

Report Assignments

(Due: the beginning of the lect.)
Mon. 2016.11.7

1a) In the last proof, in the construction of $\eta: \text{id} \Rightarrow GF$

$$\varepsilon: FG \Rightarrow \text{id}$$

from $\left(C(x, GA) \xrightarrow[\varphi_{x,A}]{\cong} D(Fx, A) \right)_{x \in C, A \in A}$

prove that the family

$$\left(\eta_x: x \longrightarrow GFX \stackrel{\mathbb{C}}{\longleftarrow} \right)_{x \in C}$$

indeed satisfies the naturality requirement, and hence constitutes a natural transformation.

Proof. We have to show the following commutativity.

$$\begin{array}{ccc}
 \mathcal{C} \downarrow & X & \xrightarrow{\eta_X} & GFX \\
 & \downarrow F & & \downarrow GFF \\
 \mathcal{C} \downarrow & X & \xrightarrow{\eta_X} & GFX \\
 & \downarrow F & \text{aim} & \downarrow GFF \\
 \mathcal{C} \downarrow & X' & \xrightarrow{\eta_{X'}} & GFX'
 \end{array}$$

Towards this goal we use the following naturality of the bijective correspondence.

$$\begin{array}{ccc}
 \mathcal{C} \downarrow & \mathcal{D} \downarrow & \text{Sets} \\
 X & FX & \mathcal{C}(X, GFX) \xrightarrow{\cong} \mathcal{D}(FX, FX) \\
 \uparrow \text{id}_X & \downarrow FF & \downarrow \mathcal{C}(\text{id}_X, GFF) \text{ // naturality of } \varphi \downarrow \mathcal{D}(\text{id}_{FX}, FF) \\
 X & FX' & \mathcal{C}(X, GFX') \xrightarrow{\cong} \mathcal{D}(FX, FX') \\
 & & \varphi_{X, FX'}
 \end{array}$$

\Downarrow

Therefore

$$\begin{array}{ccc}
 \text{Sets} & & \varphi_{X, FX}^{-1}(\text{id}_{FX}) \xleftarrow{\cong} \text{id}_{FX} \\
 \mathcal{C}(X, GFX) & \xleftarrow{\varphi_{X, FX}^{-1}} & \mathcal{D}(FX, FX) \\
 \downarrow \eta_X & & \downarrow \mathcal{C}(\text{id}_X, GFF) \\
 \varphi_{X, FX}^{-1}(\text{id}_{FX}) & \xleftarrow{\eta_X} & (\text{id}_{FX} \downarrow) \\
 & & \text{by def. of } \eta \\
 & & (GFF) \circ \eta_X \\
 & & \parallel \\
 & & \varphi_{X, FX'}^{-1}(FF) \xleftarrow{\varphi^{-1}} FF \\
 & & \downarrow \mathcal{D}(\text{id}, FF) \\
 & & FF
 \end{array}$$

Thus we've shown

$$(GFf) \circ \eta_x = \varphi_{x, Fx'}^{-1} (Ff). \quad (1)$$

Here we invoke on another naturality square:

$$\begin{array}{ccc}
 \text{C} \downarrow & \text{D} \downarrow & \text{Sets} \\
 X' & FX' & \text{C}(X', GFX') \xrightarrow[\varphi]{\cong} \text{D}(FX', FX') \\
 \uparrow f & \downarrow \text{id}_{FX'} & \downarrow \text{C}(f, \text{id}) \qquad \downarrow \text{C}(Ff, \text{id}) \\
 X & FX' & \text{C}(X, GFX') \xrightarrow[\varphi_{x, FX'}]{\cong} \text{D}(FX, FX')
 \end{array}$$

Hence

$$\begin{array}{ccc}
 \eta_{X'} & \xleftarrow[\cong]{\varphi^{-1}} & \text{id}_{FX'} \\
 \downarrow \text{C}(f, \text{id}) & & \downarrow \text{C}(Ff, \text{id}) \\
 \eta_{X'} \circ f & & \\
 \parallel & & \\
 \varphi_{x, FX'}^{-1} (Ff) & \xleftarrow[\cong]{\varphi^{-1}} & Ff
 \end{array}$$

Thus we have

$$\eta_{X'} \circ f = \varphi_{x, FX'}^{-1} (Ff) \quad (2)$$

Combining ①, ② yield

$$(GFf) \circ \eta_x = \eta_{x'} \circ f$$

that is the desired commutativity. \square

1b) Show that the same η, ε as above satisfy the equational axioms in the def. of adjunction. in 1a

Proof. We shall first show

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{D} \\ \eta \downarrow & & \downarrow \varepsilon \\ \mathcal{C} & \xrightarrow{F} & \mathcal{D} \end{array} \begin{array}{l} \text{id} \\ \nearrow \\ \searrow \end{array} \begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{D} \\ \parallel & & \parallel \\ \mathcal{C} & \xrightarrow{F} & \mathcal{D} \end{array} \quad (*)$$

that is, for each $x \in \mathcal{C}$,

$$\begin{array}{ccc} FX & \xrightarrow{F\eta_x} & FGFX \\ \text{id} \searrow & \parallel \text{aim} & \downarrow \varepsilon_{FX} \\ & & FX \end{array}$$

for $\alpha: F \Rightarrow G$, a nat. trans. $\alpha_x: FX \rightarrow GX$ is its x -Component

Convince yourself that $\begin{array}{ccc} FX & \xrightarrow{F\eta} & FGFX \\ \downarrow \varepsilon_{FX} & & \downarrow \varepsilon_{FX} \end{array}$ is indeed the x -Component of the LHS of $(*)$

Since

$$\mathcal{D}(FX, FX) \xrightarrow[\cong]{\varphi_{x,FX}^{-1}} \mathcal{C}(x, GFX)$$

is an isomorphism in Sets (i.e. a bijection),

it suffices to show that

$$\varphi_{x, FX}^{-1}(\text{id}_{FX}) = \varphi_{x, FX}^{-1}(\varepsilon_{FX} \circ F\eta_x) \quad (**)$$

Now

$$(LHS) = (x \xrightarrow{\eta_x} GFX) \quad (\text{by def. of } \eta)$$

$$(RHS) = \varphi_{x, FX}^{-1} \left(FX \xrightarrow{F\eta_x} FGFX \xrightarrow{\varepsilon_{FX}} FX \right) \quad \begin{array}{l} \text{def. of } \varepsilon \\ \parallel \end{array}$$

$$\stackrel{\varphi}{=} \eta_x \quad \begin{array}{c} \mathcal{C}(x, GFX) \\ \uparrow \varphi^{-1} \\ \mathcal{D}(FX, FX) \end{array} \quad \varphi_{GFX, FX}(\text{id}_{GFX})$$

naturality of φ , that is,

$$\begin{array}{ccccc} \mathcal{C} & \text{SETS} & & & \\ \mathcal{C}(GFX, GFX) & \xrightarrow{\varphi_{GFX, FX}} & \mathcal{D}(FGFX, FX) & & \\ \uparrow \mathcal{C}(\eta_x, \text{id}) & & \downarrow \mathcal{C}(F\eta_x, \text{id}) & & \\ \mathcal{C}(x, GFX) & \xrightarrow{\varphi_{x, FX}} & \mathcal{D}(FX, FX) & & \\ \uparrow \eta_x & & & & \\ x & & & & \end{array}$$

from which it follows that

$$\begin{array}{ccc} \text{id}_{GFX} & \xrightarrow{\varphi} & \varphi(\text{id}_{GFX}) \stackrel{\text{def. of } \varepsilon}{=} \varepsilon_{FX} \\ \downarrow \mathcal{C}(\eta_x, \text{id}) & & \downarrow \mathcal{C}(F\eta_x, \text{id}) \\ \eta_x = \varphi^{-1}(\varepsilon_{FX} \circ F\eta_x) & \xleftarrow{\varphi_{x, FX}^{-1}} & \varepsilon_{FX} \circ F\eta_x \end{array}$$

claim

Therefore we have $(LHS) = (RHS)$ in $(**)$.

The other eq. axiom is proved similarly.

2a) Let $(F, G, \eta, \varepsilon)$ form an adjunction and define

$$\varphi_{X,A} : \mathcal{C}(X, GA) \longrightarrow \mathcal{D}(FX, A)$$

by

$$\left(X \xrightarrow{h} GA \right) \longmapsto \left(FX \xrightarrow{Fh} FGA \xrightarrow{\varepsilon_A} A \right)$$

Then φ is natural in X and A .

Proof.

We have to show the following.

\mathcal{C}	\mathcal{D}	Sets	
X	A	$\mathcal{C}(X, GA)$	$\xrightarrow{\varphi} \mathcal{D}(FX, A)$
$\downarrow f$	$\downarrow g$	\downarrow	$\Downarrow \text{aim} \quad \downarrow \mathcal{D}(Ff, g)$
X'	A'	$\mathcal{C}(X', GA')$	$\xrightarrow{\varphi} \mathcal{D}(FX', A')$

That is, for each $h : X \rightarrow GA$,

$$\begin{array}{ccc}
 h & \longmapsto & \varepsilon_A \circ Fh \\
 \downarrow & & \downarrow \\
 & & g \circ \varepsilon_A \circ Fh \circ Ff \quad \text{aim} \\
 & & \parallel \\
 (Gg \circ h \circ f) & \longmapsto & (FX' \xrightarrow{F(Gg \circ h \circ f)} FGA' \xrightarrow{\varepsilon_{A'}} A')
 \end{array}$$

Now we have

$$\begin{array}{ccccc}
 FX' \xrightarrow{FF} FX & \xrightarrow{Fh} & FGA & \xrightarrow{\varepsilon_A} & A \\
 & & \downarrow FGg & \parallel & \downarrow g \\
 & & FGA' & \xrightarrow{\varepsilon_{A'}} & A'
 \end{array}$$

naturality
 $\downarrow \varepsilon$

Therefore

$$g \circ \varepsilon_A \circ Fh \circ FF = \varepsilon_{A'} \circ FGg \circ Fh \circ FF$$

and this proves the claim. \square

2b) In the above setting of **2a)**,
 $\varphi_{x,A}$ is an isomorphism.

Proof. We define

$$\psi_{x,A} : \mathcal{D}(FX, A) \longrightarrow \mathcal{C}(x, GA)$$

$$(FX \xrightarrow{i} A) \longmapsto (x \xrightarrow{\eta_x} GFX \xrightarrow{Gi} GA)$$

We shall show that $\varphi_{x,A}$ and $\psi_{x,A}$ are
 inverse to each other.

$$\boxed{\varphi \circ \psi = \text{id}}$$

$$(\varphi \circ \psi)(i) = \varphi(Gi \circ \eta_x)$$

$$= \left(FX \xrightarrow{F\eta_X} FGFX \xrightarrow{FGi} FGA \xrightarrow{\epsilon_A} A \right)$$

$\begin{array}{c} \nearrow \parallel \\ \epsilon_{FX} \\ \searrow \parallel \text{ nat. of } \epsilon \\ FX \xrightarrow{\quad} A \\ \nearrow i \end{array}$

$\begin{array}{c} \circlearrowleft \\ \text{eq. axiom} \\ \text{for } \eta, \epsilon \\ \searrow \text{id} \end{array}$

$$= i.$$

$$\boxed{\eta \circ \varphi = \text{id}}$$

Similar.

□