### Completeness results for the zx-calculus

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Approximate completeness for single qubits

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Completeness for stabilizer diagrams

**Conclusions & Outlook** 

- zx-calculus is incomplete, even for single qubits; not obvious how to complete it
- instead: look for fragments of the general calculus which are complete
- approximate universality: small sets of operators suffice to approximate arbitrary unitaries to any accuracy
- approximate completeness: completeness for such a set

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Completeness for stabilizer diagrams

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## The single-qubit Clifford+T group

An approximately universal group, generated by:

• single-qubit Clifford group  $C_1 = \langle S, H \rangle$ , where

$$S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix} \qquad \qquad \phi \pi/2$$
$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \qquad \qquad \blacksquare$$
$$T \text{ gate} \qquad \qquad T = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix} \qquad \qquad \phi \pi/4$$

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The zx-calculus for the single-qubit Clifford+T group

Generated by  $\phi \pi/2$ ,  $\mathbf{H}$ , and  $\phi \pi/4$  —or  $\phi \pi/4$  and  $\phi \pi/2$ 

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single-qubit: restrict diagrams to line graphs

• *Clifford+T group*: restrict phases to multiples of  $\pi/4$  (Ignore global phases.)

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single-qubit: restrict diagrams to line graphs

 Clifford+T group: restrict phases to multiples of π/4 (Ignore global phases.)
 Rules:



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and  $\oint 2n\pi = |$  for integer *n* 

# Single-qubit Clifford diagrams

Any single-qubit Clifford diagram can be written uniquely as one of the following, with  $\alpha, \beta, \gamma \in \{0, \pi/2, \pi, -\pi/2\}$ :



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## Single-qubit Clifford diagrams

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or, equivalently, one of



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where  $a, b \in \{0, 1\}$  and  $\beta, \gamma$  as above.

### The normal form for single-qubit Clifford+T diagrams

Following [Matsumoto & Amano 2008], any single-qubit Clifford+T diagram is either pure Clifford or it can be written as

$$\begin{array}{c}
\overbrace{V_{1}}{V_{n}} \\
\overbrace{V_{1}}{V_{1}} \\
\overbrace{V}{V_{1}} \\
\overbrace{V}{V} \\
\overbrace{V}{V_{1}} \\
\overbrace{V}$$

with *n* a non-negative integer and  $\alpha, \beta, \gamma \in \{0, \pi/2, \pi, -\pi/2\}$ .

Diagrams are rewritten into normal form by pushing phases towards the bottom of the diagram as much as possible.

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- can push  $\phi \alpha$  and  $\phi \pi$  past  $\phi \pi/4$
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- can push  $\phi \pi$  and  $\phi \pi$  past  $\bigvee \in \left\{ \begin{array}{c} \phi \pi/4 \\ \phi \pi/2 \end{array} \right\}$

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# Rewriting diagrams into normal form, starting with a Clifford diagram

The single-qubit Clifford+T group is generated by  $\oint \pi/4$  and  $\oint \pi/2$  so it suffices to check what happens when a pure Clifford operator or a normal form diagram is composed with one of these.

Now for any Clifford unitary C,

$$\begin{array}{c} \bullet \pi/2 \\ \hline C \\ \hline \end{array} \text{ is pure Clifford, and } \bullet \pi/4 \\ \hline C \\ \hline \end{array} = \underbrace{U}_{-} \in \left\{ \begin{array}{c} \bullet \pi/4 + \alpha \\ \bullet \pi/2 \\ \bullet \pi/2 \\ \bullet \pi/2 \end{array} \right\}.$$

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# Rewriting diagrams into normal form, starting with a normal form diagram

For any 
$$\underline{W} \in \left\{ \left| \begin{array}{c} \bullet \pi/2 \\ \bullet \pi/2 \\ \hline \pi/2 \\ \hline \pi/2 \\ \hline \end{array} \right\}$$
 and  $\underline{V}_{n} \in \left\{ \begin{array}{c} \bullet \pi/4 \\ \bullet \pi/2 \\ \hline \bullet \pi/2 \\ \hline \end{array} \right\},$ 
$$\begin{pmatrix} \bullet \pi/2 \\ \bullet \pi/2 \\ \hline \hline V_{n} \\ \hline \hline V_{n} \\ \hline \end{array} \right\} = \begin{array}{c} \underline{C} \\ \hline V_{n} \\ \hline \hline V_{n} \\ \bullet b\pi \\ \hline \end{array}$$

and

$$\stackrel{\blacklozenge}{\underbrace{W}}_{V_n}^{\pi/4} = \left\{ \begin{array}{c} \stackrel{\bullet}{\underbrace{C'}} & \text{or} & \begin{array}{c} \stackrel{\bullet}{\underbrace{V_{n+1}}} \\ \stackrel{\bullet}{\underbrace{V_n}} \end{array} \right\}$$

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$$\begin{pmatrix} \bullet \pi/2 \\ \bullet \pi/2 \\ \hline \hline V_{n} \\ \hline \hline V_{n} \\ \hline \end{array} \right\} = \begin{array}{c} \underline{C} \\ \hline V_{n} \\ \hline \hline V_{n} \\ \bullet b\pi \\ \hline \end{array}$$

and

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Case n = 0 is similar.

# No normal form diagram represents the identity operator

Can write any single-qubit density operator as xX + yY + zZ, where  $x, y, z \in \mathbb{R}$  and X, Y, Z are the Pauli matrices. Clifford unitaries act on the vectors (x, y, z) by permuting the elements and adding minus signs; T sends

$$(x,y,z)\mapsto \frac{1}{\sqrt{2}}\left(x-y,x+y,z\sqrt{2}\right).$$

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States resulting from application of a normal form operator to  $|0\rangle$  will have vectors of the form

$$\frac{1}{\sqrt{2^m}} \left( x_1 + x_2 \sqrt{2}, y_1 + y_2 \sqrt{2}, z_1 + z_2 \sqrt{2} \right)$$

where  $x_1, x_2, y_1, y_2, z_1, z_2 \in \mathbb{Z}$ . By parity arguments, can show none of them represent  $|0\rangle$ , thus no normal form operator is equal to the identity.

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We have 
$$U^{\dagger} = \begin{pmatrix} C \\ \phi \pi/4 \end{pmatrix}$$
 and  $V^{\dagger} = V \\ \phi \pi/2 \end{pmatrix}$ , so

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We have 
$$U^{\dagger}_{\downarrow} = \dot{C}_{\phi \pi/4}$$
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 and  $V^{\dagger} = V$ , so  $\phi \pi/2$ , so



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$$U^{\dagger} = \overset{\frown}{\overset{\frown}{\bullet}} \pi/4$$
 and  $\overset{\frown}{V_{\uparrow}} = \overset{\frown}{\overset{\lor}{\bullet}} \pi/2$ , so  
 $\overset{\downarrow}{\overset{\downarrow}{\bullet}} \pi/2$   
 $\overset{\downarrow}{\overset{\downarrow}{\bullet}} \pi/2$ 

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$$= \frac{V_0}{V_{n-1}}$$

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Given two normal form diagrams that are not identical, composing one diagram with the inverse of the other will yield a non-trivial diagram.

- Assume diagrams are such that the topmost nodes differ (in colour or phase, or both).
- Inverting a diagram does not change the number of π/4 phases.
- Distinguish cases to see that the resulting diagram never collapses to the trivial one.



Approximate completeness for single qubits

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Completeness for stabilizer diagrams

**Conclusions & Outlook** 

### Stabilizer quantum mechanics

Stabilizer operations:

- preparation of qubits in state  $|0\rangle$
- Clifford unitaries, generated by

$$S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}, \quad H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}, \quad \Lambda X = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

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measurements in computational basis

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measurements in computational basis
 ZX-calculus:

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where  $\alpha,\beta$  are multiples of  $\pi/{\rm 2}$ 

### Graph states in the zx-calculus

#### Definition

Let *G* be a finite simple undirected graph. The zx-calculus diagram for the corresponding graph state consists of:

- for each node in G, a green node with one output, and
- ▶ for each edge in *G*, an edge with a Hadamard node on it.

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Theorem (Van den Nest et. al, 2004)

Any stabilizer state is local Clifford-equivalent to some graph state.

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A local complementation about a vertex v inverts the subgraph generated by the neighbourhood of v: e.g.

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## Stabilizer completeness proof (overview)

[Duncan & Perdrix, 2009] show that local complementations can be derived from the rules of the zx-calculus. Use this to show that the results from [Van den Nest et al., 2004] hold in the zx-calculus:

Any stabilizer zx-calculus diagram can be rewritten into a graph state with local Clifford operators, e.g.



There exists a terminating algorithm that, given two stabilizer diagrams, will rewrite them to be identical if possible.

See arXiv:1307.7025 for details.

### **Conclusions & Outlook**

- ZX-calculus is not complete in general but fragments of it are complete, e.g.
  - ► line graphs where all phases are multiples of π/4 (single-qubit Clifford+T group)
  - diagrams where all phases are multiples of π/2 (stabilizer quantum mechanics)

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- can these results be combined to multi-qubit Clifford+T operators?
- can we introduce some notion of approximate equality in the zx-calculus?

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Thank you!