prepare some definitions. In the next definition, when $V \in [C, \mathbb{C}]$ is given, $\mathbb{C}_C^{\text{m}}(V, V)$ gives the underlying C -signature of a monoidal category \mathbb{C} . Below, we use the notations V^* and θ^* given in Def. B.13

Definition B.16 Given a monoidal category \mathbb{C} , we define $\mathbb{C}_C^{\text{m}} : [C, \mathbb{C}]_g^{\text{op}} \times [C, \mathbb{C}]_g \rightarrow \mathbf{Set}^{C^* \times C^*}$ as follows: for $(V', V) \in [C, \mathbb{C}]_g^{\text{op}} \times [C, \mathbb{C}]_g$ and $(w, u) \in C^* \times C^*$,

$$\mathbb{C}_C^{\text{m}}(V', V)(w, u) := \mathbb{C}(V'^*(w), V^*(u)).$$

For $\theta' : W' \rightarrow V'$ and $\theta : V \rightarrow W$ in $[C, \mathbb{C}]_g$ and $(w, u) \in C^* \times C^*$, we define

$$\mathbb{C}_C^{\text{m}}(\theta', \theta)(w, u) := \mathbb{C}(\theta'_w, \theta_u) : \mathbb{C}(V'^*(w), V^*(u)) \rightarrow \mathbb{C}(W'^*(w), W^*(u)).$$

We also have the compact closed version of the above notion; below we use the notations V^* and θ^* given in Def. B.4.

Definition B.17 Given a compact closed category \mathbb{C} , we define $\mathbb{C}_C^{\text{c}} : [C, \mathbb{C}]_g^{\text{op}} \times [C, \mathbb{C}]_g \rightarrow \mathbf{Set}^{(C+C)^* \times (C+C)^*}$ as follows: for $(V', V) \in [C, \mathbb{C}]_g^{\text{op}} \times [C, \mathbb{C}]_g$ and $(w, u) \in (C + C)^* \times (C + C)^*$,

$$\mathbb{C}_C^{\text{c}}(V', V)(w, u) := \mathbb{C}(V'^*(w), V^*(u)).$$

For $\theta' : W' \rightarrow V'$ and $\theta : V \rightarrow W$ in $[C, \mathbb{C}]_g$ and $(w, u) \in (C + C)^* \times (C + C)^*$, we define

$$\mathbb{C}_C^{\text{c}}(\theta', \theta)(w, u) := \mathbb{C}(\theta'_w, \theta_u) : \mathbb{C}(V'^*(w), V^*(u)) \rightarrow \mathbb{C}(W'^*(w), W^*(u)).$$

The following is a variant of the notion of the category of elements:

Definition B.18 (category of elements of mixed-variant functor) Given a category \mathbb{C} and a functor $F : \mathbb{C}^{\text{op}} \times \mathbb{C} \rightarrow \mathbf{Set}$, we define the category $\text{el}^{C \in \mathbb{C}} F(C, C)$ as follows: An object is a pair of $X \in \mathbb{C}$ and $x \in F(C, C)$. A morphism $f : (X, x) \rightarrow (Y, y)$ is a morphism $f : X \rightarrow Y$ in \mathbb{C} such that $F(X, f)(x) = F(f, Y)(y)$. The identity and the composition is given by those of \mathbb{C} .

The construction $F \mapsto \text{el}^{C \in \mathbb{C}} F(C, C)$ extends to a functor from $[\mathbb{C}^{\text{op}} \times \mathbb{C}, \mathbf{Set}]$ to \mathbf{Cat} . Therefore:

Lemma B.19 *If F and G are isomorphic, then so are $\text{el}^{C \in \mathbb{C}} F(C, C)$ and $\text{el}^{C \in \mathbb{C}} G(C, C)$. \square*

Next we reformulate the notion of valuation.

Lemma B.20 *For a C -signature Σ and a traced symmetric monoidal category \mathbb{C} , we have an isomorphism*

$$[\Sigma, \mathbb{C}]_g^{\text{tr}} \cong \text{el}^{V \in [C, \mathbb{C}]_g} \mathbf{Set}^{C^* \times C^*}(\Sigma, \mathbb{C}_C^{\text{m}}(V, V)).$$

Also, for a C -CCS Σ and a compact closed category \mathbb{C} , we have an isomorphism

$$[\Sigma, \mathbb{C}]_g^{\text{cpt}} \cong \text{el}^{V \in [C, \mathbb{C}]_g} \mathbf{Set}^{(C+C)^* \times (C+C)^*}(\Sigma, \mathbb{C}_C^{\text{c}}(V, V)).$$

\square

Now we show Cor. B.15:

Proof. We only show the equivalence $\mathbb{F}(T_\Sigma) \simeq \text{Int}(\mathbb{F}_{\text{tr}}(\Sigma^{\text{int}}))$ in **CpCC**, from which the remaining parts follow easily.

We have

$$\begin{aligned}
 & \mathbf{CpCC}(\text{Int}(\mathbb{F}_{\text{tr}}(\Sigma^{\text{int}})), \mathbb{C}) \\
 & \simeq \mathbf{TrSMC}_g(\mathbb{F}_{\text{tr}}(\Sigma^{\text{int}}), \mathbb{C}) && (\text{Thm. A.7: Int is a left biadjoint}) \\
 & \simeq [\Sigma^{\text{int}}, \mathbb{C}]_g^{\text{tr}} && (\text{Thm. B.14: } \mathbb{F}_{\text{tr}} \text{ is the free TSMC construction}) \\
 & \simeq \text{el}^{V \in [C, \mathbb{C}]_g} \mathbf{Set}^{C^* \times C^*}(\Sigma^{\text{int}}, \mathbb{C}_C^{\text{m}}(V, V)). && (\text{Lem. B.20: reformulation of valuations}) \\
 & \simeq \text{el}^{V \in [C, \mathbb{C}]_g} \mathbf{Set}^{(C+C)^* \times (C+C)^*}(\Sigma, \mathbf{Set}^{\text{int}_C}(\mathbb{C}_C^{\text{m}}(V, V))). && (\Sigma^{\text{int}} = L_{\text{int}_C}(\Sigma) \text{ and } L_{\text{int}_C} \dashv \mathbf{Set}^{\text{int}_C})
 \end{aligned}$$

and also we have

$$\begin{aligned}
 & \mathbf{CpCC}(\mathbb{F}(T_\Sigma), \mathbb{C}) \\
 & \simeq [\Sigma, \mathbb{C}]_g^{\text{cpt}} && (\text{Thm. B.6: } \mathbb{F}(T_{(-)}) \text{ is the free CpCC construction}) \\
 & \simeq \text{el}^{V \in [C, \mathbb{C}]_g} \mathbf{Set}^{(C+C)^* \times (C+C)^*}(\Sigma, \mathbb{C}_C^{\text{c}}(V, V)). && (\text{Lem. B.20: reformulation of valuations})
 \end{aligned}$$

Thus, by Lem. B.19, we only need to show the following isomorphism

$$\begin{aligned}
 \mathbf{Set}^{\text{int}_C}(\mathbb{C}_C^{\text{m}}(V, V)) & \cong \mathbb{C}_C^{\text{c}}(V, V), && \text{i.e.,} \\
 \mathbb{C}(V^*(w_+ \cdot u_-), V^*(u_+ \cdot w_-)) & \cong \mathbb{C}(V^*(w), V^*(u)) && \text{where } (w_+ \cdot u_-, u_+ \cdot w_-) = \text{int}_C(w, u).
 \end{aligned}$$

This can be shown by using the internal adjointness of duals X^\perp in \mathbb{C} (recall that the dual X^\perp can be seen as an internal left adjoint to X in \mathbb{C} , where a strict monoidal category \mathbb{C} can be seen as the single-object 2-category [28]). \square

We also have $\text{inl} := (\text{in}_1)^* \times (\text{in}_1)^* : C^* \times C^* \rightarrow (C+C)^* \times (C+C)^*$ in the converse direction of int_C , and have the left adjoint L_{inl} to $\mathbf{Set}^{\text{inl}} : \mathbf{Set}^{(C+C)^* \times (C+C)^*} \rightarrow \mathbf{Set}^{C^* \times C^*}$. This gives a similar result to Cor. B.15 for a given C -signature Σ and $L_{\text{inl}}(\Sigma)$ (rather than for a given C -CCS Σ and $L_{\text{int}_C}(\Sigma)$). Especially, given a C -signature Σ , the full subcategory of $\mathbb{F}(T_{L_{\text{inl}}(\Sigma)})$ whose objects are those in $(\text{in}_1(C))^*$ ($\subseteq (C+C)^*$) gives an alternative (up-to isomorphic) construction of $\mathbb{F}_{\text{tr}}(\Sigma)$.

B.3 Fullness of the Functor \mathcal{R}_M

We prove Thm. 3.15, which says that the functor $\mathcal{R}_M : \mathbb{F}(\Sigma_M^{\text{opg}}, E_M^{\text{opg}}) \rightarrow \mathbf{OPG}_M$ is full. For this, it is sufficient to show that $\mathcal{R}_M^r : \mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}}) \rightarrow \mathbf{OPG}_M^r$ is full, because, if a traced symmetric strong monoidal functor F is full, so is $\text{Int}(F)$.

Let \mathcal{A} be a rightward open parity game from m to n . Let k be the size of E . We define a graphical game $f \in \mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}})(\mathbb{r}^{k+m}, \mathbb{r}^{k+n})$ so that $\mathcal{R}_M^r(\text{tr}_{k;m,n}^{\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}})}(f)) \sim \mathcal{A}$. We first take enumerations of positions $[m] + [n] + Q = \{s_1, s_2, \dots\}$ and of edges $E = \{e_1, e_2, \dots\}$. Then for each position s_i , we prepare a fresh node $\langle s_i \rangle$ (of the form $\mathbb{n}_{r,p}^{w,u}$) in the graphical language if s_i is an internal position, and prepare an identity game $\langle s_i \rangle$ on \mathbb{r} if s_i is an open end. Then we vertically compose all the games $\langle s_i \rangle$. For each internal position s_i , the role r and priority p of $\langle s_i \rangle = \mathbb{n}_{r,p}^{w,u}$ are obviously inherited from those of s_i in \mathcal{A} , and the domain w and codomain u will be determined below.

For each $\ell = 1, \dots, k$, if $e_\ell = (s_{i_\ell}, s_{j_\ell})$, then we add an edge from a fresh exit position of $\langle s_{i_\ell} \rangle$ to the exit position k of f , and similarly, we add an edge to a fresh entry position of $\langle s_{j_\ell} \rangle$ from the entry position ℓ of f . Also, for each $a \in [m]$, we add an edge from the entry position $k+a$ of f to the unique entry position of the identity game $\langle a \rangle$ prepared above, and similarly, for each $b \in [n]$, we add an edge to the exit position $k+b$ of f from the unique exit position of the identity game $\langle b \rangle$. Above, connecting edge can be done by using suitable swap games. Then we can show that $\mathcal{R}_M^r(\text{tr}_{k;m,n}^{\mathbb{F}_{\text{tr}}(\Sigma_M^{\text{int}})}(f)) \sim \mathcal{A}$, where an isomorphism for the equivalence relation \sim is given obviously from the above construction.

C Fixpoint Operator on FinScottL_1

Here we give the definition of the fixpoint operator introduced in [16, 17], with our notation.

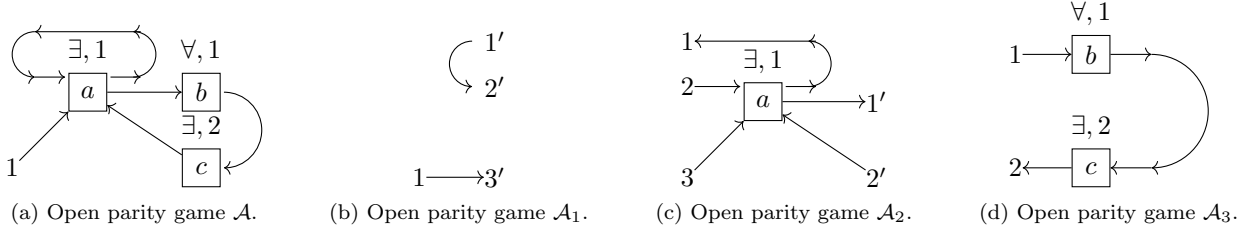


Fig. D.1. A detailed example.

Definition C.1 (semantic run-tree for fix^{GM} [16, 17]) Let $R \in \mathbf{FinScottL}_1(A + B, B)$ and $b \in |B|$; then especially, $R \subseteq P(\mathbb{N}_M \times (|A| + |B|)) \times |B|$, and, intuitively, think of b as an element of the codomain of R , rather than the domain. A *semantic run-tree* ψ (for fixpoint operator) is a possibly infinite $(\mathbb{N}_M \times (|A| + |B|))$ -labeled tree ψ that satisfies the following conditions:

- (i) The label of the root of ψ is $(0, b)$.
- (ii) Any non-leaf node of ψ has the label in $\mathbb{N}_M \times |B|$.
- (iii) For any non-leaf node of ψ with label $(p, b') \in \mathbb{N}_M \times |B|$, let $X \subseteq \mathbb{N}_M \times (|A| + |B|)$ be the set of the labels of all the children of the node. Then $(X, b') \in R$.
- (iv) For any leaf node of ψ such that its label belongs to $\mathbb{N}_M \times |B|$ (rather than $\mathbb{N}_M \times |A|$) and is (p, b') , we have $(\emptyset, b') \in R$.

We write $\text{SRT}^{\text{fix}}(A, B, R, b)$ for the set of semantic run-trees for fixpoint operator with respect to A, B, R and b . For a semantic run-tree $\psi \in \text{SRT}^{\text{fix}}(A, B, R, b)$, we define $\text{leaves}^{\text{fix}}(\psi) \in |\mathcal{P}(\square_M(A))| = P(\mathbb{N}_M \times |A|)$ as the set of elements $(p, a) \in \mathbb{N}_M \times |A|$ such that there exists a leaf ℓ of ψ such that: (i) the label of leaf ℓ is (p', a) for some $p' \in \mathbb{N}_M$ and (ii) p is the maximal priority encountered on the path from the leaf ℓ to the root of ψ .

Definition C.2 (fixpoint operator fix^{GM} [16, 17]) For every $A, B \in \mathbf{FinScottL}_1$, we define a fixpoint operator $\text{fix}_{A,B}^{\text{GM}} : \mathbf{FinScottL}_1(A + B, B) \rightarrow \mathbf{FinScottL}_1(A, B)$ as follows:

$$\text{fix}_{A,B}^{\text{GM}}(R) := \{(\text{leaves}^{\text{fix}}(\psi), b) \mid \psi \in \text{SRT}^{\text{fix}}(A, B, R, b) \text{ and every infinite path on } \psi \text{ meets the parity condition}\}$$

where an infinite sequence $((p_i, b_i))_{i \in \mathbb{N}}$ of labels meets the *parity condition* if the maximal priority that infinitely occurs in the sequence is even (i.e., $\max\{q \mid \#\{i \mid p_i = q\} = \infty\}$ is even).

D A Detailed Example

Here we demonstrate the use of our categorical theory by exhibiting the compositional solution of an open parity game. Such compositionality is enabled by the fact that the functors \mathcal{R}_M , $\llbracket - \rrbracket_M$, and \mathcal{W}_M preserve suitable structures.

Consider the open parity game \mathcal{A} shown in Fig. D.1(a), where the open ends are labelled using prop-style ordering. In reality, the numbering differs in $\text{Int}(\mathbf{FinScottL}_{!M}^{\text{op}})$ (because the Int construction numbers ends differently from props), but we make sure to refer to ends by their labels in Fig. D.1 throughout the explanation.

Our goal is to check whether the entry position 1 is a winning position or not in \mathcal{A} . We do so compositionally, i.e. by solving \mathcal{A} 's subgames and propagating those solutions. Concretely, \mathcal{A} is divided into \mathcal{A}_1 , \mathcal{A}_2 , and \mathcal{A}_3 shown in Fig. D.1(b), Fig. D.1(c), and Fig. D.1(d), respectively.

Let us start with \mathcal{A}_1 , which can be decomposed as $d_0 \oplus \text{id}_{\tau}$, whence $\llbracket \mathcal{A}_1 \rrbracket_M = d_{\llbracket 0 \rrbracket_M} \oplus \text{id}_{\llbracket \tau \rrbracket_M}$ by compact closedness of $\llbracket - \rrbracket_M$. When translated back to the level of $\mathbf{FinScottL}_{!M}^{\text{op}}$, the unit d_0 in $\text{Int}(\mathbf{FinScottL}_{!M}^{\text{op}})$ is a morphism from $\{1'\}$ to $\{2'\}$, defined as the identity (up to isomorphism here, since we changed the names to reflect those in Fig. D.1). By definition of Kleisli categories, this identity is $\epsilon_{\llbracket 1' \rrbracket_M}^{\llbracket 1' \rrbracket_M} = \{(T, 1') \mid (0, 2') \in T\}$. Similarly $\text{id}_{\llbracket \tau \rrbracket_M} = \{(T, 1) \mid (0, 3') \in T\}$. Therefore, the interpretation $\llbracket \mathcal{A}_1 \rrbracket_M$ is the following:

$$\llbracket \mathcal{A}_1 \rrbracket_M = \{(T, 1) \mid (0, 3') \in T\} \cup \{(T, 1') \mid (0, 2') \in T\}.$$

While \mathcal{A}_2 is a slightly more complex game than \mathcal{A}_1 (involving a generator, a parallel composition, and a sequential composition), because the counit is basically an identity in the Int construction, the parallel and

sequential compositions amount to bureaucratic index tracking, and we get that the interpretation is that of the generator:

$$\llbracket \mathcal{A}_2 \rrbracket_M = \{(T, i) \mid i \in \{2, 3, 2'\}, \exists j \in \{1, 1'\}, (1, j) \in T\}.$$

The last subgame, \mathcal{A}_3 , is defined as the composition of two generators (up to a unit). Here, the computation involves a non-trivial composition in $\mathbf{FinScottL}_{!M}^{op}$, for which we need a direct and explicit definition of $\delta^{!M}$, which is a bit more involved and can be found in [16], so we skip here. Intuitively, $\delta^{!M}$ takes care of both the non-determinism and of registering the highest priority seen along a path. In our case, we get $\llbracket \mathcal{A}_3 \rrbracket_M = \{(T, 1') \mid (2, 2) \in T\} \circ \{(T, 1) \mid (1, 1') \in T\}$ (for some fictious position $1'$ between the two generators), so:

$$\llbracket \mathcal{A}_3 \rrbracket_M = \{(T, 1) \mid (2, 2) \in T\}.$$

Next, we turn our attention to the first sequential composition that is not trivial in $\text{Int}(\mathbf{FinScottL}_{!M}^{op})$. Indeed, up to this point, none of the computations required to compute a trace (because the trace was over the tensor unit and therefore trivial), but the interpretation of $\mathcal{A}_2 \circ \mathcal{A}_1$ will. By definition, if $f: (X_+, X_-) \rightarrow (Y_+, Y_-)$ and $g: (Y_+, Y_-) \rightarrow (Z_+, Z_-)$ in $\text{Int}(\mathbf{FinScottL}_{!M}^{op})$, then $g \circ f$ is computed (in $\mathbf{FinScottL}_{!M}^{op}$) as the trace over Y_- of:

$$((\sigma_{Z_+, Y_-} \circ g) \oplus \text{id}_{X_-}) \circ (\text{id}_{Y_+} \oplus \sigma_{X_-, Z_-}) \circ ((f \circ \sigma_{Y_-, X_+}) \oplus \text{id}_{Z_-}).$$

Here, however, $X_- = I$, so $\sigma_{X_-, Z_-} = \text{id}_{Z_-}$, $\llbracket \mathcal{A}_1 \rrbracket_M \circ \sigma_{Y_-, X_+} = \text{id}_{Y_-, X_+}$ (\mathcal{A}_1 is a swap), and $\sigma_{Z_+, Y_-} \circ \llbracket \mathcal{A}_2 \rrbracket_M = \llbracket \mathcal{A}_2 \rrbracket_M$ (\mathcal{A}_2 is “symmetric” in 1 and $1'$), we get that we need to compute the trace of $\llbracket \mathcal{A}_2 \rrbracket_M$.

To this end, we compute the set of semantic run-trees ψ that correspond to it. Here, $D = \{2\} = \{1\}$ (we need to resort to such notation because we kept the names from Fig. D.1, but this does give the right intuition that 2 and 1 get “connected” by the trace in \mathcal{A}_2), $A = \{3, 2'\}$, $B = \{1'\}$, and $R = \llbracket \mathcal{A}_2 \rrbracket_M$.

The root of ψ must be $(0, 3)$ or $(0, 2')$ by Def. 4.6(i). By Def. 4.6(iv) and the value of $\llbracket \mathcal{A}_2 \rrbracket_M$, only elements of the form $(p, 1')$ can be leaves, and in particular, the root cannot be a leaf. By Def. 4.6(iii), all non-leaves must have either $(1, 1)$ or $(1, 1')$ as one of its children. From this, we can deduce that, either ψ has $(1, 1')$ as one of its leaves, or it has an infinite path with only parity 1, and therefore does not meet the parity condition. This corresponds to the fact that, if \exists wants to win, then they cannot choose to go through the cycle infinitely often and have to go to $1'$ at some point. Moreover, any ψ without leaves fails to meet the parity condition, so (after renaming 3 to 1 because of the composition)

$$\llbracket \mathcal{A}_2 \circ \mathcal{A}_1 \rrbracket_M = \{(T, i) \mid i \in \{1, 2'\}, (1, 1') \in T\}.$$

The computation follows the same step for the final composition. We need to take the trace of some composite of $\llbracket \mathcal{A}_2 \circ \mathcal{A}_1 \rrbracket_M$, $\llbracket \mathcal{A}_2 \rrbracket_M$, some swaps and identities, namely $R = \{(T, i) \mid i \in \{1, 2\}, (2, 2) \in T\}$. Here, $D = \{2\} = \{2'\}$, $A = \{1\}$, and $B = \emptyset$. Contrary to the example above, however, there is a semantic run-tree with no leaves that meets the parity condition: for example the tree whose root is $(0, 1)$, and whose nodes all have $(2, 2)$ as their unique child (note that this does verify condition (iii) of Def. 4.6 because $2 = 2'$). This corresponds to the fact that this new cycle is winning for \exists . Finally, we get

$$\llbracket \mathcal{A} \rrbracket_M = \{(T, 1) \mid \text{true}\} = \{(\emptyset, 1)\},$$

which means that 1 is winning in \mathcal{A} .