

# Mixed quantum states in higher categories

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(with Chris Heunen and Jamie Vicary)

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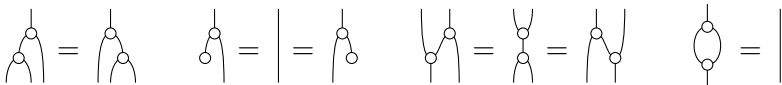
- The theory of bimodules
- The  $2(-)$  construction
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## Applications

- A unified description of teleportation and classical encryption
- A unified security proof

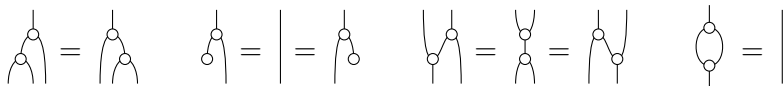
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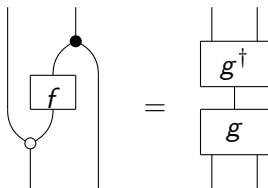


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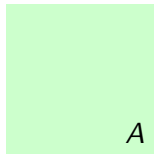
2. Completely positive maps between Frobenius algebras:  
morphisms  $f$  in  $\mathbf{C}$ , for which  $\exists g$  such that



## 2-categories and their graphical language

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*0-cells*   Regions   Classical information



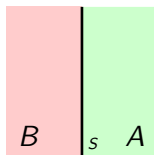
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$$B \xrightarrow{s} A$$

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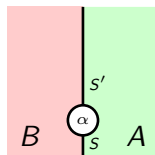
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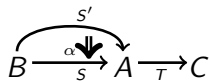
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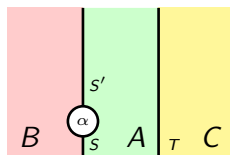
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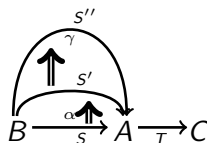
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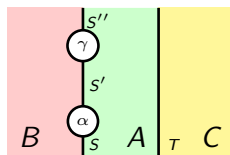
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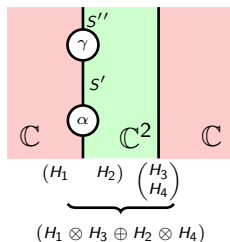
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The standard example is **2Hilb**:

- ▶ *0-cells* given by natural numbers
- ▶ *1-cells* given by matrices of finite-dimensional Hilbert spaces
- ▶ *2-cells* given by matrices of linear maps

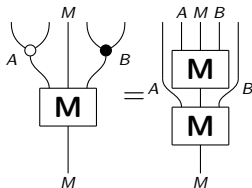
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Let  $(A, \circlearrowleft, \circlearrowright)$  and  $(B, \bullet\circlearrowleft, \bullet\circlearrowright)$  be classical structures in  $\mathbf{C}$ .

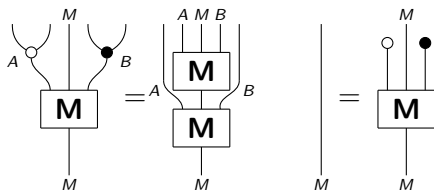
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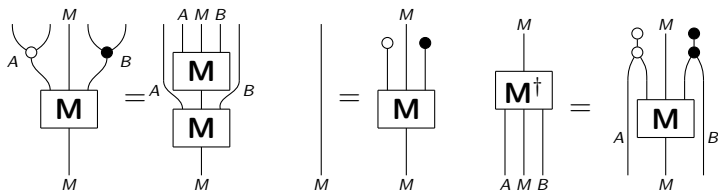
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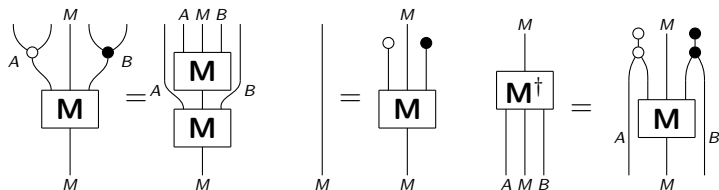
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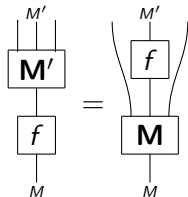
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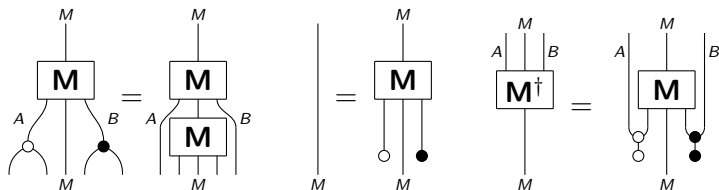
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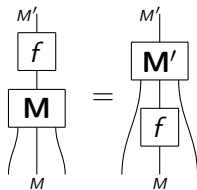
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For proofs see LW (2013), Masters's thesis, 'Categorical Models for Quantum Computing'.

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Horizontal composition is defined by the following coequaliser in  $\mathbf{C}$ :

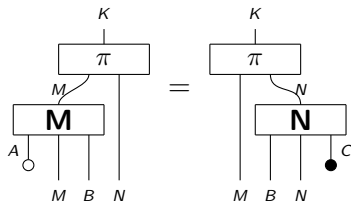
$$\begin{array}{ccccc} M \otimes B \otimes N & \xrightarrow{\mathbf{M}_B \otimes id_N} & M \otimes N & \xrightarrow{\pi} & M \otimes_B N \\ & \xrightarrow{id_M \otimes_B \mathbf{N}} & & & \downarrow \tilde{f} \\ & & & & K \end{array}$$

The diagram shows a coequaliser in the 2-category  $\mathbf{C}$ . The top row consists of three objects:  $M \otimes B \otimes N$ ,  $M \otimes N$ , and  $M \otimes_B N$ . The first arrow is labeled  $\mathbf{M}_B \otimes id_N$  and the second is labeled  $\pi$ . The bottom row consists of two objects:  $M \otimes N$  and  $K$ . The arrow from  $M \otimes N$  to  $K$  is labeled  $f$ . A dashed arrow labeled  $\tilde{f}$  points from  $M \otimes_B N$  to  $K$ .

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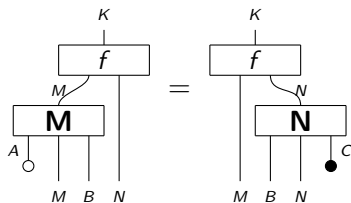
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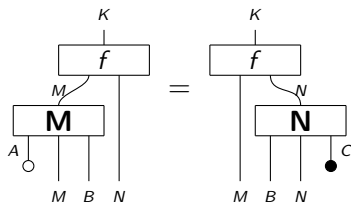
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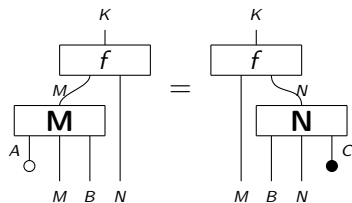


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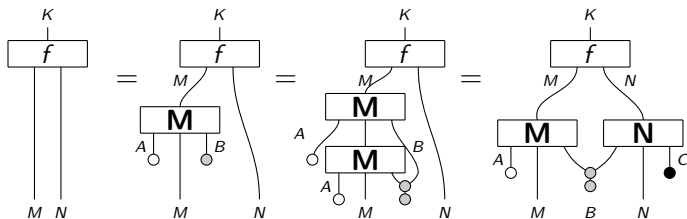


Can we find this module explicitly? Yes!



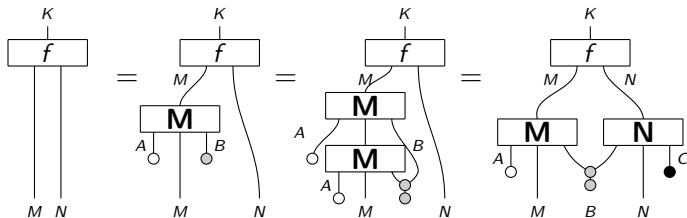
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Any such  $f$  factorizes through  $\mathbf{M} \circ \mathbf{N}$ :



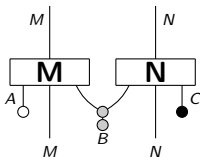
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## Theorem

*Finding the dagger coequaliser is equivalent to finding a dagger splitting of the following morphism:*



## $2(\mathbf{CP}^*(-))$

We would like to understand the 2-category ' ? '

$$\begin{array}{ccc} \mathbf{FHilb} & \xrightarrow{\mathbf{CP}^*(-)} & \mathbf{CP}^*(\mathbf{FHilb}) \\ \downarrow 2(-) & & \downarrow 2(-) \\ 2(\mathbf{FHilb}) & \overset{\text{---}}{\dashrightarrow} & ? \\ & \mathbf{CP}^*(-) & \end{array}$$

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- ▶ There is a correspondence between special dagger Frobenius algebras on classical structures in  $\mathbf{FHilb}$  and finite groupoids.
- ▶  $\mathbf{CP}^*(\mathbf{FHilb})$  does not have all coequalisers.

## Modelling POVM's

The following subcategory of  $2(\text{CP}^*(\mathbf{FHilb}))$  is a sufficient model for modelling communication protocols:

- ▶ 0-cells: natural numbers
- ▶ 1-cells: matrices of dagger Frobenius algebras
- ▶ 2-cells: matrices of completely positive maps

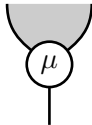


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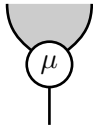


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### Theorem

*Measurements on algebras  $\mathbb{C}^n$  are exactly stochastic maps.  
Measurements on algebras  $B(H)$  are exactly POVMs.*

# Modelling POVM's

Proof.

The count preserving condition gives us

$$\left( \begin{array}{c} \text{C}^n \\ \mu \\ \text{---} \end{array} = \begin{array}{c} \circ \\ \text{---} \end{array} \right) \Leftrightarrow \left( \begin{array}{c} \text{---} \\ \mu^\dagger \\ \text{C}^n \\ \text{---} \\ \circ \end{array} = \begin{array}{c} \text{---} \\ \circ \end{array} \right)$$

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So we have the following equalities of positive elements:

$$\sum_{i=1}^n \begin{array}{c} \mu_i^\dagger \\ \text{---} \end{array} = \begin{array}{c} \mu^\dagger \\ \text{C}^n \\ \text{---} \end{array} = \begin{array}{c} \text{---} \\ \circ \end{array}$$

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$$\left( \begin{array}{c} \text{C}^n \\ \mu \\ \text{---} \end{array} = \begin{array}{c} \circ \\ \text{---} \end{array} \right) \stackrel{\dagger}{=} \left( \begin{array}{c} \mu^\dagger \\ \text{C}^n \\ \text{---} \end{array} = \begin{array}{c} \text{---} \\ \circ \end{array} \right)$$

So we have the following equalities of positive elements:

$$\sum_{i=1}^n \begin{array}{c} \text{---} \\ \mu_i^\dagger \\ \circ \end{array} = \begin{array}{c} \text{---} \\ \mu^\dagger \\ \text{C}^n \\ \text{---} \end{array} = \begin{array}{c} \text{---} \\ \circ \end{array}$$

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# Modelling POVM's

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- ▶ On  $\mathbb{C}^n$  this corresponds to a stochastic map
- ▶ On  $B(\mathbb{C}^n)$  this corresponds to a POVM



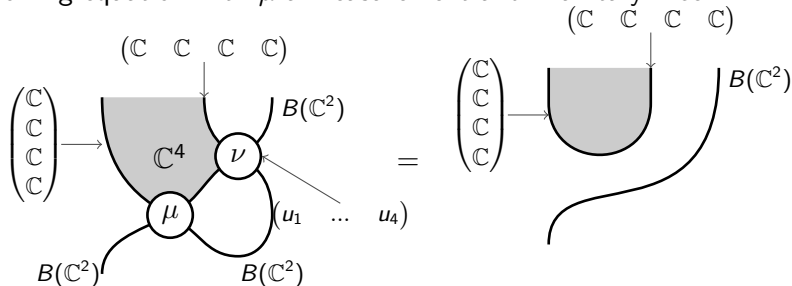
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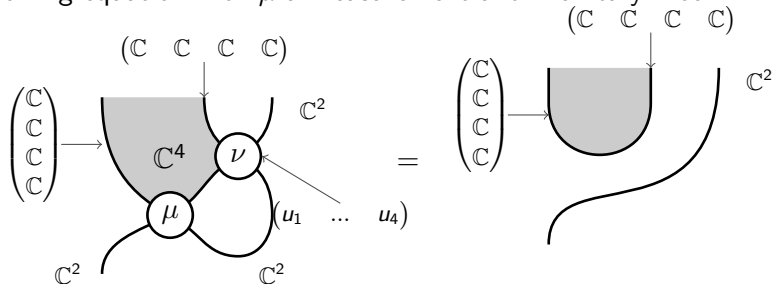
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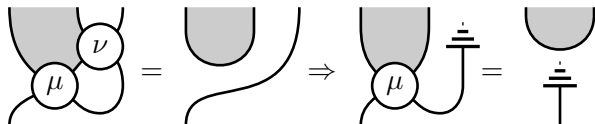


This equation corresponds to:

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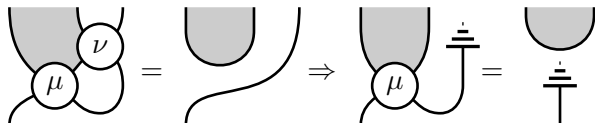
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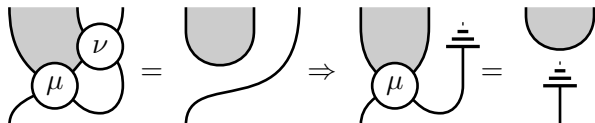
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- ▶ We apply the trace map on both sides of the equation
- ▶ On the left-hand-side:  $\nu$  is a family invertible completely positive maps, which are trace preserving.

So this give a unified security proof

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Thank you!