

Diagrammatic execution models for functional programming

Koko Muroya

(RIMS, Kyoto University
& University of Birmingham)

Steven W. T. Cheung

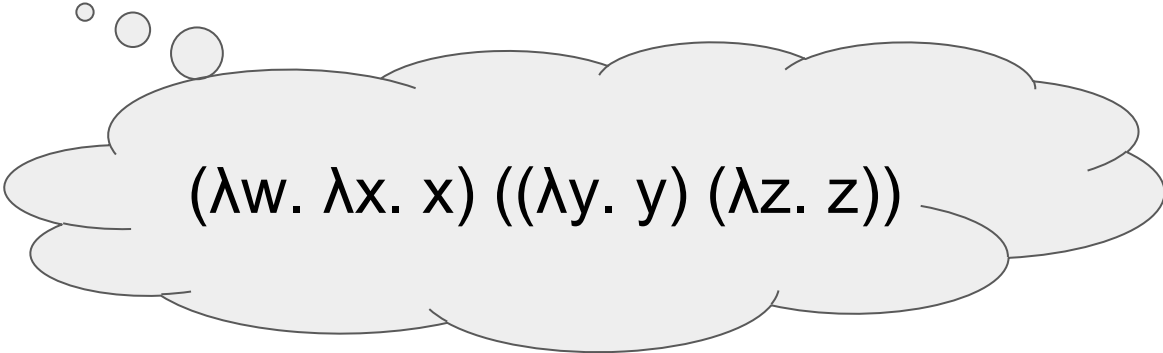
Dan R. Ghica
(University of Birmingham)



What happens when you run a functional program by hand?

Running a functional program by hand

```
(* computing the identity function in OCaml *)  
(fun _ -> fun x -> fun x) ((fun y -> y) (fun z -> z))  
  
-- computing the identity function in Haskell  
(\_ -> \x -> x) ((\y -> y) (\z -> z))
```




$(\lambda w. \lambda x. x) ((\lambda y. y) (\lambda z. z))$

Running a functional program by hand

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(*      (fun _ -> fun x -> x) ((fun y -> y) (fun z -> z)) *)  
(* => (fun _ -> fun x -> x) (fun z -> z) *)  
(* => fun x -> x *)
```

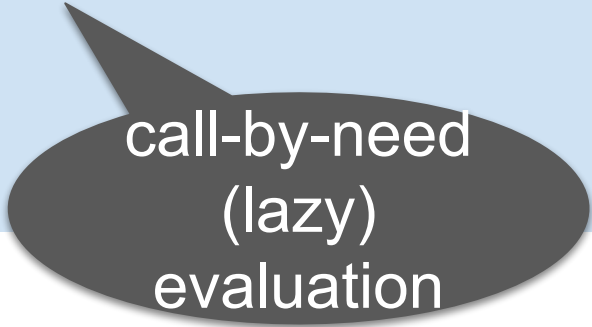


call-by-value
evaluation

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(\_ -> \x -> x) ((\y -> y) (\z -> z))

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call-by-need
(lazy)
evaluation

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call-by-need
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evaluation

***Same result,
but different steps!***

Let's run a functional program with a bit
of formality...

Running a functional program, formally

by change of configuration

- program
- “focus”

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\x -> x
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evaluation

Running a functional program, formally

by change of configuration

- program
- “focus” *determining evaluation strategy*

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Goodness of formality

prove program equivalence (validate compiler optimisations)

- “When are two programs *the same*?”
- result modelled by observable transitions & final configuration

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- “How *much time/space* does it take to run a program?”
- time cost modelled by number and cost of transitions
- space cost modelled by size of configurations

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guarantee “correctness” of implementation

- “Does a compiler work *as intended*?”
- implementation derived as abstract machine

Now we use *diagrams* to run a functional program with formality!

Conventional approaches

running a program by
change of configuration

- program
- “focus” determining evaluation strategy

Diagrammatic approach

running a program by
change of diagram configuration

- *diagram representation* of program
- “token” determining evaluation strategy

Running a functional program with *diagrams*

by change of *diagram* configuration

name-free

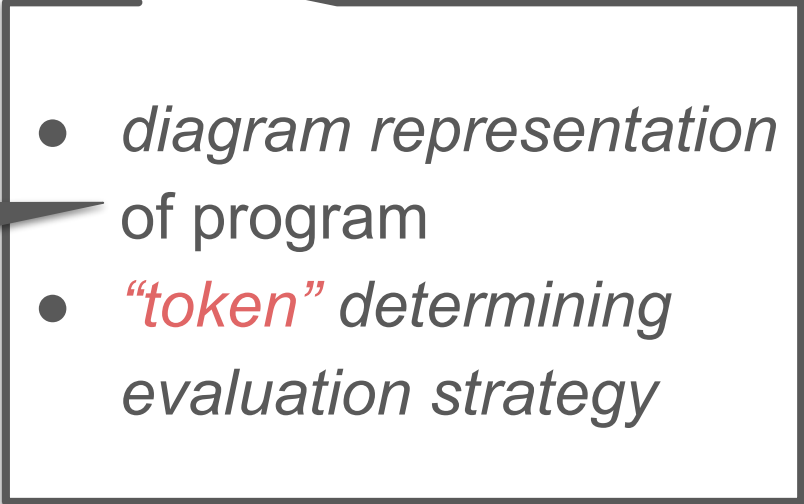
- *diagram* representation of program
- “*token*” determining *evaluation strategy*

Running a functional program with *diagrams*

by change of *diagram* configuration



name-free

- 
- *diagram* representation of program
 - “*token*” determining evaluation strategy

[DEMO]

On-line visualiser for lambda-calculus

<https://koko-m.github.io/GoI-Visualiser/>

- call-by-name, call-by-need (lazy)
- call-by-value: left-to-right, right-to-left

Goodness of *diagrams*

- name-free
- environment included

Goodness of *diagrams*

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visualise interesting properties of programs

- call-by-value vs. call-by-need (lazy)
- on-demand copying with intermediate sharing
- patterns in divergence
 - $(\lambda x. x x) (\lambda x. x x)$
 - $(\lambda x. x x x) (\lambda x. x x x)$

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answer (conventional) questions from new perspectives

- “When are two programs *the same*?”
- “How *much time/space* does it take to run a program?”
- “Does a compiler work *as intended*?”

More goodness of *diagrams*

More goodness of *diagrams*

support textual-and-visual programming

guide language design for unconventional/new programming paradigms (to be presented by Steven)

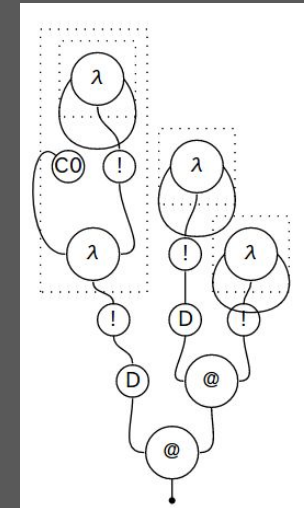
Towards textual-and-visual programming

“We’d like not just text or diagram,
but both!”

textual program

```
(λw. λx. x) ((λy. y) (λz. z))
```

diagrammatical
program



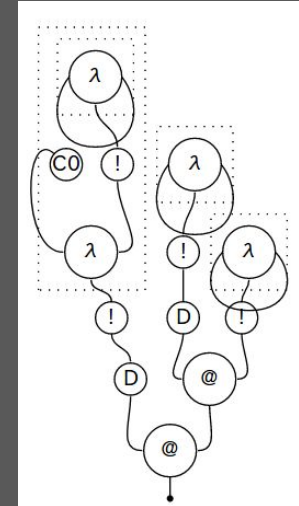
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diagrammatical
program



✓
(grammar)

form / edit

?
(validity criteria)

✓

execute

✓

✓

debug

□ !

OCaml Visual Debugger

<https://fyp.jackhughesweb.com/> by Jack Hughes

for a subset of OCaml

- arithmetic (int), comparison (bool)
- conditional (if), recursion (let rec)
- lists, pairs
- ~~pattern matching~~

features

- interactive diagram view
- go forwards/backwards & pause/resume & jump steps
- breakpoint on diagram
- stats

OCaml Visual Debugger

visualise interesting properties of programs

- sorting algorithms comparison
 - bubble-sort vs. insert-sort
<https://www.youtube.com/watch?v=bZMSwo0zLio&t=130s>
 - merge-sort vs. insert-sort
<https://www.youtube.com/watch?v=U1NI-mWeNe0>
- tail-recursion vs. non tail-recursion
<https://www.youtube.com/watch?v=R4yCV5Ts1gk&t=14s>

More goodness of *diagrams*

support textual-and-visual programming

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- We’ve got OCaml Visual Debugger
- ... and want a text-and-diagram editor!

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