

Mathematical Semantics of Computer Systems, *MSCS* (4810-1168) Handout for Lecture 3 (2016/10/17)

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(ask me for username, password)

We loosely follow [4], but it hardly serves as an introductory textbook. More beginner-friendly ones include [1, 5]; other classical textbooks include [6, 2]. nLab (ncatlab.org) is an excellent online information source.

1 Today's Goal, I

Come to understand and prove the following statement.

Proposition 1. *Let $X, Y \in \mathbb{C}$. A product of X and Y , if it exists, is unique up-to a canonical isomorphism.*

2 Today's Agenda I

2.1 Categories, Functors

Definition. Functor

Example. Monotone functions as functors (for preorders considered as categories); monoid homomorphisms as functors (for monoids).

Example. Graphs as functors. Monoid/group actions as functors.

2.2 Reasoning with Arrows

Definition. Epi, mono. Split epi, split mono.

Definition. Coproduct.

3 Today's Goal II

Identify the following framework of *abstract interpretation* [3] as an instance of adjunction. (Thanks are due to Kengo Kido for a nice introduction.)

Definition (Galois connection). Let L and \bar{L} be posets; and $\alpha: L \rightarrow \bar{L}$ and $\gamma: \bar{L} \rightarrow L$ be monotone functions. The pair (α, γ) is said to be a *Galois connection* if, for any $x \in L$ and $\bar{x} \in \bar{L}$,

$$\alpha(x) \leq_{\bar{L}} \bar{x} \quad \text{if and only if} \quad x \leq_L \gamma(\bar{x}) .$$

Example (interval domain). Let

$$L := \mathcal{P}(\mathbb{N}) \quad \text{and} \quad \bar{L} := \{\emptyset\} \cup \{[l, r] \mid l, r \in \mathbb{N} \cup \{-\infty, \infty\}, l \leq r\}$$

where each set is ordered by inclusion. Moreover,

$$\alpha(S) := [\min S, \max X] \quad \text{and} \quad \gamma(\bar{S}) := \{n \in \mathbb{N} \mid n \in \bar{S}\} .$$

Then the pair (α, γ) is a Galois connection.

4 Today's Agenda II

4.1 Natural Transformations

Definition. Natural transformation

Example. Natural transformations in graphs, and in monoid/group actions.
Natural transformations between monotone maps as functors.

Definition. Horizontal and vertical composition of natural transformation

4.2 Limits and Colimit

Definition. Diagram, cone, cocone

Definition. Limit, colimit

Proposition 2. *Limits from products and equalizers*

Corollary 1. *Concrete presentation of (co)limits in Sets*

4.3 Adjunction

Definition. Homset.

Definition. Adjunction.

Example. Free monoids.

Definition. Unit, counit.

Lemma 1. *Adjoint transposes by units and counits.*

Proposition 3. *Characterization of adjunction by: 1) the universality of η (Def. 3.2 of [Lambek & Scott], intuitive for free monoids); 2) the triangular equalities (Def. 3.1 of [Lambek & Scott]).*

Lemma 2. 1. *Adjoint functors determine each other uniquely up-to canonical natural isomorphisms.*

2. *Composition of adjoints.*

4.4 Limits as Adjoints

Definition. Functor category

Proposition 4. *Limits give rise to an adjunction.*

5 Exercises

1. Formulate and prove the following statement.

A right adjoint preserves limits.

2. Prove the following: in an adjunction $F \dashv G$, G is faithful if and only if every component of the counit ε is an epi. [6, Thm. IV.3.1]

6 Report Problems

1. Prove that: in the category **Sets** of sets and functions, an arrow is a mono if and only if it is an injective function. Similarly, an arrow is an epi if and only if it is a surjective function.
2. Give a detailed proof of the *Today's Goal, I*:

Let $X, Y \in \mathbb{C}$. A product of X and Y , if it exists, is unique up-to a canonical commuting isomorphism.

Reports are due at the beginning of the next lecture.

References

- [1] S. Awodey. *Category Theory*. Oxford Logic Guides. Oxford Univ. Press, 2006.
- [2] M. Barr and C. Wells. *Toposes, Triples and Theories*. Springer, Berlin, 1985. Available online.
- [3] P. Cousot and R. Cousot. Abstract interpretation: A unified lattice model for static analysis of programs by construction or approximation of fixpoints. In R.M. Graham, M.A. Harrison and R. Sethi, editors, *Conference Record of the Fourth ACM Symposium on Principles of Programming Languages, Los Angeles, California, USA, January 1977*, pp. 238–252. ACM, 1977.
- [4] J. Lambek and P.J. Scott. *Introduction to higher order Categorical Logic*. No. 7 in Cambridge Studies in Advanced Mathematics. Cambridge Univ. Press, 1986.
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- [6] S. Mac Lane. *Categories for the Working Mathematician*. Springer, Berlin, 2nd edn., 1998.