A Kochen-Specker system has at least 21 vertices

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A Kochen-Specker system S is a finite set of points on the (open) northern hemisphere, such that there is no 010-coloring; that is: there is no $\{0,1\}$ -valued coloring with

- 1. three pairwise orthogonal points are assigned (1,0,0), (0,1,0) or (0,0,1) and
- 2. two orthogonal points are not assigned (1,1).

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Theorem (Kochen-Specker)

There is a Kochen-Specker system. Thus: there is no non-contextual deterministic theory predicting the measurement of a SPIN-1 particle.

Arends, Wampler, Ouaknine 2009

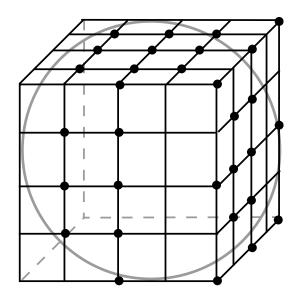
Arends, Wampler, Ouaknine 2009

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Kochen-Specker 1975 \leq 117 Penrose, Peres (indep.) 1991 \leq 33 Conway \sim 1995 \leq 31 U&W may \geq 21 Arends, Wampler, Ouaknine 2009 \geq 18
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Conway's record



Given a finite set of points S on the projective plane, its orthogonality graph $\mathcal{G}(S)$ has as vertices the points and two points are adjacent if and only if they are orthogonal.

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A 010-coloring of a graph, is a $\{0,1\}$ -vertex coloring, such that

- 1. every triangle is colored (1,0,0), (0,1,0) or (0,0,1) and
- 2. no adjacent vertices are colored both 1.

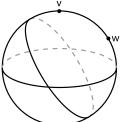
There is a Kochen-Specker system with n points if and only if there is a embeddable and non-010-colorable graph on n vertices.

Restrict the search

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(The orthogonality graph of) a minimal Kochen-Specker system is connected; \sim 10^{26.4} squarefree and \sim 10^{10.2} has minimal vertex degree 3; \sim 10^{7.5}
```

The candidates

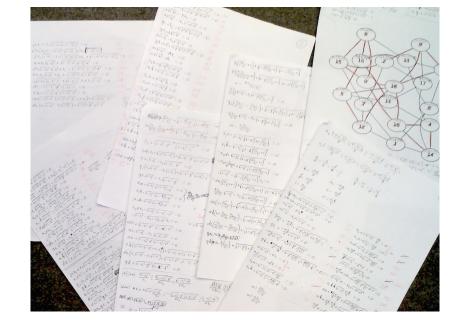
Our computation found the following number of non 010-colorable squarefree graphs with minimal vertex degree 3

#V	# candidates
≤ 16	0
17	1
18	2
19	19
20	441

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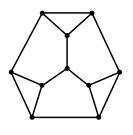
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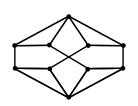
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≤ 16	0
17	1
18	2
19	19
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21	\geq 7616

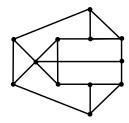


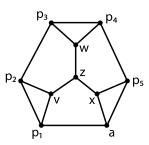
Unembeddable subgraphs

All these candidates contain as a subgraph one of these unembeddable graphs.



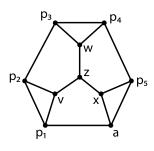






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Similarly: p_2 is collinear to $v \times (v \times a)$. And so on. We see a is collinear to $x \times (x \times (w \times (w \times (v \times (v \times a)))))$.

We may assume z = (0,0,1), x = (1,0,0), $v = (v_1, v_2, 0)$, $w = (w_1, w_2, 0)$ and $a = (0, a_2, a_3)$. We have:

$$\begin{pmatrix} 0 \\ a_2 \\ a_3 \end{pmatrix} \text{ is collinear to } \begin{pmatrix} 0 \\ -a_2v_1w_2(v_1w_1+v_2w_2) \\ -a_3(v_1^2w_1^2+v_1^2w_2^2+v_2^2w_1^2+v_2^2w_2^2) \end{pmatrix}$$

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That is:

$$v_1 w_2 \langle v, w \rangle = v_1 w_2 (v_1 w_1 + v_2 w_2)$$

$$= v_1^2 w_1^2 + v_1^2 w_2^2 + v_2^2 w_1^2 + v_2^2 w_2^2$$

$$= (v_1^2 + v_2^2) w_1^2 + (v_1^2 + v_2^2) w_2^2$$

$$= w_1^2 + w_2^2 = 1.$$

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$$= (v_1^2 + v_2^2) w_1^2 + (v_1^2 + v_2^2) w_2^2$$

$$= w_1^2 + w_2^2 = 1.$$

Since v and w are not collinear, we have by Cauchy-Schwarz $|\langle v,w\rangle|<1$. Note $|v_1|,|w_2|\leq 1$. Thus: $|v_1w_2\langle v,w\rangle|<1$. Contradiction.

Example of automized cross product chasing

```
load_package redlog;
rlset R;
procedure d(x,v):
    (first x) * (first y) +
    (second x) * (second v) +
    (third x) * (third v);
procedure k(x,v):
    {(second x)*(third y) - (third x)*(second y),
     (third x)*(first v) - (first x)*(third v).
     (first x)*(second v) - (second x)*(first v)}:
v0c1 := 1; v0c2 := 0; v0c3 := 0;
v1c1 := 0; v1c2 := 1; v1c3 := 0;
v0 := {v0c1, v0c2, v0c3}:
v1 := {v1c1, v1c2, v1c3}:
v2 := {v2c1, v2c2, v2c3};
v3 := {v3c1, v3c2, v3c3};
v2c1 := 0:
neg0 := k(v0,k(v3,v1)):
                                                        (snip)
neq29 := k(k(k(k(v3,v1),v1),v2),k(k(v3,v0),v3));
phi :=
       (first neg0 neg 0 or
        second neg0 neg 0 or
        third neg0 neg 0) and
                                                         (snip)
       (first neg29 neg 0 or
        second neg29 neg 0 or
        third neq29 neq 0) and
       d(v2.v0) = 0 and
      d(k(k(v3,v0),v3),k(k(k(k(v3,v1),v1),v2),v2)) = 0 and
        true;
rlae ex(v3c3.
     ex(v3c2.
     ex(v3c1,
     ex(v2c3.
     ex(v2c2,phi))));
```

Source code, paper and experimental results can be found at kochen-specker.info

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Some open problems:

- ▶ If G is embeddable, is there a S such that G = G(S).
- ▶ Is every embeddable graph, grid embeddable? That is: using points of the form $\left(\frac{x}{\sqrt{n}}, \frac{y}{\sqrt{n}}, \frac{z}{\sqrt{n}}\right)$ for $x, y, z, n \in \mathbb{Z}$.